

A Concurrent Logical Relation

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Abstract—We present a logical relation for showing the correctness of program transformations based on a new type-and-effect system for a concurrent extension of an ML-like language with higher-order functions, higher-order store and dynamic memory allocation.

We show how to use our model to verify a number of interesting program transformations that rely on effect annotations. In particular, we prove a Parallelization Theorem, which expresses when it is sound to run two expressions in parallel instead of sequentially. The conditions are expressed solely in terms of the types and effects of the expressions. To the best of our knowledge, this is the first such result for a concurrent higher-order language with higher-order store and dynamic memory allocation.

I. INTRODUCTION

Relational reasoning about program equivalence is useful for reasoning about the correctness of program transformations, data abstraction (representation independence), compiler correctness, etc. The standard notion of program equivalence is contextual equivalence and in recent years, there has been many improvements in reasoning methods for higher-order ML-like languages with general references, based on bisimulations, e.g., [1], [2], [3], traces [4], game semantics [5], and Kripke logical relations, e.g., [6], [7], [8], [9].

So far, these methods have focused on sequential languages. In this paper we present the first Kripke logical relation for reasoning about equivalence of a *concurrent* higher-order ML-like language with higher-order store and dynamic memory allocation.

To state and prove useful equivalences about concurrent programs, it is necessary to have some way of restricting the contexts under which one proves equivalences. This point was made convincingly in the recent paper by Liang et. al. [10], who presented a rely-guarantee-based simulation for verifying concurrent program transformations for a first-order imperative language (with first-order store). Here is a very simple example illustrating the point. Consider two expressions

$$e_1 \equiv x := 1; y := 1 \quad \text{and} \quad e_2 \equiv y := 1; x := 1.$$

Here x and y are variables of type refint . The expressions e_1 and e_2 are not contextually equivalent. (To see why, consider expression $e_3 \equiv x := 0; y := 0$, and note that running e_1 in parallel with e_3 may result in a state with $!x = 0$ and $!y = 1$, but that cannot be the case when we run e_2 in parallel with e_3 .) The issue is, of course, that the context may also modify the references x and y . On the other hand, if we know that no other threads have access to x or y , then it should be the case

that e_1 and e_2 are equivalent. We can express this restriction on the contexts using a refined region-based type-and-effect system.

We first recall that a type-and-effect system is a type system that classifies programs according to which side effects the programs may have. A variety of effect systems have been proposed for higher-order programming languages, e.g., [11], [12], [13], see [14] for a recent overview. Effect systems can often be understood as specifying the results of a static analysis, in the sense that it is possible to automatically infer types and effects. Effect systems can be used for different purposes: they were originally proposed by Lucassen and Gifford [12] for parallelization purposes but they have also, e.g., been used as the basis for implementing ML using a stack of regions for memory management [13], [15]. In a recent series of papers, Benton et. al. have argued that another important point of effect systems is that they can be used as the basis for effect-based program transformations, e.g., compiler optimizations, [16], [17], [18], [19], see also [20]. The idea is that certain program transformations are only sound under additional assumptions about which effects program phrases may, or rather may not, have.

Now, returning to our example, we refine the types of x and y to be $\text{ref}_\rho \text{int}$ and $\text{ref}_\sigma \text{int}$, respectively. Intuitively, this expresses that x and y are references in different regions, but it does not put any restrictions on whether other threads may access x or y . Thus, when we type e_1 and e_2 we will, moreover, use two contexts of region variables, one for public regions that can be used by other concurrently running threads, and one for private regions that are under the control of the present thread. This idea is inspired by recent work on concurrent separation logic, e.g., [21], [22], [23], [24]. We use a vertical bar to separate public and private regions: the typing context

$$\rho, \sigma \mid \emptyset \mid x : \text{ref}_\rho \text{int}, y : \text{ref}_\sigma \text{int}$$

expresses that ρ and σ are public regions, whereas the typing context

$$\emptyset \mid \rho, \sigma \mid x : \text{ref}_\rho \text{int}, y : \text{ref}_\sigma \text{int}$$

expresses that ρ and σ are private regions. The expressions e_1 and e_2 are well-typed in the latter context and, with this refined typing, they are indeed contextually equivalent, because our type-and-effect system guarantees that no well-typed context can access regions ρ or σ . (The expressions are also well-typed

in the former context, but not contextually equivalent with that refined typing.)

In this paper we present a step-indexed Kripke logical relations model of a type-and-effect system with public and private regions for a concurrent higher-order language with general references. Our model is constructed over the operational semantics of the programming language, and builds on recent work by Thamsborg and Birkedal on logical relations for the sequential sub-language [20]. Note that the type-and-effect annotations are just annotations; the operational semantics of the language is standard and regions only exist in our semantic model, not in the operational semantics.

As an important application of our model we prove a Parallelization Theorem, which expresses when it is sound to run two expressions in parallel instead of sequentially. To the best of our knowledge, this is the first such result for a higher-order language with higher-order store and dynamic memory allocation. Here is a very simple instance of the theorem. Consider two expressions

$$\begin{aligned} e_1 &\equiv y := !x + !y \\ e_2 &\equiv z := !x + !z, \end{aligned}$$

each well-typed in a context

$$\emptyset \mid \rho_x, \rho_y, \rho_z \mid x : \text{ref}_{\rho_x} \text{int}, y : \text{ref}_{\rho_y} \text{int}, z : \text{ref}_{\rho_z} \text{int},$$

i.e., where x , y , and z are references in distinct private regions. In this context, running e_1 and e_2 sequentially is contextually equivalent to running e_1 and e_2 in parallel. Intuitively, this also makes sense: e_1 and e_2 update references in distinct regions, and it is unproblematic that they both read (but not write) from the same region.

As mentioned, this was a simple instance of the Parallelization Theorem. We stress that the theorem is expressed solely in terms of the type and effects of the expressions e_1 and e_2 , so a compiler may automatically infer that it is safe to parallelize two expressions by looking at the inferred effect types, and without reasoning about all interleavings. Moreover, the theorem applies to contexts and expressions with general higher types (not just with references to integers and unit types). Note that the distinction between private and public regions is also crucial here (parallelization would not be sound if the effects of the expressions were on public regions).

Our type-and-effect system crucially also includes a region-masking rule. Traditionally, this rule has been used to hide local effects on regions, which makes it possible to view a computation as pure even if it uses effects locally and makes the effect system stronger, in the sense that it can justify more program transformations. Here we also observe that the masking rule can be used for introducing private regions, since the masking rule intuitively guarantees that effects on a region are not leaked to the context. It is well-known that region-masking makes the model construction for a sequential language technically challenging, see the extensive discussion in [20]. Here it is yet more challenging because of concurrency; we explain how our model ensures soundness of the masking rule in Section III.

The extension with concurrency also means that when we define the logical relation for contextual approximation and relate two computations e_1 and e_2 , then we cannot simply require relatedness after e_1 has completed evaluation (as in the sequential case), since other threads should be allowed to execute as well. We explain our approach to relating concurrent computations in Section III; it is informed by recent soundness proofs of unary models of concurrent separation logic [25], [26].

Another challenge arises from the fact that since our language includes dynamically allocated general references, the existence of the logical relation is non-trivial; in particular, the set of Kripke worlds must be recursively defined. Here we build on our earlier work [27] and define the worlds as a solution to a recursive metric-space equation. Indeed, to focus on the essential new aspects due to the extension with concurrency, we deliberately choose to use the exact same notion of worlds as we used for the sequential sub-language in [20]. In the same vein, we here consider a monomorphically typed higher-order programming language with general references, but leave out universal and existential types as well as recursive types. However, we want to stress that since our semantic techniques (step-indexed Kripke logical relations over recursively defined worlds) do indeed scale well to universal, existential, and recursive types, e.g. [27], [9], it is possible to extend our model to a language with such types. We conjecture that it is also possible to extend our model to richer effect systems involving region and effect polymorphism, but we have not done so yet.

II. LANGUAGE AND TYPING

We consider a standard call-by-value lambda calculus with general references, and extended with parallel composition and an atomic construct. We assume countably infinite, pairwise disjoint sets of *region variables* \mathcal{RV} (ranged over by ρ), *locations* \mathcal{L} (ranged over by l) and program variables (ranged over by x, y, f). As usual, the reduction relation is between configurations, $(e \mid h) \mapsto (e' \mid h')$ where heaps \mathcal{H} are finite maps from locations to values. Figures 1 and 2 give the syntax and operational semantics; we denote the set of expressions \mathcal{E} and the set of values \mathcal{V} . The evaluation contexts allow parallel evaluation inside par expressions, and there is a new primitive reduction covering the case when the two subcomputations have terminated. For technical simplicity, we allow an atomic e expression to reduce to itself, possibly introducing more divergence than the diverging behaviours of e . The syntax is kept minimal; in examples we may use additional syntactic sugar, e.g., writing $\text{let } x = e_1 \text{ in } e_2$ for $(\text{fun } f(x).e_2) e_1$ for some fresh f . For $e \in \mathcal{E}$, we write $\text{FV}(e)$ and $\text{FRV}(e)$ for the sets of free program variables and region variables, respectively; also we define $\text{rds } \varepsilon = \{\rho \in \mathcal{RV} \mid \text{rd}_\rho \in \varepsilon\}$ and similarly for writes and allocation.

The form of the judgments of our type-and-effect system is standard with one important refinement: regions are partitioned into *public* and *private* regions, with the purpose of restricting interference from the environment. In greater detail, a typing

$$\begin{aligned}
\pi &::= rd_\rho \mid wr_\rho \mid al_\rho \\
\varepsilon &::= \pi_1, \dots, \pi_n \\
\tau &::= \mathbf{1} \mid \text{int} \mid \tau_1 \times \tau_2 \mid \tau_1 \xrightarrow{\Pi, \Lambda} \tau_2 \mid \text{ref}_\rho \tau \\
v &::= x \mid \langle \rangle \mid \langle v_1, v_2 \rangle \mid \text{fun } f(x).e \mid l \\
e &::= v \mid \text{proj}_i v \mid v e \mid \text{ref } v \mid v_1 := v_2 \mid !v \mid \text{par } e_1 \text{ and } e_2 \\
&\quad \mid \text{cas } (v_1, v_2, v_3) \mid \text{atomic } e \\
E &::= [] \mid v E \mid \text{par } E \text{ and } e_2 \mid \text{par } e_1 \text{ and } E
\end{aligned}$$

