

Advanced Database Technology

February 13, 2003

DATA STORAGE

(Lecture based on [GUW 11.2–11.5], [Sanders03, 1.3–1.5],

and [MaheshwariZeh03, 3.1–3.2])

Slides based on

Notes 02: Hardware

for Stanford CS 245, fall 2002

by Hector Garcia–Molina

Today

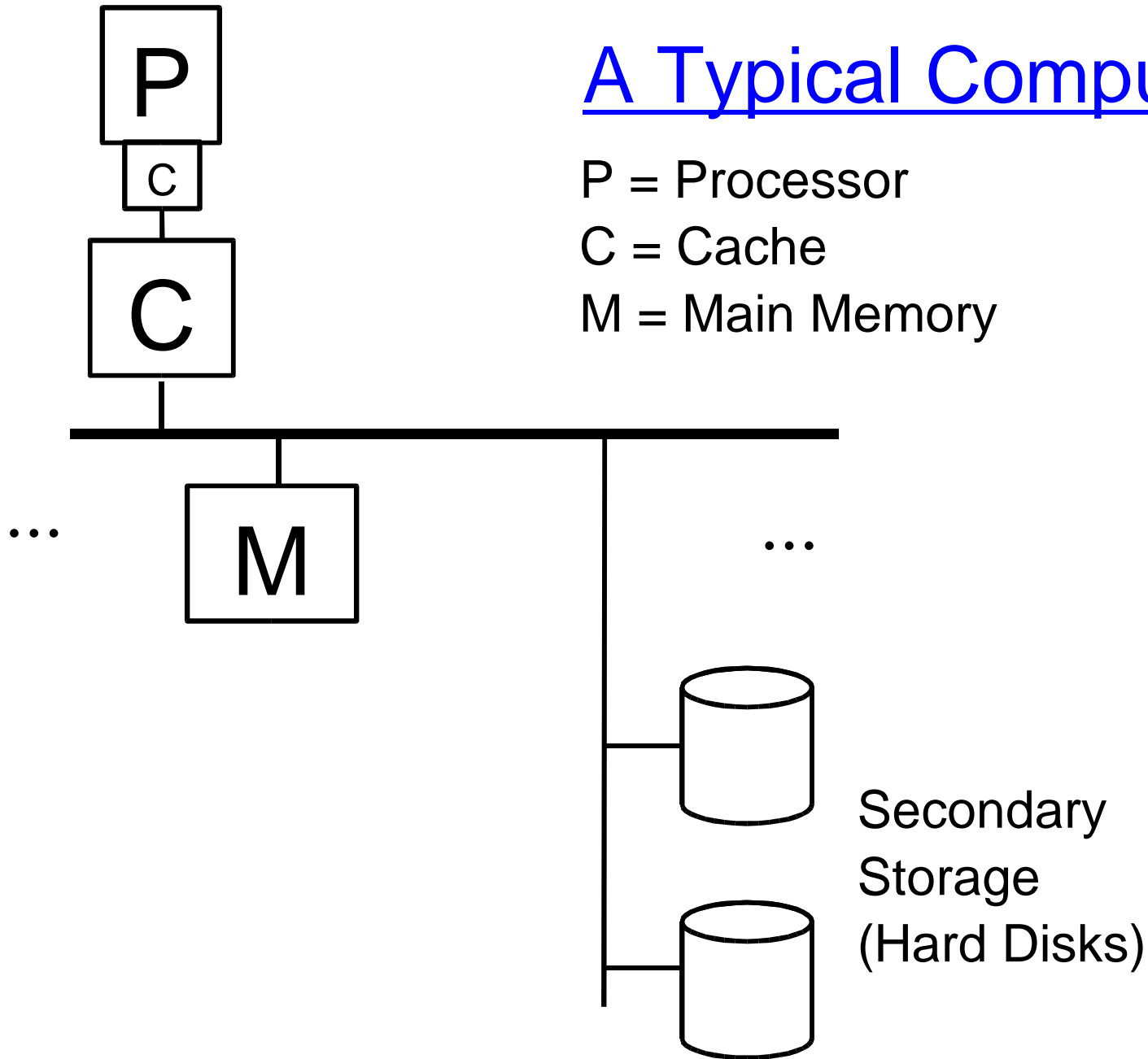
- Hardware, external memory
- Disk access times
- The I/O model
- Case study: Sorting
- Lower bound for sorting
- Optimizing disk usage