Fig. 1. Syntax

$$\begin{aligned}
(E[\text{proj}_i \langle v_1, v_2 \rangle] \mid h) &\longmapsto (E[v_i] \mid h) \\
(E[(\text{fun } f(x).e) v] \mid h) &\longmapsto (E[e[f := \text{fun } f(x).e, x := v]] \mid h) \\
(E[\text{ref } v] \mid h) &\longmapsto (E[l \mid h[l \mapsto v]] \mid h) \quad \text{if } l \notin \text{dom}(h) \\
(E[l := v] \mid h) &\longmapsto (E[\langle \rangle \mid h[l := v]] \mid h) \quad \text{if } l \in \text{dom}(h) \\
(E[!l] \mid h) &\longmapsto (E[h(l)] \mid h) \quad \text{if } l \in \text{dom}(h) \\
(E[\text{par } v_1 \text{ and } v_2] \mid h) &\longmapsto (E[\langle v_1, v_2 \rangle] \mid h) \\
(E[\text{cas } (l, n_1, n_2)] \mid h) &\longmapsto (E[\mathbf{1} \mid h[l := n_2]] \mid h) \\
&\quad \text{if } l \in \text{dom}(h) \text{ and } h(l) = n_1 \\
(E[\text{cas } (l, n_1, n_2)] \mid h) &\longmapsto (E[\mathbf{0}] \mid h) \\
&\quad \text{if } l \in \text{dom}(h) \text{ and } h(l) \neq n_1 \\
(E[\text{atomic } e] \mid h) &\longmapsto (E[v] \mid h') \\
&\quad \text{if } (e \mid h) \mapsto^* (v \mid h') \\
(E[\text{atomic } e] \mid h) &\longmapsto (E[\text{atomic } e] \mid h)
\end{aligned}$$

Fig. 2. Operational semantics

judgement looks like this:

$$\Pi \mid \Lambda \mid \Gamma \vdash e : \tau, \varepsilon.$$

The Γ , e and τ are the usual: the *variable context* Γ assigns types to program variables; if you substitute properly typed values for variables in the expression e and run the resulting expression, you will get a value of type τ or possibly diverge. To get an idea of — or rather an upper bound of — the side-effects of e , we furthermore split the heap into *regions*; these are listed in Π and Γ . And we track memory accesses by adding a set ε of *effects* of the form rd_ρ , wr_ρ and al_ρ , where ρ is a region. Roughly, a computation with effect rd_ρ may read one or more locations in region ρ , and similarly for writes and allocation. This setup goes back to Lucassen and Gifford [12].

The novelty, as mentioned in the Introduction, is our partition of regions into the *public* ones Π and the *private* ones Λ . As opposed to the rest of the judgement, this public-private division does not make promises about the behavior of e . Instead, it states the expectations that e has of the environment: threads running in parallel with e may — in a well-typed manner — read, write and allocate in the public regions but must leave the private regions untouched.

When running parallel threads, the private regions of the parent are shared between the children, and so are public from their point of view; this is reflected in the typing rule for

parallel composition:

$$\frac{\Pi, \Lambda \mid \cdot \mid \Gamma \vdash e_1 : \tau_1, \varepsilon_1 \quad \Pi, \Lambda \mid \cdot \mid \Gamma \vdash e_2 : \tau_2, \varepsilon_2}{\Pi \mid \Lambda \mid \Gamma \vdash \text{par } e_1 \text{ and } e_2 : \tau_1 \times \tau_2, \varepsilon_1 \cup \varepsilon_2}.$$

Note that the parent thread only continues once both children have terminated; as a consequence, the parent regains ownership of its private regions before it goes on. Running an expression atomically temporarily makes all regions private:

$$\frac{\cdot \mid \Pi, \Lambda \mid \Gamma \vdash e : \tau, \varepsilon}{\Pi \mid \Lambda \mid \Gamma \vdash \text{atomic } e : \tau, \varepsilon} \quad (\text{als } \varepsilon \subseteq \text{rds } \varepsilon \cap \text{wrs } \varepsilon).$$

The side condition is a technical necessity. Finally, new, private regions are introduced by the so-called masking rule:

$$\frac{\Pi \mid \Lambda, \rho \mid \Gamma \vdash e : \tau, \varepsilon}{\Pi \mid \Lambda \mid \Gamma \vdash e : \tau, \varepsilon - \rho} \quad (\rho \notin \text{FRV}(\Gamma, \tau))$$

The subtraction of ρ in the conclusion removes any read, write or allocation effects tagged with ρ . The reading of the masking rule is that we make a brand new, empty region ρ for e to use, but once e has terminated we forget about ρ again; this works out since the side condition prevents e from leaking locations from ρ . Traditionally, the masking rule has been used to do memory-management [13] as well as a means of hiding local effects to facilitate effect-based program transformations [17], [20]. Here we make another use of the rule: we observe that, moreover, e cannot leak locations from ρ while running and so ρ is a *private* region for the duration of e . After all, the only means of inter-thread communication is shared memory.

All the typing rules are in Figure 3. We just remark here, that reference types are tagged with the region where the location resides and that function arrows are tagged with the latent effects as well as with the public and private regions that the function expects; the latter is natural once we remember that a function is basically just a suspended, well-typed expression.

Because of the nondeterminism arising from *par* and shared references, the definition of contextual equivalence could take into account both may- and must-convergence. In this paper we only consider may-equivalence and formally we define (may-) contextual approximation by:

Definition 1. $\Pi \mid \Lambda \mid \Gamma \vdash e \lesssim_\downarrow e' : \tau, \varepsilon$ if and only if for all h and C typed such that $\cdot \mid \cdot \mid \cdot \vdash C[e], C[e'] : \text{int}, \emptyset$, whenever $(C[e] \mid h) \downarrow$ then $(C[e'] \mid h) \downarrow$.

Here, as usual, $(e \mid h) \downarrow$ means that $(e \mid h) \mapsto^* (v \mid h')$ for some value v and some h' .

Contextual equivalence, $\Pi \mid \Lambda \mid \Gamma \vdash e \approx e' : \tau, \varepsilon$, is then defined as $\Pi \mid \Lambda \mid \Gamma \vdash e \lesssim_\downarrow e' : \tau, \varepsilon$ and $\Pi \mid \Lambda \mid \Gamma \vdash e' \lesssim_\downarrow e : \tau, \varepsilon$.

Note that the diverging behaviours introduced by our operational semantics of atomic e do not influence may-contextual equivalence. (The notion of may-contextual equivalence would be the same if we had omitted the last reduction rule in Figure 2.)

$$\begin{array}{c}
\frac{}{\Pi | \Lambda | \Gamma, x : \tau \vdash x : \tau, \emptyset} \quad \frac{}{\Pi | \Lambda | \Gamma \vdash \langle \rangle : \mathbf{1}, \emptyset} \quad \frac{\Pi | \Lambda | \Gamma \vdash v_1 : \tau_1, \varepsilon_1 \quad \Pi | \Lambda | \Gamma \vdash v_2 : \tau_2, \varepsilon_2}{\Pi | \Lambda | \Gamma \vdash \langle v_1, v_2 \rangle : \tau_1 \times \tau_2, \varepsilon_1 \cup \varepsilon_2} \\
\frac{\Pi | \Lambda | \Gamma \vdash v : \tau_1 \times \tau_2, \varepsilon}{\Pi | \Lambda | \Gamma \vdash \text{proj}_i v : \tau_i, \varepsilon} \quad \frac{\Pi | \Lambda | \Gamma, f : \tau_1 \xrightarrow{\Pi, \Lambda} \tau_2, x : \tau_1 \vdash e : \tau_2, \varepsilon}{\Pi | \Lambda | \Gamma \vdash \text{fun } f(x).e : \tau_1 \xrightarrow{\Pi, \Lambda} \tau_2, \emptyset} \quad \frac{\Pi | \Lambda | \Gamma \vdash v : \tau_1 \xrightarrow{\Pi, \Lambda} \tau_2, \varepsilon_1 \quad \Pi | \Lambda | \Gamma \vdash e : \tau_1, \varepsilon_2}{\Pi | \Lambda | \Gamma \vdash v e : \tau_2, \varepsilon_1 \cup \varepsilon_2 \cup \varepsilon} \\
\frac{\Pi | \Lambda | \Gamma \vdash v : \tau, \varepsilon \quad \rho \in \Pi, \Lambda}{\Pi | \Lambda | \Gamma \vdash \text{ref } v : \text{ref}_\rho \tau, \varepsilon \cup \{al_\rho\}} \quad \frac{\Pi | \Lambda | \Gamma \vdash v_1 : \text{ref}_\rho \tau, \varepsilon_1 \quad \Pi | \Lambda | \Gamma \vdash v_2 : \tau, \varepsilon_2}{\Pi | \Lambda | \Gamma \vdash v_1 := v_2 : \mathbf{1}, \varepsilon_1 \cup \varepsilon_2 \cup \{wr_\rho\}} \quad \frac{\Pi | \Lambda | \Gamma \vdash v : \text{ref}_\rho \tau, \varepsilon}{\Pi | \Lambda | \Gamma \vdash !v : \tau, \varepsilon \cup \{rd_\rho\}} \\
\frac{\Pi, \Lambda | \cdot | \Gamma \vdash e_1 : \tau_1, \varepsilon_1 \quad \Pi, \Lambda | \cdot | \Gamma \vdash e_2 : \tau_2, \varepsilon_2}{\Pi | \Lambda | \Gamma \vdash \text{par } e_1 \text{ and } e_2 : \tau_1 \times \tau_2, \varepsilon_1 \cup \varepsilon_2} \quad \frac{\cdot | \Pi, \Lambda | \Gamma \vdash e : \tau, \varepsilon}{\Pi | \Lambda | \Gamma \vdash \text{atomic } e : \tau, \varepsilon} \quad (\text{als } \varepsilon \subseteq \text{rds } \varepsilon \cap \text{wrs } \varepsilon) \\
\frac{\Pi | \Lambda, \rho | \Gamma \vdash e : \tau, \varepsilon \quad (\rho \notin \text{FRV}(\Gamma, \tau))}{\Pi | \Lambda | \Gamma \vdash e : \tau, \varepsilon - \rho} \quad \frac{\Pi | \Lambda | \Gamma \vdash v_1 : \text{ref}_\rho \text{int}, \varepsilon_1 \quad \Pi | \Lambda | \Gamma \vdash v_2 : \text{int}, \varepsilon_2 \quad \Pi | \Lambda | \Gamma \vdash v_3 : \text{int}, \varepsilon_3}{\Pi | \Lambda | \Gamma \vdash \text{cas } (v_1, v_2, v_3) : \text{int}, \{wr_\rho, rd_\rho\} \cup \varepsilon_1 \cup \varepsilon_2 \cup \varepsilon_3} \\
\frac{\Pi | \Lambda | \Gamma \vdash e : \tau_1, \varepsilon_1 \quad \Pi, \Lambda \vdash \tau_1 \leq \tau_2 \quad \varepsilon_1 \subseteq \varepsilon_2}{\Pi | \Lambda | \Gamma \vdash e : \tau_2, \varepsilon_2} \quad (\text{FRV}(\varepsilon_2) \subseteq \Pi, \Lambda) \\
\frac{}{\Theta \vdash \tau \leq \tau} \quad (\text{FRV}(\tau) \subseteq \Theta) \quad \frac{\Theta \vdash \tau_1 \leq \tau'_1 \quad \Theta \vdash \tau_2 \leq \tau'_2}{\Theta \vdash \tau_1 \times \tau_2 \leq \tau'_1 \times \tau'_2} \\
\frac{\Theta \vdash \tau'_1 \leq \tau_1 \quad \Theta \vdash \tau_2 \leq \tau'_2 \quad \varepsilon_1 \subseteq \varepsilon_2 \quad \Pi_1 \subseteq \Pi_2 \quad \Lambda_1 \subseteq \Lambda_2}{\Theta \vdash \tau_1 \xrightarrow{\Pi_1, \Lambda_1} \tau_2 \leq \tau'_1 \xrightarrow{\Pi_2, \Lambda_2} \tau'_2} \quad (\text{FRV}(\varepsilon_2) \cup \Pi_2 \cup \Lambda_2 \subseteq \Theta)
\end{array}$$

Fig. 3. Typing and subtyping relations. Notice that for a typing judgement $\Pi | \Lambda | \Gamma \vdash e : \tau, \varepsilon$ we always have $\text{FRV}(\Gamma, \tau, \varepsilon) \subseteq \Pi \cup \Gamma$.

III. DEFINITION OF THE LOGICAL RELATION

Semantic Types and Worlds We give a Kripke or world-indexed logical relation. This is a fairly standard approach to modeling dynamic allocation; in combination with higher-order store, however, it comes with a fairly standard problem: the *type-world circularity*. Roughly, semantic types are indexed over worlds and worlds contain semantic types, so both need to be defined before the other. A specific instance of this circularity was solved recently by Thamsborg and Birkedal [20] based on metric-space theory developed by Birkedal et. al. [27]; we re-use that solution here. Semantic types (and worlds) are constructed as a fixed-point of an endofunctor on a certain category of metric-spaces. We do not care about that, though; we just give the result of the construction. In addition, we largely ignore the fact that we actually deal in metric spaces and not just plain sets; the little metric machinery we need is deferred to Appendix A.

There is a set \mathbf{T} of *semantic types* and a set \mathbf{W} of *worlds*; types are world-indexed relations on values and worlds describe the regions and type-layouts of heaps, roughly speaking. Take a type $\mu \in \mathbf{T}$ and apply it to a world $w \in \mathbf{W}$ and you get an indexed relation on values, i.e.,

$$\mu(w) \subseteq \mathbb{N} \times \mathcal{V} \times \mathcal{V}.$$

These relations are downwards closed in the first coordinate; we read $(k, v_1, v_2) \in \mu(w)$ as saying that v_1 and v_2 are related at type μ up to approximation k assuming world w .

We assume a countably infinite set of region names \mathcal{RN} ; a world $w \in \mathbf{W}$ contains finitely many such $|w| \subseteq_{\text{fin}} \mathcal{RN}$. Some of these $\text{dom}(w) \subseteq |w|$ are *live* and the rest are *dead*. To each live region $r \in \text{dom}(w)$ we associate a finite partial bijection $w(r)$ on locations decorated with types, i.e.,

$$w(r) \subseteq_{\text{fin}} \mathcal{L} \times \mathcal{L} \times \hat{\mathbf{T}}$$

such that for $(l_1, l_2, \mu), (m_1, m_2, \nu) \in w(r)$ we have that both $l_1 = m_1$ and $l_2 = m_2$ imply $l_1 = m_1, l_2 = m_2$ and $\mu = \nu$. We write $\text{dom}_1(w(r))$ for the set of left hand side locations in the bijection and $\text{dom}_2(w(r))$ for the right hand side ones; different regions must have disjoint left and right hand side locations. For convenience, we set

$$\text{dom}_1^A(w) = \bigcup_{r \in A \cap \text{dom}(w)} \text{dom}_1(w(r))$$

whenever $A \subseteq |w|$, and we write $\text{dom}_1(w)$ for $\text{dom}_1^{|w|}(w)$, i.e., the set of all left hand side locations. Similarly for the right hand side.

Worlds evolve and types adapt. Triples of two locations and a type can be added to a live region, as long as different regions remain disjoint. Orthogonal to this, one can add a fresh, i.e., neither live nor dead, region name with an associated empty partial bijection. And one can kill any live region, rendering it dead and losing the associated the partial bijection in the process. The reflexive, transitive closure of all three combined is a preorder \sqsubseteq on worlds; it is a crucial property of types that they respect this, i.e., that

$$w \sqsubseteq w' \implies \mu(w) \subseteq \mu(w')$$

for any two $w, w' \in \mathbf{W}$ and any $\mu \in \mathbf{T}$. This is *type monotonicity* and it prevents values from fleeing types over time.

Finally, to tie the knot, there is an isomorphism $\iota : \hat{\mathbf{T}} \rightarrow \mathbf{T}$ from the odd types stored in worlds to proper types. Whenever a type is extracted from a world it needs to be coerced by this isomorphism before it can be applied to some world.

The Logical Relation and Interpretation of Types Often, a logical relation goes like this: two computations are related

if they (from related heaps) reduce to related values (and heaps); this is the *extensional* view: we do not care about the intermediate states. As we consider concurrency, however, a computation can be interrupted and so we need to start caring. In our setup, public regions are accessible from the environment. To address this, we assume that before each reduction step, the public regions hold related values; in return, we promise related values after the step. In other words, the *granularity of extensionality* is just one step for the public regions. For private regions, however, there is no interference and the granularity is an entire computation as usual. This is the fundamental idea; it is how we propose to stay extensional in the face of currency.

Without further ado, let us look into the cornerstone of our model: the safety relation defined in Figure 7; auxiliary relations are defined and named in Figure 8. What does it mean to have

$$(k, h_1^\circ, h_2^\circ, e_1, e_2, h_1, h_2) \in \mathbf{safe}_{\tau, \varepsilon}^{\Pi, \Lambda, A, R} w^\circ, w?$$

Overall, it says that after environment interference, we can match the behavior of e_1 , i.e., termination or any one-step reduction, by zero or more steps of e_2 ; match in the sense of (re-)establishing certain relations, including safety itself. Safety is a *local* property of a pair of computations, this is crucial: it has no knowledge of computations running concurrently and h_1 and h_2 are the *local heaps*, i.e., the parts of the global heaps that the e_1 respectively e_2 control exclusively. The computations consider $R(\Pi)$ to be their public, $R(\Lambda)$ to be their private and A to be their *anonymous* regions. The latter intuitively are private regions that have been masked out: they exist only for the duration of these computations, but we have to track them to deny the environment access; this is another difficulty imposed by concurrency. Safety is indexed by a world w as well; note that worlds are *global* things: all concurrent threads share one world, i.e., they agree about the division of the heap into regions and the types associated to locations. Finally k is intuitively the number of steps we are safe for, h_1° and h_2° are the (private parts of) the initial local heaps, τ is the expected return type, ε the effects and w° the initial world.

We unroll the definition in writing. The first pair of big square bracket — the *prerequisites* — translates to ‘the environment interferes’. This yields a new world w' subject to the constraints of the environment transition relation: no public, private or anonymous regions are killed, and the latter two see no allocation either. These are the expectations that we have of the environment. The actual contents of the public regions are unknown, but we are free to assume that they hold related values of the proper type, at least where we have read effects; this is the *public heaps* g_1 and g_2 in the precondition relation. In addition we have *frames* f_1 and f_2 that cover the remainder of the world and a triple-split relation that ensures coherence between the domains of corresponding parts of the world and the heaps, see Figure 4.

The left hand side is irreducible in the *termination branch* and takes one step in the *progress branch*. In either case, we

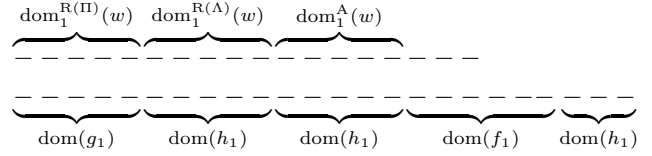


Fig. 4. The left hand side of the triple-split relation. The top dashed line is $\text{dom}_1(w)$, the bottom dashed line $\text{dom}(g_1 \cdot h_1 \cdot f_1)$. The local heap h_1 has a *private* part matching the private regions, an *anonymous* part matching the anonymous regions and an *off-world* part outside the domain of the world. The frame f_1 must cover regions that are neither public, private nor anonymous.

must match this in zero or more steps on the right hand side, not touching the frame; this means finding a future world w'' and relating a number of things. The choice of future world is restricted by the self transition relation: we must not kill private or public regions, but we can allocate in them, and regions that we know nothing about must be left untouched; this is our promise to the environment. In the termination branch, we are furthermore required to kill off all anonymous regions as the computation is done; any new regions added in the progress branch go to the set of anonymous regions. In both branches, the changes made to the public heap must be well-typed and permitted by the effects and, if we are done, we check the changes made to the (private part of) the local heaps as well; the fact that the public heaps are compared across a single stage and the (private parts of) the local heaps are compared across the entire computations is the crux of the idea of having different granularities of extensionality.

In addition to performing actual allocation, we have the possibility of moving existing locations from, say, the off-world part of the local heap into the public heap or the private part of the public heap; this is a subtle point that permits the actual allocation of new locations and the corresponding extension of the world to be temporarily out of sync.

We have glossed over one aspect of safety: the right hand side takes steps in the ordinary operational semantics, but the left hand side works in the *instrumented operational semantics* of Figure 5. A reduction $(e | h) \rightarrow_\mu^n (e' | h')$ in the latter implies a similar reduction in former; in addition it counts the steps of a reduction with all atomic commands ‘unfolded’ (with unfolding itself counting one step) and it records all heap accesses. We need the former for compatibility of the atomic typing rule in Lemma 6 below: atomic commands really unfold as they execute, hence we must count the number of ‘unfolded’ steps. It is less immediate that we must test the actual reads, writes and allocations, recorded by μ , against the effects described by ε , as done in the progress branch of safety; it is an intensional weed in our garden of extensional flowers. But if omitted, our present proof of the Parallelization Theorem falls short; the issue is technical, but it boils down to the fact that our present proof relies on the following simple, but crucial commutation property:

Lemma 2. *If we have $rd_l \notin \mu$, $wr_l \notin \mu$ and $al_l \notin \mu$ and $(e | h) \rightarrow_\mu^n (e' | h')$, then $(e | h[l \mapsto v]) \rightarrow_\mu^n (e' | h'[l \mapsto v])$.*