A Typical Computer

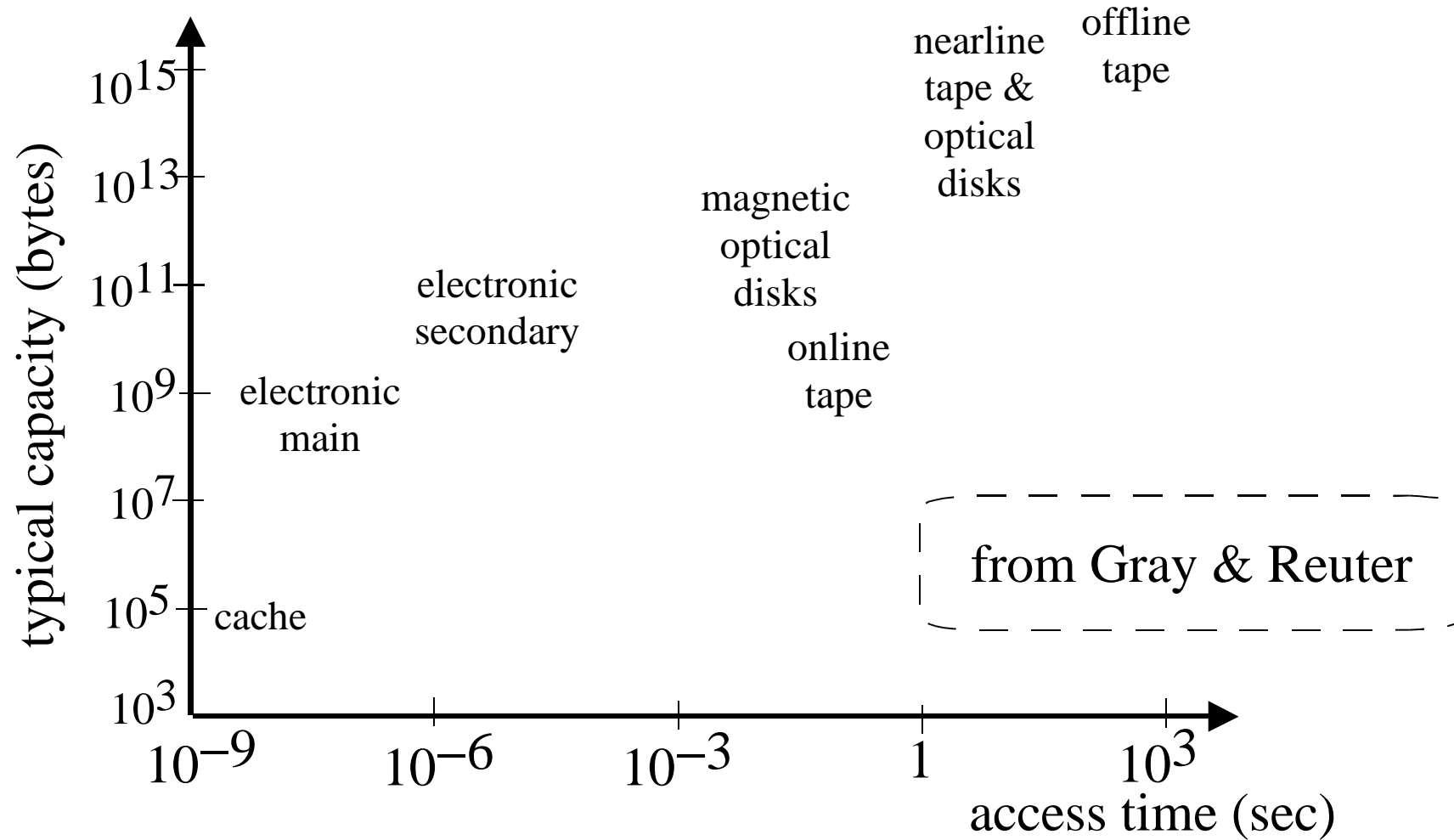
P = Processor

C = Cache

M = Main Memory

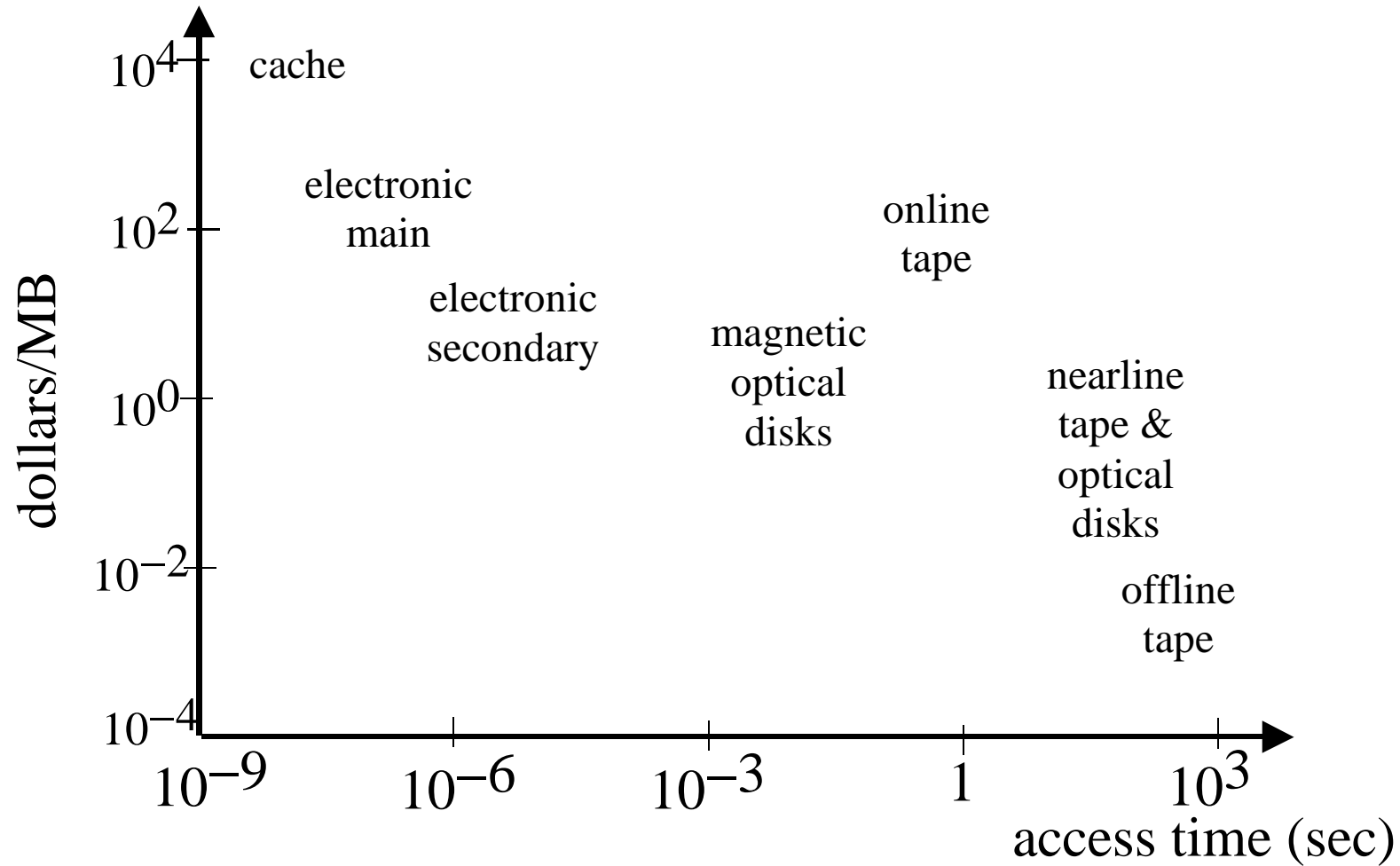


Storage Capacity



Storage Cost

from Gray & Reuter



Caching

Cache: Fast memory, holding frequently used parts of a slower, larger memory.

Examples:

- A small (L1) cache holds a few kilobytes of the memory "most recently used" by the processor.
- Most operating systems keep the most recently used "pages" of memory in main memory and put the rest on disk.

Virtual memory

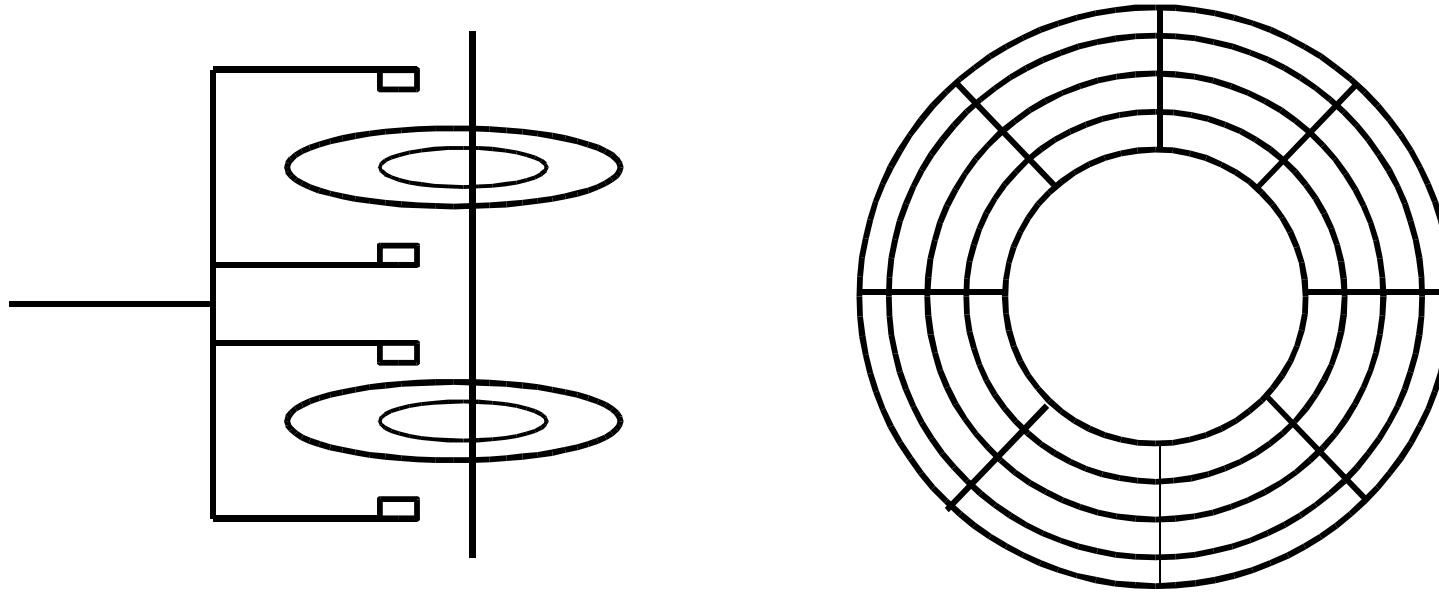
- In most operating systems, programs don't know if they access main memory or a page that resides on secondary memory.
- This is called **virtual memory** (the book is a little fuzzy on this).
- Database systems usually take explicit control over secondary memory accesses.

Secondary storage

Many flavors:

- Disk: Floppy (hard, soft)
Winchester
Ram disks
Optical, CD-ROM...
- Tape Reel, cartridge
Robots

Focus on: “Typical Disk”



- Terms: – Platter, Head,
Cylinder, Track
Sector (all physical).
– Block (logical).