The actual logical relation is given in Figure 9. The

$$\begin{aligned}
& (E[\text{proj}_i \langle v_1, v_2 \rangle] | h) \rightarrow_{\emptyset}^1 (E[v_i] | h) \\
& (E[(\text{fun } f(x).e) v] | h) \rightarrow_{\emptyset}^1 (E[e[f := \text{fun } f(x).e, x := v]] | h) \\
& \quad (E[\text{ref } v] | h) \rightarrow_{\{al_i\}}^1 (E[l | h[l \mapsto v]] \text{ if } l \notin \text{dom}(h)) \\
& \quad (E[l := v] | h) \rightarrow_{\{wr_l\}}^1 (E[\langle \rangle] | h[l := v]) \text{ if } l \in \text{dom}(h) \\
& \quad (E[! l] | h) \rightarrow_{\{rd_l\}}^1 (E[h(l)] | h) \text{ if } l \in \text{dom}(h) \\
& (E[\text{par } v_1 \text{ and } v_2] | h) \rightarrow_{\emptyset}^1 (E[\langle v_1, v_2 \rangle] | h) \\
& (E[\text{cas } (l, n_1, n_2)] | h) \rightarrow_{\{rd_l, wr_l\}}^1 (E[1] | h[l := n_2]) \\
& \quad \text{if } l \in \text{dom}(h) \text{ and } h(l) = n_1 \\
& (E[\text{cas } (l, n_1, n_2)] | h) \rightarrow_{\{rd_l\}}^1 (E[0] | h) \\
& \quad \text{if } l \in \text{dom}(h) \text{ and } h(l) \neq n_1 \\
& (E[\text{atomic } e] | h) \rightarrow_{\mu}^{n+1} (E[v] | h') \text{ if } (e | h) \xrightarrow{\mu}^n (v | h') \\
& (E[\text{atomic } e] | h) \rightarrow_{\emptyset}^1 (E[\text{atomic } e] | h) \\
& (e | h) \xrightarrow{\emptyset}^0 (e' | h') \iff e = e' \wedge h = h' \\
& (e | h) \xrightarrow{\mu}^n (e' | h') \iff (e | h) \rightarrow_{\mu_1}^{n_1} (e' | h') \wedge \\
& \quad (e' | h') \xrightarrow{\mu_2}^{n_2} (e'' | h'') \wedge \\
& \quad n = n_1 + n_2 \wedge \mu = \mu_1 \cup \mu_2.
\end{aligned}$$

Fig. 5. Instrumented operational semantics. $(e | h) \rightarrow_{\mu}^n (e' | h')$ is defined for $e, e' \in \mathcal{E}$ with $\text{FV}(e, e') = \emptyset$, $h, h' \in \mathcal{H}$, $n \geq 1$ and $\mu \subseteq \{al_i \mid l \in \mathcal{L}\} \cup \{wr_l \mid l \in \mathcal{L}\} \cup \{rd_l \mid l \in \mathcal{L}\}$; the starred version additionally permits a zero as superscript. Note that $(e | h) \mapsto (e' | h')$ if and only if there are n and μ such that $(e | h) \rightarrow_{\mu}^n (e' | h')$.

existentially quantified $a \in \mathbb{N}$ is the minimal number of anonymous regions required to run; otherwise it builds on safety in a straightforward way. There is some asymmetry to these definitions: the anonymous regions A are required to exist (and be empty) in the world beforehand, but are killed off in the termination branch; also the precondition on the (private parts of) the initial local heaps is in the logical relation whereas the postcondition lives in the termination branch. The interpretation of types is in Figure 6. Interpreting the function type looks daunting, but a function is just a suspended expression with a single free variable, hence we have to stuff most of the logical relation into the definition. Apart from that, we just remark that the $R(\rho) \not\subseteq \text{dom}(w)$ case of reference interpretation is part of an approach to handling dangling pointers (due to region masking) proposed recently in [20]; similarly for the $R(\text{FRV}(\varepsilon)) \not\subseteq \text{dom}(w)$.

As conclusion to this subsection we give a theorem that, combined with the upcoming compatibility, means that logical relatedness implies contextual may-approximation. The proof is in Appendix B and it is not hard, but it is worth noting that we need a proof at all: with sequential languages, this is a result that one reads off the definition of the logical relation.

Theorem 3 (May-Equivalence). *Assume that $\cdot | \cdot | \cdot \models e_1 \preceq e_2 : \text{int}, \emptyset$ holds. Take any $h_1, h_2 \in \mathcal{H}$. If there are e'_1, h'_1 with $(e_1 | h_1) \xrightarrow{*} (e'_1 | h'_1)$ such that $\text{irr}(e'_1 | h'_1)$ holds, then there is $n \in \mathbb{Z}$ such that $e'_1 = \underline{n}$ and h'_2 such that $(e_2 | h_2) \xrightarrow{*} (\underline{n}, h'_2)$.*

$$\begin{aligned}
& \mathbf{envtran}^{\Pi, \Lambda, A, R} w, w' \iff \\
& \quad w \sqsubseteq w' \wedge \forall r \in \text{dom}(w) \cap (R(\Pi \cup \Lambda) \cup A). r \in \text{dom}(w') \\
& \quad \wedge \forall r \in \text{dom}(w) \cap (R(\Lambda) \cup A). w(r) = w'(r).
\end{aligned}$$

$$\begin{aligned}
& \mathbf{selftran}^{\Pi, \Lambda, A, R} w, w' \iff \\
& \quad w \sqsubseteq w' \wedge \forall r \in \text{dom}(w) \setminus A. r \in \text{dom}(w') \\
& \quad \wedge \forall r \in \text{dom}(w) \setminus (R(\Pi \cup \Lambda) \cup A). w(r) = w'(r).
\end{aligned}$$

$$\begin{aligned}
& (g_1, h_1, f_1, g_2, h_2, f_2) \in \mathbf{splits}^{\Pi, \Lambda, A, R} w \iff \\
& \quad \text{dom}(h_1) \# \text{dom}(g_1) \# \text{dom}(f_1) \wedge \\
& \quad \text{dom}(h_2) \# \text{dom}(g_2) \# \text{dom}(f_2) \wedge \\
& \quad \text{dom}_1^{R(\Pi)}(w) = \text{dom}(g_1) \wedge \text{dom}_1^{R(\Lambda) \cup A}(w) \subseteq \text{dom}(h_1) \wedge \\
& \quad \text{dom}_1^{\text{dom}(w) \setminus (R(\Pi \cup \Lambda) \cup A)}(w) \subseteq \text{dom}(f_1) \wedge \\
& \quad \text{dom}_2^{R(\Pi)}(w) = \text{dom}(g_2) \wedge \text{dom}_2^{R(\Lambda) \cup A}(w) \subseteq \text{dom}(h_2) \wedge \\
& \quad \text{dom}_2^{\text{dom}(w) \setminus (R(\Pi \cup \Lambda) \cup A)}(w) \subseteq \text{dom}(f_2).
\end{aligned}$$

$$\begin{aligned}
& (k, h_1, h_2) \in \mathbf{P}_{\varepsilon}^{\Theta, R} w \iff \\
& \quad \text{dom}(h_1) = \text{dom}_1^{R(\Theta)}(w) \wedge \text{dom}(h_2) = \text{dom}_2^{R(\Theta)}(w) \wedge \\
& \quad \forall r \in R(\Theta) \cap \text{dom}(w). \forall (l_1, l_2, \mu) \in w(r). \\
& \quad r \in R(\text{rds } \varepsilon) \Rightarrow k > 0 \Rightarrow (k - 1, h_1(l_1), h_2(l_2)) \in (\iota \mu)(w).
\end{aligned}$$

$$\begin{aligned}
& (k, h_1, h_2, h'_1, h'_2) \in \mathbf{Q}_{\varepsilon}^{\Theta, R} w, w' \iff \\
& \quad \text{dom}(h_1) = \text{dom}_1^{R(\Theta)}(w) \wedge \text{dom}(h_2) = \text{dom}_2^{R(\Theta)}(w) \wedge \\
& \quad \text{dom}(h'_1) = \text{dom}_1^{R(\Theta)}(w') \wedge \text{dom}(h'_2) = \text{dom}_2^{R(\Theta)}(w') \wedge \\
& \quad (\forall r \in R(\Theta) \cap \text{dom}(w). \forall (l_1, l_2, \mu) \in w(r). \\
& \quad [h_1(l_1) = h'_1(l_1) \wedge h_2(l_2) = h'_2(l_2)] \vee [r \in R(\text{wrs } \varepsilon) \wedge \\
& \quad k > 0 \Rightarrow (k - 1, h'_1(l_1), h'_2(l_2)) \in (\iota \mu)(w')]) \wedge \\
& \quad (\forall r \in R(\Theta) \cap \text{dom}(w). \\
& \quad \forall (l_1, l_2, \mu) \in w'(r) \setminus w(r). r \in R(\text{als } \varepsilon) \wedge \\
& \quad k > 0 \Rightarrow (k - 1, h'_1(l_1), h'_2(l_2)) \in (\iota \mu)(w')).
\end{aligned}$$

$$\begin{aligned}
& \mu \in \mathbf{effs}_{\varepsilon, h}^{A, R} w \iff \\
& \quad \{l \mid rd_l \in \mu\} \cap \text{dom}_1(w) \subseteq \text{dom}_1^{R(\text{rds } \varepsilon) \cup A}(w) \wedge \\
& \quad \{l \mid wr_l \in \mu\} \cap \text{dom}_1(w) \subseteq \text{dom}_1^{R(\text{wrs } \varepsilon) \cup A}(w) \wedge \\
& \quad \{l \mid al_i \in \mu\} \cap \text{dom}_1(w) \subseteq \text{dom}_1^{R(\text{als } \varepsilon) \cup A}(w) \wedge \\
& \quad \{l \mid rd_l \in \mu \vee wr_l \in \mu \vee al_i \in \mu\} \setminus \text{dom}_1(w) \subseteq \text{dom}(h).
\end{aligned}$$

Fig. 8. Six auxiliary definitions. The *environment transition* and *self transition* relations are defined for $\Pi \# \Lambda, R : \Pi \cup \Lambda \leftrightarrow |w|, A \subseteq \text{dom}(w)$ and $R(\Pi \cup \Lambda) \# A$. The *triple-split* relation has the same prerequisites. The *precondition* relation is defined for $R : \mathcal{RV} \xrightarrow{\text{fin}} |w|$ injective with $\Theta \cup \text{FRV}(\varepsilon) \subseteq \text{dom}(R)$. The *postcondition* relation additionally requires $w' \sqsupseteq w$ such that $\text{dom}(w) \cap R(\Theta) \subseteq \text{dom}(w')$. Finally the *actual-effects* relation expects $R : \mathcal{RV} \xrightarrow{\text{fin}} |w|$ injective with $\text{FRV}(\varepsilon) \subseteq \text{dom}(R)$ and $A \subseteq \text{dom}(w)$.