“Typical” Numbers

Diameter: 1 inch → 15 inches

Cylinders: 40 (floppy) → 20000

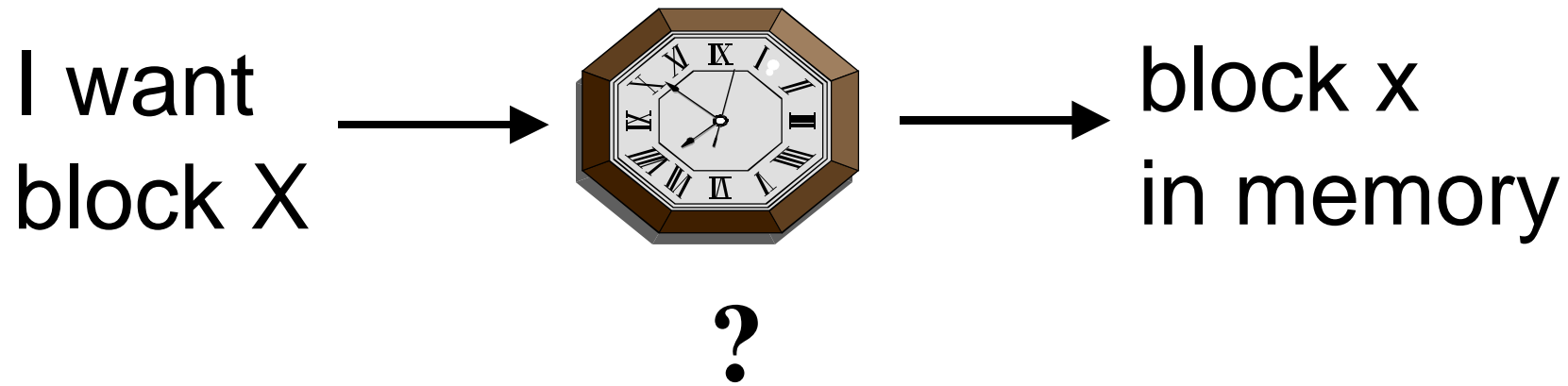
Surfaces: 1 (CDs) →

2 (floppies) → 30

Sector Size: 512B → 50K

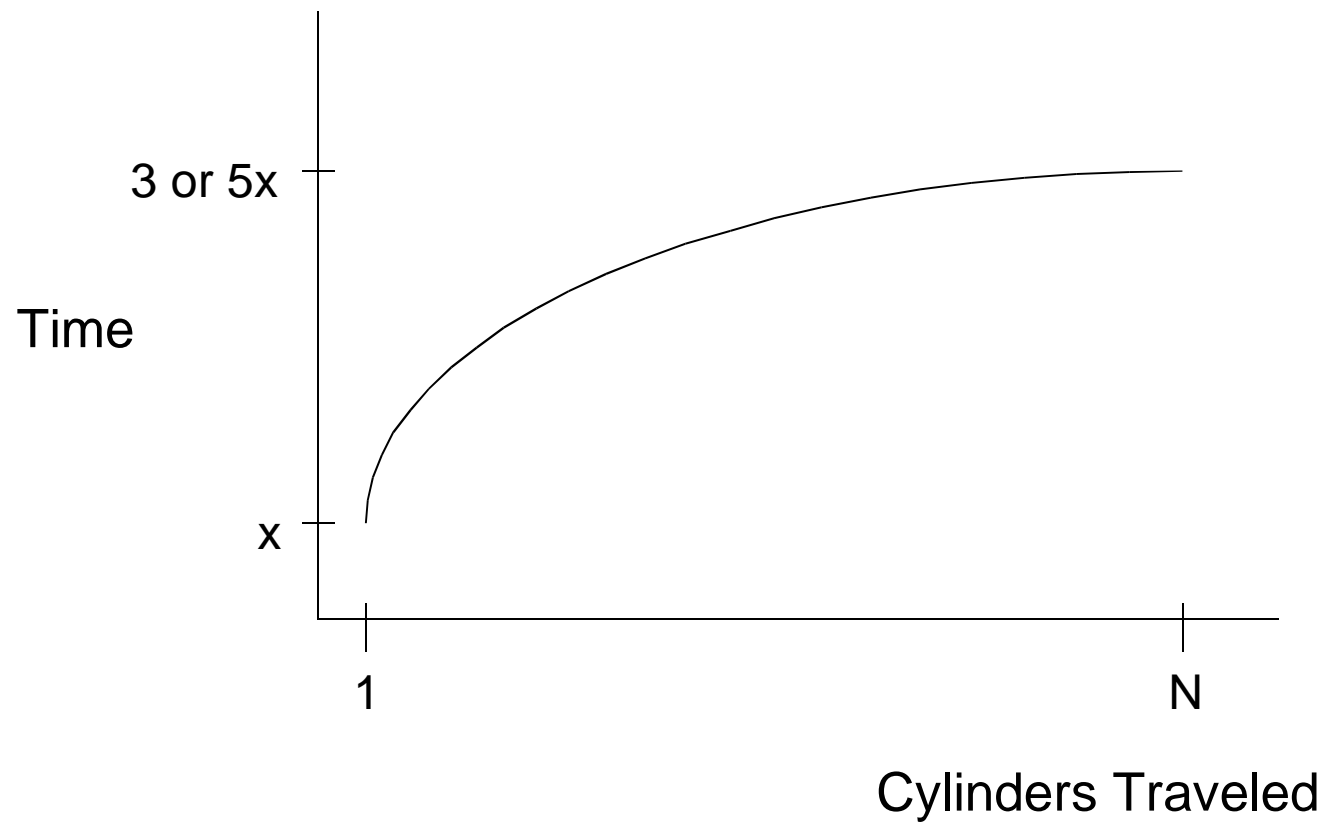
Capacity: 360 KB (old floppy)
→ 30 GB (I use)

Disk Access Time



Time = Seek Time +
Rotational Delay +
Transfer Time +
Other

Seek Time

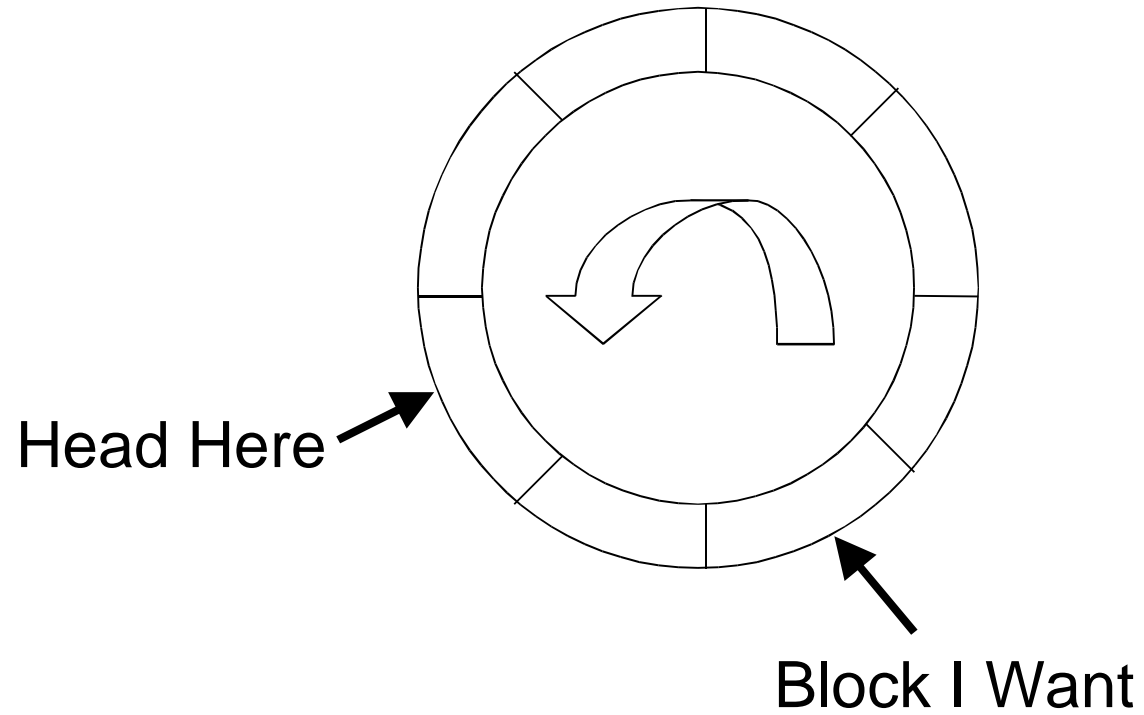


Average Random Seek Time

$$S = \frac{\sum_{i=1}^N \sum_{\substack{j=1 \\ j \neq i}}^N \text{SEEKTIME } (i \rightarrow j)}{N(N-1)}$$

“Typical” S: 10 ms → 40 ms
10s of millions clock cycles!