$$\begin{aligned}
\llbracket \mathbf{1} \rrbracket^R w &= \{(k, (), ()) \mid k \in \mathbb{N}\} & \llbracket \mathbf{int} \rrbracket^R w &= \{(k, n, n) \mid k \in \mathbb{N} \wedge n \in \mathbb{Z}\} \\
\llbracket \tau_1 \times \tau_2 \rrbracket^R w &= \{(k, (v_{11}, v_{21}), (v_{12}, v_{22})) \mid (k, v_{11}, v_{12}) \in \llbracket \tau_1 \rrbracket^R w \wedge (k, v_{21}, v_{22}) \in \llbracket \tau_2 \rrbracket^R w\} \\
\llbracket \mathbf{ref}_{\rho} \tau \rrbracket^R w &= \begin{cases} \{(k, l_1, l_2) \mid \exists \mu \in \widehat{\mathbf{T}}. (l_1, l_2, \mu) \in w(R(\rho)) \wedge \forall w' \sqsupseteq w. \llbracket \tau \rrbracket^R w' \stackrel{k}{=} (\iota \mu)(w')\} & R(\rho) \in \text{dom}(w) \\ \{(k, v_1, v_2) \mid k \in \mathbb{N} \wedge v_1, v_2 \in \mathcal{V}\} & R(\rho) \notin \text{dom}(w) \end{cases} \\
\llbracket \tau_1 \xrightarrow{\Pi, \Lambda} \tau_2 \rrbracket^R w &= \left\{ \begin{array}{l} (k, \text{fun } f(x).e_1, \text{fun } f(x).e_2) \mid \exists a \in \mathbb{N}. \forall j < k. \forall w' \sqsupseteq w. \\ \forall A \subseteq \text{dom}(w'). \forall v_1, v_2 \in \mathcal{V}. \forall h_1, h_2, h'_1, h'_2 \in \mathcal{H}. \\ \left[\begin{array}{l} R(\text{FRV}(\varepsilon)) \subseteq \text{dom}(w') \wedge A \# R(\Pi \cup \Lambda) \wedge |A| \geq a \wedge \forall r \in A. w'(r) = \emptyset \wedge \\ (j, v_1, v_2) \in \llbracket \tau_1 \rrbracket^R w' \wedge h'_1 \subseteq h_1 \wedge h'_2 \subseteq h_2 \wedge (j, h'_1, h'_2) \in \mathbf{P}_{\varepsilon}^{\Lambda, R} w' \\ (j, h'_1, h'_2, e_1[\text{fun } f(x).e_1/f, v_1/x], e_2[\text{fun } f(x).e_2/f, v_2/x], h_1, h_2) \in \mathbf{safe}_{\tau_2, \varepsilon}^{\Pi, \Lambda, A, R} w', w' \end{array} \right] \Rightarrow \\ \{(k, v_1, v_2) \mid k \in \mathbb{N} \wedge v_1, v_2 \in \mathcal{V}\} \end{array} \right\} \begin{array}{l} R(\text{FRV}(\varepsilon)) \subseteq \text{dom}(w) \\ R(\text{FRV}(\varepsilon)) \not\subseteq \text{dom}(w) \end{array}
\end{aligned}$$

Fig. 6. Interpretation of types. We require $R : \mathcal{RV} \rightarrow_{fin} \mathcal{RN}$ injective with $\text{FRV}(\tau) \subseteq \text{dom}(R)$. We assume $R(\text{FRV}(\tau)) \subseteq |w|$ above, otherwise we define $\llbracket \tau \rrbracket^R w$ to be the empty set. We get that $\llbracket \tau \rrbracket^R \in \mathbf{T}$.

$$\begin{aligned}
(k, h_1^\circ, h_2^\circ, e_1, e_2, h_1, h_2) &\in \mathbf{safe}_{\tau, \varepsilon}^{\Pi, \Lambda, A, R} w^\circ, w \\
&\iff \\
\forall j \leq k. \forall w', g_1, g_2, f_1, f_2. & \\
\left[\mathbf{envtran}^{\Pi, \Lambda, A, R} w, w' \wedge (j, g_1, g_2) \in \mathbf{P}_{\varepsilon}^{\Pi, R} w' \wedge (g_1, h_1, f_1, g_2, h_2, f_2) \in \mathbf{splits}^{\Pi, \Lambda, A, R} w' \right] \Rightarrow & \\
\left[\text{irr}(e_1 | g_1 \cdot h_1 \cdot f_1) \Rightarrow \right. & \\
\exists e'_2, w'', h'_1, h'_2, g'_1, g'_2. & \\
(e_2 | g_2 \cdot h_2 \cdot f_2) \mapsto^* (e'_2 | g'_2 \cdot h'_2 \cdot f_2) \wedge \mathbf{selftran}^{\Pi, \Lambda, A, R} w', w'' \wedge \emptyset = (A \cap \text{dom}(w'')) \cup (\text{dom}(w'') \setminus \text{dom}(w')) \wedge & \\
g_1 \cdot h_1 = g'_1 \cdot h'_1 \wedge (g'_1, h'_1, f_1, g'_2, h'_2, f_2) \in \mathbf{splits}^{\Pi, \Lambda, \emptyset, R} w'' \wedge (j, e_1, e_2) \in \llbracket \tau \rrbracket^R (w'') \wedge & \\
(j, g_1, g_2, g'_1, g'_2) \in \mathbf{Q}_{\varepsilon}^{\Pi, R} w', w'' \wedge \exists h''_1 \subseteq h'_1, h''_2 \subseteq h'_2. (j, h_1^\circ, h_2^\circ, h''_1, h''_2) \in \mathbf{Q}_{\varepsilon}^{\Lambda, R} w^\circ, w'' \left. \right] \wedge & \\
\left[\forall e'_1, h_1^\dagger, \mu, n \leq j. (e_1 | g_1 \cdot h_1 \cdot f_1) \rightarrow_{\mu}^n (e'_1 | h_1^\dagger) \Rightarrow \right. & \\
\exists e'_2, w'', A', h'_1, h'_2, g'_1, g'_2. & \\
(e_2 | g_2 \cdot h_2 \cdot f_2) \mapsto^* (e'_2 | g'_2 \cdot h'_2 \cdot f_2) \wedge \mathbf{selftran}^{\Pi, \Lambda, A, R} w', w'' \wedge A' = (A \cap \text{dom}(w'')) \cup (\text{dom}(w'') \setminus \text{dom}(w')) \wedge & \\
h_1^\dagger = g'_1 \cdot h'_1 \cdot f_1 \wedge \mu \in \mathbf{effs}_{\varepsilon, h_1^\dagger}^{A', R} w'' \wedge (g'_1, h'_1, f_1, g'_2, h'_2, f_2) \in \mathbf{splits}^{\Pi, \Lambda, A', R} w'' \wedge & \\
(j - n, g_1, g_2, g'_1, g'_2) \in \mathbf{Q}_{\varepsilon}^{\Pi, R} w', w'' \wedge (j - n, h_1^\circ, h_2^\circ, e'_1, e'_2, h'_1, h'_2) \in \mathbf{safe}_{\tau, \varepsilon}^{\Pi, \Lambda, A', R} w^\circ, w'' \left. \right]
\end{aligned}$$

Fig. 7. Safety. The predicate is defined by well-founded induction. Nontrivial requirements are: $\Pi \# \Lambda$, $\text{FRV}(\tau, \varepsilon) \subseteq \Pi \cup \Lambda$, $\text{FV}(e_1, e_2) = \emptyset$, $R : \Pi \cup \Lambda \hookrightarrow |w^\circ|$, $R(\text{FRV}(\varepsilon)) \subseteq \text{dom}(w^\circ)$ and $w \sqsupseteq w^\circ$ with $\text{dom}(w^\circ) \cap R(\Pi \cup \Lambda) \subseteq \text{dom}(w)$, $A \subseteq \text{dom}(w)$ and $A \# R(\Pi \cup \Lambda)$. See Figure 8 for auxiliary definitions. We refer to the contents of the big square brackets as the *prerequisites*, the *termination branch* and the *progress branch*, respectively.

$$\begin{aligned}
\Pi \mid \Lambda \mid \Gamma \models e_1 \preceq e_2 : \tau, \varepsilon & \\
&\iff \\
\exists a \in \mathbb{N}. \forall k \in \mathbb{N}. \forall w \in \mathbf{W}. \forall R : \Pi \cup \Lambda \hookrightarrow |w|. \forall A \subseteq \text{dom}(w). \forall \gamma_1, \gamma_2 \in \mathcal{V}^{|\Gamma|}. \forall h_1, h_2, h'_1, h'_2 \in \mathcal{H}. & \\
\left[R(\text{FRV}(\varepsilon)) \subseteq \text{dom}(w) \wedge A \# R(\Pi \cup \Lambda) \wedge |A| \geq a \wedge \forall r \in A. w(r) = \emptyset \wedge \right. & \\
(k, \gamma_1, \gamma_2) \in \llbracket \Gamma \rrbracket^R w \wedge h'_1 \subseteq h_1 \wedge h'_2 \subseteq h_2 \wedge (k, h'_1, h'_2) \in \mathbf{P}_{\varepsilon}^{\Lambda, R} w \left. \right] \Rightarrow & \\
(k, h'_1, h'_2, e_1[\gamma_1/\Gamma], e_2[\gamma_2/\Gamma], h_1, h_2) \in \mathbf{safe}_{\tau, \varepsilon}^{\Pi, \Lambda, A, R} w, w. &
\end{aligned}$$

Fig. 9. The logical relation with anonymous regions. We require that $\Pi \# \Lambda$, that $\text{FRV}(\Gamma, \tau, \varepsilon) \subseteq \Pi \cup \Lambda$ and, as always, that $\text{FV}(e_1, e_2) \in |\Gamma|$.

Compatibility of the Logical Relation The logical relation is compatible, i.e., respects all typing rules. This is a *sine quo non* of logical relations; it implies the fundamental lemma stating that every well-typed expression is related to itself. And, as discussed just above, it makes the logical relation approximate contextual may-approximation:

Theorem 4. $\Pi | \Lambda | \Gamma \models e_1 \preceq e_2 : \tau, \varepsilon$ implies $\Pi | \Lambda | \Gamma \vdash e \lesssim_{\downarrow} e' : \tau, \varepsilon$.

Compatibility means that each typing rule induces a lemma by reading the (unary) typing judgments as the corresponding (binary) logical relations. The three most interesting of these have to do with concurrency and the divide between public and private regions; they are listed here and proofs are given in Appendix D:

Lemma 5. $\Pi | \Lambda, \rho | \Gamma \models e_1 \preceq e_2 : \tau, \varepsilon$ implies $\Pi | \Lambda | \Gamma \models e_1 \preceq e_2 : \tau, \varepsilon - \rho$ provided that $\rho \notin \text{FRV}(\Gamma, \tau)$.

Lemma 6. $\cdot | \Pi, \Lambda | \Gamma \models e_1 \preceq e_2 : \tau, \varepsilon$ implies $\Pi | \Lambda | \Gamma \models$ atomic $e_1 \preceq$ atomic $e_2 : \tau, \varepsilon$ if $\text{als} \varepsilon \subseteq \text{rds} \varepsilon \cap \text{wrs} \varepsilon$.

Lemma 7. $\Pi, \Lambda | \cdot | \Gamma \models e_1 \preceq e_2 : \tau, \varepsilon$ and $\Pi, \Lambda | \cdot | \Gamma \models e_1^{\dagger} \preceq e_2^{\dagger} : \tau^{\dagger}, \varepsilon^{\dagger}$ together imply $\Pi | \Lambda | \Gamma \models \text{par } e_1$ and $e_1^{\dagger} \preceq \text{par } e_2$ and $e_2^{\dagger} : \tau \times \tau^{\dagger}, \varepsilon \cup \varepsilon^{\dagger}$.