Rotational Delay



Average Rotational Delay

R = 1/2 revolution

“typical” R = 4.16 ms (7200 RPM)

Transfer Rate: t

- “typical” t : 10 → 50 MB/second
- transfer time: $\frac{\text{block size}}{t}$
- E.g., block size 32 kb, $t=32$ MB/second gives transfer time 1 ms.

Other Delays

- CPU time to issue I/O
- Contention for controller
- Contention for bus, memory

“Typical” Value: 0

- **So far:** Random Block Access
- **What about:** Reading “Next” block?

If we don't need to change cylinder

$$\text{Time to get block} = \frac{\text{Block Size}}{t} + \text{Negligible}$$



- switch track
- once in a while,
next cylinder

Rule of Thumb	Random I/O: Expensive
	Sequential I/O: Less expensive

- Ex: 1 KB Block
 - » Random I/O: ~ 20 ms.
 - » Sequential I/O: ~ 1 ms.

However, the relative difference is smaller if we use larger blocks.

Cost for Writing similar to Reading

.... unless we want to verify!

need to add (full) rotation + $\frac{\text{Block size}}{t}$

To Modify a Block?

- (a) Read Block
- (b) Modify in Memory
- (c) Write Block
- [(d) Verify?]

Problem session

We usually express the running time of algorithms as the number of operations.

- Give an example of algorithms (using external memory) where this is misleading.
- What should we count instead?
- How do we do the counting?

The I/O model of computation

- Count the number of disk blocks read or written by an algorithm (I/Os).
- Ignore the number of operations! (?)
- Explicit control of which blocks are in main memory.
- **Notation:** M = size of main memory
 B = size of disk blocks
 N = size of data set

Problem session

Name a few sorting algorithms that you know:

- What are the (worst case) running times in internal memory?
- What is the number of I/Os if we use the algorithm on external memory
 - if no caching is done?
 - when storing the M/B most recently accessed blocks in main memory?

External memory sorting

Let's see if we can do better..!

(Perhaps as well as in [MaheshwariZeh03, 3.2])

Sorting in GUW

- More practical viewpoint: Two passes over the data enough unless data is huge.
- **TPMMS** = Two-Pass Multiway MergeSort.
- More general treatment in [MaheshwariZeh03, 3.2]

External memory sorting

Let's see if we did as well as we possibly could:

**A lower bound on the number of I/Os
for sorting**

(based on [Sanders03, 1.5])

Problem session

Consider the sorting problem in example 11.7 (page 527 in GUW):

- What are the values of N , B , and M ?
- How many I/Os will our sorting algorithm use on this problem?
- What lower bound does the formula from [Sanders03, Theorem 1] give?

Optimizations

Some practically oriented ways of speeding up external memory algorithms:

- **Disk Scheduling**
 - e.g., using the "elevator algorithm"
 - reduces average seek time when there are multiple simultaneous "unpredictable" requests
- **Track buffer / cylinder buffer**
 - good when data are arranged and accessed "by cylinder"

- **Pre-fetching**
 - Speeds up access when the needed blocks are known, but the order of requests is data dependent
- **Disk arrays**
 - Increases the rate at which data can be transferred
 - Striping: Blocks from each disk are grouped to form a large logical block
- **Mirrored disks**
 - Same speed for writing
 - Several blocks can be read in parallel

Block Size Selection?

- Big Block → Amortize I/O Cost



Unfortunately...

- Big Block ⇒ Read in more useless stuff!
And takes longer to read
- **Trend:** Blocks get bigger

Problem session

Which of the mentioned optimizations apply to (parts of) the sorting algorithm we developed earlier?

- Disk scheduling
- Cylinder buffer
- Pre-feeding
- Disk arrays
- Mirrored disks

Summary

- External memories are complicated.
- Essential features captured by the I/O model (to be used henceforth).
- We saw matching I/O upper and lower bounds for sorting.
- A bit (and the last bit) about optimizing constant factors in external memory access.