IV. APPLICATIONS

A. Parallelization Theorem: Disjoint Concurrency

We now explain our Parallelization Theorem, which gives us an easy way to prove properties about the common case of disjoint concurrency, where disjointness is captured using private regions and effect annotations.

Theorem 8 (Parallelization). Assuming that

- 1) $\Pi, \Lambda | \cdot | \Gamma \vdash e_1 : \tau_1, \varepsilon_1$,
- 2) $\Pi, \Lambda | \cdot | \Gamma \vdash e_2 : \tau_2, \varepsilon_2$,
- 3) $\text{rds } \varepsilon_1 \cup \text{wrs } \varepsilon_1 \cup \text{rds } \varepsilon_2 \cup \text{wrs } \varepsilon_2 \subseteq \Lambda$,
- 4) $\text{rds } \varepsilon_1 \cap \text{wrs } \varepsilon_2 = \text{rds } \varepsilon_2 \cap (\text{wrs } \varepsilon_1 \cup \text{als } \varepsilon_1) = \text{wrs } \varepsilon_1 \cap \text{wrs } \varepsilon_2 = \emptyset$,

the following property holds:

$$\Pi | \Lambda | \Gamma \models \langle e_1, e_2 \rangle \cong \text{par } e_1 \text{ and } e_2 : \tau_1 \times \tau_2, \varepsilon_1 \cup \varepsilon_2.$$

Intuitively, item 3 keeps the environment from detecting anything, and item 4 prevents the two computations from talking among themselves, thereby making them independent; the $\text{als } \varepsilon_1$ in item 4 is a technicality that we cannot do without. We showed a concrete simple application of this theorem in the Introduction. More generally, example usage includes situations where we operate on two imperative data structures (say linked lists or graphs); if we only mutate parts of the data structures that are in different regions, then we may safely parallelize operations on the data structures.

The masking rule makes it possible to do more optimizations via the Parallelization Theorem: Consider, for simplicity, the familiar example of an efficient implementation *fib* of the Fibonacci function using two local references. We can use the masking rule to give it type and effect $\text{int} \rightarrow_{\emptyset} \text{int}, \emptyset$. This

allows us to view the imperative implementation as pure, and thus by Theorem 8 we find that it is sound to optimize two sequential calls to *fib* to two parallel calls. This may sound like a simple optimization, but the point is that a compiler can perform it automatically, just based on the effect types.

The proof of the Parallelization Theorem is quite tricky. We now sketch the idea of the proof; please see the appendix for the technical details.

The left-to-right direction of the theorem is not too hard, intuitively since sequential sequencing clearly can be mimicked by parallel reduction. Thus, the hard part of the proof is to show the direction

$$\Pi | \Lambda | \Gamma \models \text{par } e_1 \text{ and } e_2 \lesssim_{\downarrow} \langle e_1, e_2 \rangle : \tau_1 \times \tau_2, \varepsilon_1 \cup \varepsilon_2. \quad (1)$$

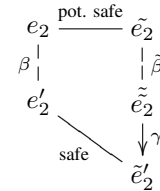
The challenge is that if we take a step of reduction on the left, then we cannot always immediately mimick it on the right, because if it is a step in (a derivative of) e_2 then we might not be ready to make it on the right, since we haven't yet started reducing in e_2 on the right (we only start reducing e_2 on the right when we are done with all the e_1 reductions).

We address this issue by formulating a notion of *potential safety* to relate (a derivative of) e_2 on the left with (a derivative of) e_2 on the right. Safety, and also potential safety, really involves configurations (expressions and heaps) rather than just expression), but to keep the overview, we will continue just using expression notation in the high-level sketch here. Potential safety of e_2 and \tilde{e}_2 then allows \tilde{e}_2 to *catch up* by doing some reductions to reach \tilde{e}'_2 , and then e_2 and \tilde{e}'_2 should be safe (in the usual sense).

The idea of the proof of (1) is then to show (for suitable $e_1, \tilde{e}_1, e_2, \tilde{e}_2$ satisfying conditions corresponding to those in the parallelization theorem) that

$$\text{safe}(e_1, \tilde{e}_1) \wedge \text{potentially-safe}(e_2, \tilde{e}_2) \implies \text{safe}(\langle \text{par } e_1 \text{ and } e_2 \rangle, \langle \tilde{e}_1, \tilde{e}_2 \rangle). \quad (2)$$

Now, for this to work, potential safety actually needs to allow for some transitions by e_1 and \tilde{e}_1 — and the context — but regulated by the overall assumptions regarding effects, etc., of the parallelization theorem. Thus potential safety of e_2 and \tilde{e}_2 actually says that for any good (in the sense that it is governed by overall assumptions regarding effects, etc.) pair of transitions, β and $\tilde{\beta}$, going from e_2 to e'_2 and \tilde{e}_2 to \tilde{e}'_2 respectively, there is some sequence of catching-up reductions, call them γ , that reaches a \tilde{e}'_2 such that e'_2 and \tilde{e}'_2 are safe. Diagrammatically:



To show (2), we proceed by induction on the index used in the safety predicate. There are three cases to consider, corresponding to possible reductions of $\text{par } e_1$ and e_2 . The case where e_1 reduces is not too hard since we have defined

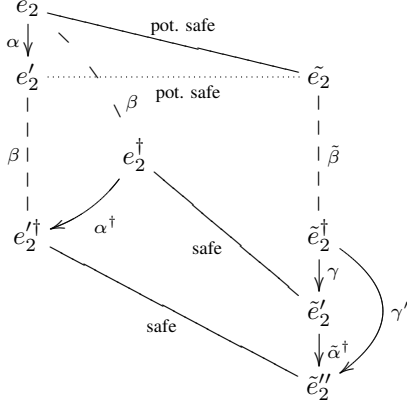


Fig. 10. Proving parallelization: recovering potential safety after α transition.

potential safety to accommodate such reductions (the β 's). The case where e_1 and e_2 are both values is not too hard either. The case where e_2 reduces, say to e_2' by a reduction α , is more tricky. We then show potential safety of e_2' and \tilde{e}_2 (argument outline follows below), and that e_1 and \tilde{e}_1 stay safe (this follows from the assumption on effects, etc.), which allows us to conclude by induction.

The argument for the potential safety of e_2' and \tilde{e}_2 is outlined in Figure 10. By assumption e_2 and \tilde{e}_2 are potentially safe, as depicted by the topmost arrow, and we need to show the potential safety depicted by the dashed arrow. Hence we consider any β and $\tilde{\beta}$ as shown, and we need to show that there exists a catching-up reduction γ' from \tilde{e}_2^\dagger to \tilde{e}_2'' . By properties of the instrumented operational semantics, we can commute α and β to get β and α^\dagger , as shown in the diagram (essentially, this is Lemma 2). Hence by the assumption of potential safety of e_2 and \tilde{e}_2 , we get safety of e_2^\dagger and \tilde{e}_2^\dagger , as shown. Using this safety and the α^\dagger reduction, we get a reduction $\tilde{\alpha}^\dagger$ and the final required safety of e_2^\dagger and \tilde{e}_2'' . Thus the required catching-up reduction γ' is γ followed by $\tilde{\alpha}^\dagger$. Note that the latter makes intuitive sense; we did a reduction on the left in e_2 and now we have correspondingly extended the catching-up reduction to be done on the right.

Formally, the Catch-up lemma in Appendix E gives the precise formulation and proof of (2); item 3 in that lemma gives the precise technical formulation of potential safety.

B. Non-disjoint Concurrency

We now exemplify how our logical relations model can also be used to reason compositionally about equivalences of fine-grained concurrent programs operating on public regions.

Consider the following type

$$\tau \equiv \text{ref}_\rho \text{int} \rightarrow_{\{\text{rd}_\rho, \text{wr}_\rho\}}^{\rho, \emptyset} 1$$

for functions that take an integer reference in a public region, possibly reads and writes from the reference, and returns unit.

The following two functions

$$\begin{aligned} \text{fun inc}_1(x). & \text{ let } y = !x \text{ in let } z = y + 1 \text{ in} \\ & \text{ if cas } (x, y, z) \text{ then } \langle \rangle \text{ else inc}_1(x) \\ \text{fun inc}_2(x). & \text{ atomic } (x := !x + 1) \end{aligned}$$

both have type τ . (We have allowed ourselves to use a standard conditional expression; 1 corresponds to true and 0 to false.) Both functions increment the integer given in their reference arguments; inc_1 uses the fine-grained compare-and-swap to do it atomically, whereas inc_2 uses the brute-force atomic operation. Using our logical relations model, we can prove that inc_1 and inc_2 are contextually equivalent:

$$\rho \mid \cdot \mid \cdot \vdash \text{inc}_1 \approx \text{inc}_2 : \tau, \emptyset. \quad (3)$$

Hence, replacing inc_2 with inc_1 in any well-typed client gives two contextually equivalent expressions. Thus our logical relation models a form of *data abstraction for concurrency* (where we abstract over the granularity of concurrency in the module).

We now show how to use the equivalence of inc_1 and inc_2 to derive equivalences of two different clients using the fine-grained concurrency implementation inc_1 .

To this end, consider the following two client programs of type

$$\begin{aligned} \sigma &\equiv \tau \rightarrow_{\emptyset}^{\rho, \emptyset} \text{ref}_\rho \text{int} \rightarrow_{\{\text{rd}_\rho, \text{wr}_\rho\}}^{\rho, \emptyset} \text{int} \\ \text{fun } c_1(\text{inc}). & \lambda n. \text{inc } n; \text{inc } n; !n \\ \text{fun } c_2(\text{inc}). & \lambda n. (\text{par } \text{inc } n \text{ and } \text{inc } n); !n \end{aligned}$$

Note that c_1 makes two sequential calls to inc , whereas c_2 runs the two calls in parallel. Because of the use of compare-and-swap in inc_1 , we would hope that the $c_1 \text{inc}_1$ and $c_2 \text{inc}_1$ are contextually equivalent (in typing context $\rho \mid \emptyset \mid \emptyset$). We can prove that this is indeed the case using compositional reasoning as follows. Using our logical relation, we prove that $c_1 \text{inc}_2$ is contextually equivalent to $c_2 \text{inc}_2$, i.e.,

$$\rho \mid \cdot \mid \cdot \vdash c_1 \text{inc}_2 \approx c_2 \text{inc}_2 : \text{ref}_\rho \text{int} \rightarrow_{\{\text{rd}_\rho, \text{wr}_\rho\}}^{\rho, \emptyset} \text{int}, \emptyset. \quad (4)$$

Finally, we conclude that $c_1 \text{inc}_1$ is contextually equivalent to $c_2 \text{inc}_1$ by transitivity of contextual equivalence:

$$\begin{aligned} c_1 \text{inc}_1 &\approx c_1 \text{inc}_2 && \text{by (3)} \\ &\approx c_2 \text{inc}_2 && \text{by (4)} \\ &\approx c_2 \text{inc}_1 && \text{by (3)}. \end{aligned}$$

Note that this proof illustrates the following important point. To show equivalence of two clients of a module implemented using fine-grained concurrency, it suffices to show that the clients are equivalent wrt. a coarse-grained implementation, and that the coarse-grained implementation is equivalent to the fine-grained implementation. The latter is often a lot simpler than trying to show the equivalence of the clients wrt. the fine-grained implementation directly. We can think of the coarse-grained implementation of the module (here inc_2) as the *specification* of the module and the fine-grained implementation (here inc_1) as the *implementation* of the module.

The formal proofs of (3) and (4) are mostly straightforward; for this direction

$$\rho \mid \cdot \mid \cdot \models \text{inc}_1 \preceq \text{inc}_2 : \tau, \emptyset$$

of (3), we use induction on the number of steps to prove the functions related, and only reduce the body of inc_2 once the compare-and-swap instruction in the body of inc_1 succeeds, while in the other case we use the induction hypothesis.

V. DISCUSSION

Gifford and Lucassen [11], [12] originally proposed type-and-effect systems as a static analysis for determining which parts of a higher-order imperative program could be implemented using parallelism. Here we are able to express the formal correctness of these ideas in a succinct way by having a parallel construct in our programming language and establishing the Parallelization Theorem.

In Section IV-B we showed how contextual equivalence can be used to *state* that compare-and-swap can be used to implement a simple form of locking, and how our logical relations model could be used to *prove* this statement. We believe that it should be possible to give similar succinct statements and proofs of other implementations of synchronization. For instance, we are currently working on a similar relational specification and correctness proof of Peterson’s mutual exclusion algorithm, which involves (for succinctness of specification) extending the language with a primitive notion of critical section.

As mentioned earlier, we have deliberately used the same definition of worlds here as in [20]. As discussed there [20, Section 8.2], this notion of world has somewhat limited expressiveness: the only heap invariants we can state are those that relate values at two locations by a semantic type. To increase expressiveness, it would thus be interesting to extend our model using ideas from [9], where worlds are defined using state-transition systems, and then investigate more examples of equivalences.

Recently, Liang et al. [10] have proposed RGSim, a simulation based on rely-guarantee, to verify program transformations in a concurrent setting. Their actual definition [10, Definition 4] bears some resemblance to our safety relation; indeed, an early draft of *loc.cit.* was a source of inspiration. They have no division of the heap into public and private parts, instead they give a pair of a rely and a guarantee that, respectively, constrains the interference of the environment and the actions of the computation in question. Their approach is essentially untyped; one point of view is that we ‘auto-instantiate’ the many parameters of their simulation based on our typing information. The languages they consider are first-order, languages with ground store; this obviously keeps life simple, but the example equivalences they give are not.

VI. CONCLUSION AND FUTURE WORK

We have presented a logical relations model of a new type-and-effect system for a concurrent higher-order ML-like language with general references. We have shown how to use the model for reasoning about both disjoint and non-disjoint concurrency. In particular, we have proved the first automatic Parallelization Theorem for such a rich language.

In this paper, we have focused on may contextual equivalence. Future work includes investigating models for must contextual equivalence. Since our language allows the encoding of countable nondeterminism, must equivalence is non-trivial, and will probably involve indexing over ω_1 rather than ω [28]. Future work also includes extending the model to

region and effect polymorphism, as well as the extension to more expressive worlds, and to other concurrency constructs such as fork-join.

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A. Well-definedness of Interpretation of Types

As explained in the main text, we import the worlds and types of Thamsborg and Birkedal [20] wholesale; hence their non-trivial construction is for free. The details of the construction, including metric prerequisites, can be found in Appendices A.1 and A.2 of the long version of *loc. cit.*; it is available online.¹ Thus it only remains to verify that interpreting a syntactic type actually gives a semantic ditto.

In metric terms, this means that for any syntactic type τ and any $R : \mathcal{RV} \rightarrow_{fin} \mathcal{RN}$ with $FRV(\tau) \subseteq \text{dom}(R)$ we must have

$$\llbracket \tau \rrbracket^R \in \mathbf{T} = \mathbf{W} \rightarrow_{mon} URel(\mathcal{V}),$$

i.e., it should map worlds to indexed, downwards closed relations on values in a non-expansive and monotone manner. To prove this, we need two lemmas on the pre- and postcondition relations:

Lemma 9. *We have $\mathbf{P}_\varepsilon^{\Theta,R} w \in URel(\mathcal{H})$. And for $w_1, w_2 \in \mathbf{W}$ with $w_1 \stackrel{n}{=} w_2$ we have $\mathbf{P}_\varepsilon^{\Theta,R} w_1 \stackrel{n}{=} \mathbf{P}_\varepsilon^{\Theta,R} w_2$ as well.*

Lemma 10. *We have $\mathbf{Q}_\varepsilon^{\Theta,R} w, w' \in URel(\mathcal{H} \times \mathcal{H})$. And for $w_1, w'_1, w_2, w'_2 \in \mathbf{W}$ with $w_1 \stackrel{n}{=} w_2$ and $w'_1 \stackrel{n}{=} w'_2$ we have $\mathbf{Q}_\varepsilon^{\Theta,R} w_1, w'_1 \stackrel{n}{=} \mathbf{Q}_\varepsilon^{\Theta,R} w_2, w'_2$ too.*

The proofs are quite straightforward. Then, by simultaneous induction, we can prove the following two propositions and we are done:

Proposition 11. *We have $\llbracket \tau \rrbracket^R \in \mathbf{T}$.*

Proposition 12. *For $w_1^\circ, w_1, w_2^\circ, w_2 \in \mathbf{W}$ we have that $w_1^\circ \stackrel{n}{=} w_2^\circ$ together with $w_1 \stackrel{n}{=} w_2$ implies $\mathbf{safe}_{\tau,\varepsilon}^{\Pi,\Lambda,A,R} w_1^\circ, w_1 \stackrel{n}{=} \mathbf{safe}_{\tau,\varepsilon}^{\Pi,\Lambda,A,R} w_2^\circ, w_2$.*

B. Proof of May-Equivalence Theorem

Theorem 3. *Assume that $\cdot | \cdot | \cdot \models e_1 \preceq e_2 : \text{int}, \emptyset$ holds. Take any $h_1, h_2 \in \mathcal{H}$. If there are e'_1, h'_1 with $(e_1 | h_1) \mapsto^* (e'_1 | h'_1)$ such that $\text{irr}(e'_1 | h'_1)$ holds, then there is $n \in \mathbb{Z}$ such that $e'_1 = \underline{n}$ and h'_2 such that $(e_2 | h_2) \mapsto^* (\underline{n}, h'_2)$.*

Proof. Pick, by assumption, $n \in \mathbb{N}$ such that $(e_1 | h_1) \mapsto^n (e'_1 | h'_1)$. We name the configurations in this multi-step reduction $(e_1^0 | h_1^0), (e_1^1 | h_1^1), \dots, (e_1^n | h_1^n)$ in order, i.e., such that

$$\begin{aligned} (e_1 | h_1) &= (e_1^0 | h_1^0) \mapsto (e_1^1 | h_1^1) \mapsto \dots \\ &\mapsto (e_1^n | h_1^n) = (e'_1 | h'_1). \end{aligned}$$

For each of these standard reductions, there is a corresponding instrumented reduction. This means, that we can pick $n_1, n_2, \dots, n_n \in \mathbb{N}$ non-zero and $\mu_1, \mu_2, \dots, \mu_n$ such that

$$(e_1^0 | h_1^0) \rightarrow_{\mu_1}^{n_1} (e_1^1 | h_1^1) \rightarrow_{\mu_2}^{n_2} \dots \rightarrow_{\mu_n}^{n_n} (e_1^n | h_1^n).$$

For convenience, we furthermore write $N_m = \sum_{i=m+1}^n n_i$ for each $0 \leq m \leq n$.

Now let $a \in \mathbb{N}$ be the required number of anonymous regions according to the definition of the logical relation,

and let $w \in \mathbf{W}$ be any world with a empty, live regions and nothing else. We now prove by induction that for all $0 \leq m \leq n$ the following holds:

$$\begin{aligned} \exists e'_2, h'_2, w'. (e_2 | h_2) \mapsto^* (e'_2 | h'_2) \wedge \\ \text{dom}(h_1^m) \supseteq \text{dom}_1(w') \wedge \text{dom}(h'_2) \supseteq \text{dom}_2(w') \wedge \\ (N_m, \llbracket \cdot \rrbracket, \llbracket \cdot \rrbracket, e_1^m, e'_2, h_1^m, h'_2) \in \mathbf{safe}_{\text{int}, \emptyset}^{\emptyset, \emptyset, \text{dom}(w'), \emptyset} w, w'. \end{aligned}$$

The base case is easy by the definition of the logical relation: Pick $e'_2 = e_2, h'_2 = h_2$, and $w' = w$ and apply the assumption that $\cdot | \cdot | \cdot \models e_1 \preceq e_2 : \text{int}, \emptyset$.

For the inductive case, let $0 < m \leq n$ and unravel the assumptions for $m-1$: we have e'_2, h'_2, w' with $(e_2 | h_2) \mapsto^* (e'_2 | h'_2)$ and $\text{dom}(h_1^{m-1}) \supseteq \text{dom}_1(w')$, $\text{dom}(h'_2) \supseteq \text{dom}_2(w')$, and

$$(N_{m-1}, \llbracket \cdot \rrbracket, \llbracket \cdot \rrbracket, e_1^{m-1}, e'_2, h_1^{m-1}, h'_2) \in \mathbf{safe}_{\text{int}, \emptyset}^{\emptyset, \emptyset, \text{dom}(w'), \emptyset} w, w'.$$

By the overall setup, we know that $(e_1^{m-1} | h_1^{m-1}) \rightarrow_{\mu_m}^{n_m} (e_1^m | h_1^m)$ and the progress branch of safety provides for us: There are e''_2 and h''_2 with $(e'_2 | h'_2) \mapsto^* (e''_2 | h''_2)$ and hence $(e_2 | h_2) \mapsto^* (e''_2 | h''_2)$. There is w'' with $\text{dom}(h_1^m) \supseteq \text{dom}_1(w'')$ and $\text{dom}(h''_2) \supseteq \text{dom}_2(w'')$ and, finally,

$$(N_m, \llbracket \cdot \rrbracket, \llbracket \cdot \rrbracket, e_1^m, e''_2, h_1^m, h''_2) \in \mathbf{safe}_{\text{int}, \emptyset}^{\emptyset, \emptyset, \text{dom}(w''), \emptyset} w, w''.$$

We have now finished the induction proof; all that remains is to observe that the property proved implies the overall goal in the case $m = n$ by the termination branch of safety. ■

C. Properties of the Pre- and Postcondition Relations

Lemma 13 (Precondition Separation). *Assume that we have $\Pi \# \Lambda$, $h_1 = f_1 \upharpoonright_{\text{dom}_1^{\text{R}(\Lambda)}(w)}$, $g_1 = f_1 \upharpoonright_{\text{dom}_1^{\text{R}(\Pi)}(w)}$, $h_2 = f_2 \upharpoonright_{\text{dom}_2^{\text{R}(\Lambda)}(w)}$ and $g_2 = f_2 \upharpoonright_{\text{dom}_2^{\text{R}(\Pi)}(w)}$. Then it holds that*

$$\begin{aligned} (k, h_1, h_2) \in \mathbf{P}_\varepsilon^{\Lambda,R} w \wedge (k, g_1, g_2) \in \mathbf{P}_\varepsilon^{\Pi,R} w \\ \iff \\ (k, f_1, f_2) \in \mathbf{P}_\varepsilon^{\Pi \cup \Lambda, R} w \end{aligned}$$

Lemma 14 (Postcondition Separation). *Assume that we have $\Pi \# \Lambda$, $h_1 = f_1 \upharpoonright_{\text{dom}_1^{\text{R}(\Lambda)}(w)}$, $g_1 = f_1 \upharpoonright_{\text{dom}_1^{\text{R}(\Pi)}(w)}$, $h_2 = f_2 \upharpoonright_{\text{dom}_2^{\text{R}(\Lambda)}(w)}$, $g_2 = f_2 \upharpoonright_{\text{dom}_2^{\text{R}(\Pi)}(w)}$, $h'_1 = f'_1 \upharpoonright_{\text{dom}_1^{\text{R}(\Lambda)}(w')}$, $g'_1 = f'_1 \upharpoonright_{\text{dom}_1^{\text{R}(\Pi)}(w')}$, $h'_2 = f'_2 \upharpoonright_{\text{dom}_2^{\text{R}(\Lambda)}(w')}$ and $g'_2 = f'_2 \upharpoonright_{\text{dom}_2^{\text{R}(\Pi)}(w')}$. Then*

$$\begin{aligned} (k, h_1, h_2, h'_1, h'_2) \in \mathbf{Q}_\varepsilon^{\Lambda,R} w, w' \wedge \\ (k, g_1, g_2, g'_1, g'_2) \in \mathbf{Q}_\varepsilon^{\Pi,R} w, w' \\ \iff \\ (k, f_1, f_2, f'_1, f'_2) \in \mathbf{Q}_\varepsilon^{\Pi \cup \Lambda, R} w, w'. \end{aligned}$$

Lemma 15 (Precondition Composition).

$$\begin{aligned} (k, h_1, h_2) \in \mathbf{P}_\varepsilon^{\Theta,R} w \wedge (k, h_1, h_2, h'_1, h'_2) \in \mathbf{Q}_\varepsilon^{\Theta,R} w, w' \\ \implies \\ (k, h'_1, h'_2) \in \mathbf{P}_\varepsilon^{\Theta,R} w'. \end{aligned}$$

¹www.itu.dk/people/thamsborg/longcarnival.pdf

Lemma 16 (Postcondition Composition).

$$\begin{aligned} (k, h_1, h_2, h'_1, h'_2) &\in \mathbf{Q}_\varepsilon^{\Theta, R} w, w' \wedge \\ (k, h'_1, h'_2, h''_1, h''_2) &\in \mathbf{Q}_\varepsilon^{\Theta, R} w', w'' \\ \implies \\ (k, h_1, h_2, h''_1, h''_2) &\in \mathbf{Q}_\varepsilon^{\Theta, R} w, w''. \end{aligned}$$

D. Some Cases of Proof of Compatibility

Lemma 5. $\Pi \mid \Lambda, \rho \mid \Gamma \models e_1 \preceq e_2 : \tau, \varepsilon$ implies $\Pi \mid \Lambda \mid \Gamma \models e_1 \preceq e_2 : \tau, \varepsilon - \rho$ provided that $\rho \notin \text{FRV}(\Gamma, \tau)$.

Proof: Let $a \in \mathbb{N}$ be the required number of anonymous regions from the assumption. To prove the desired, we unroll the definition of the logical relation: choose $a + 1$, pick arbitrary $k^\circ \in \mathbb{N}$, $w^\circ \in \mathbf{W}$, $R : \Pi \cup \Lambda \hookrightarrow |w^\circ|$, $A^\circ \subseteq \text{dom}(w^\circ)$, $\gamma_1, \gamma_2 \in \mathcal{V}^{|\Gamma|}$ and $h_1^\circ, h_2^\circ, h_1'^\circ, h_2'^\circ \in \mathcal{H}$. Assume that $R(\text{FRV}(\varepsilon - \rho)) \subseteq \text{dom}(w^\circ)$, $A^\circ \# R(\Pi \cup \Lambda)$, $|A^\circ| \geq a$, $\forall r \in A^\circ. w^\circ(r) = \emptyset$, $(k^\circ, \gamma_1, \gamma_2) \in \llbracket \Gamma \rrbracket^{R, w^\circ}$, $h_1'^\circ \subseteq h_1^\circ$, $h_2'^\circ \subseteq h_2^\circ$, and $(k^\circ, h_1'^\circ, h_2'^\circ) \in \mathbf{P}_{\varepsilon - \rho}^{\Lambda, R} w^\circ$. We must show that $(k^\circ, h_1'^\circ, h_2'^\circ, e_1[\gamma_1/\Gamma], e_2[\gamma_2/\Gamma], h_1'', h_2'') \in \mathbf{safe}_{\tau, \varepsilon - \rho}^{\Pi, \Lambda, A^\circ, R} w^\circ, w^\circ$.

We need, obviously, to make use of the assumption of the lemma. To do so, we claim that to have

$$(k, h_1'^\circ, h_2'^\circ, e_1, e_2, h_1, h_2) \in \mathbf{safe}_{\tau, \varepsilon - \rho}^{\Pi, \Lambda, A, R} w^\circ, w \quad (5)$$

it suffices to know that

$$(k, h_1'^\circ, h_2'^\circ, e_1, e_2, h_1, h_2) \in \mathbf{safe}_{\tau, \varepsilon}^{\Pi, \Lambda, \rho, A \setminus \{r\}, R[\rho \mapsto r]} w^\circ, w \quad (6)$$

whenever $r \in A$ holds. Observe first, that if we can prove this, then we are done by an easy application of the assumption of the lemma. So all that remains is to prove the claim; this we do by well-founded induction on $k \in \mathbb{N}$.

Assume that the claim holds for all naturals strictly less than $k \in \mathbb{N}$; we will try to prove it for k . To prove (5) we proceed to unroll the definition of safety. Pick arbitrary $j \leq k$, $w' \in \mathbf{W}$ and $g_1, g_2, f_1, f_2 \in \mathcal{H}$. Assume that the prerequisites hold, i.e., that we have $\mathbf{envtran}^{\Pi, \Lambda, A, R} w, w'$, $(j, g_1, g_2) \in \mathbf{P}_{\varepsilon - \rho}^{\Pi, R} w'$ and $(g_1, h_1, f_1, g_2, h_2, f_2) \in \mathbf{splits}^{\Pi, \Lambda, A, R} w'$. Before we look into the termination and progress branches, observe that the w' and the g_1, g_2, f_1, f_2 match the prerequisites of safety for (6) as well, since $R(\Lambda) \cup A = R[\rho \mapsto r](\Lambda, \rho) \cup A \setminus \{r\}$.

So we follow the termination branch and assume that $\text{irr}(e_1 | g_1 \cdot h_1 \cdot f_1)$ holds. From (6) we get $e'_2, w'', h'_1, h'_2, g'_1$ and g'_2 with properties aplenty:

- $(e_2 | g_2 \cdot h_2 \cdot f_2) \xrightarrow{*} (e'_2 | g'_2 \cdot h'_2 \cdot f_2)$.
- $\mathbf{selftran}^{\Pi, \Lambda, \rho, A \setminus \{r\}, R[\rho \mapsto r]} w', w''$.
- $\emptyset = ((A \setminus \{r\}) \cap \text{dom}(w'')) \cup (\text{dom}(w'') \setminus \text{dom}(w'))$.
- $g_1 \cdot h_1 = g'_1 \cdot h'_1$.
- $(g'_1, h'_1, f_1, g'_2, h'_2, f_2) \in \mathbf{splits}^{\Pi, \Lambda, \rho, \emptyset, R[\rho \mapsto r]} w''$.
- $(j, e_1, e_2) \in \llbracket \Gamma \rrbracket^{R[\rho \mapsto r]} (w'')$.
- $(j, g_1, g_2, g'_1, g'_2) \in \mathbf{Q}_\varepsilon^{\Pi, R[\rho \mapsto r]} w', w''$.
- $\exists h''_1 \subseteq h'_1, h''_2 \subseteq h'_2. (j, h''_1, h''_2, h''_1, h''_2) \in \mathbf{Q}_\varepsilon^{\Lambda, \rho, R[\rho \mapsto r]} w^\circ, w''$.

Now note that we must have $r \in \text{dom}(w'')$ and so we define the world w''' by $w'' \rightarrow_{\text{mask}(r)} w'''$, i.e., we kill off region

r . So we need to discharge a number of obligations; most are straightforward, but the type of the resulting expressions as well as the postcondition on the local heaps take a bit of scrutiny: Note first that we have

$$\llbracket \tau \rrbracket^{R[\rho \mapsto r]} (w'') \subseteq \llbracket \tau \rrbracket^{R[\rho \mapsto r]} (w''') = \llbracket \tau \rrbracket^R (w'')$$

by type monotonicity and since $\rho \notin \text{FRV}(\tau)$; this means that the resulting expressions are indeed well-typed. Now define heaps $h''_1 \subseteq h'_1$ and $h''_2 \subseteq h'_2$ by demanding that $\text{dom}(h''_1) = \text{dom}_1^{R(\Lambda)}(w''')$ and $\text{dom}(h''_2) = \text{dom}_2^{R(\Lambda)}(w''')$, i.e., we restrict to the private parts of the local heaps, not including the locations (formerly) in region r . One easily verifies that this gives

$$(j, h_1^\circ, h_2^\circ, h''_1, h''_2) \in \mathbf{Q}_{\varepsilon - \rho}^{\Lambda, R} w^\circ, w''$$

and the termination branch is done.

Finally we get to the progress branch: we assume that there are e'_1, h''_1, μ and $n \leq j$ such that $(e_1 | g_1 \cdot h_1 \cdot f_1) \rightarrow_\mu^n (e'_1 | h''_1)$. From (6) we get $e'_2, w'', A', h'_1, h'_2, g'_1, g'_2$ with a range of properties:

- $(e_2 | g_2 \cdot h_2 \cdot f_2) \xrightarrow{*} (e'_2 | g'_2 \cdot h'_2 \cdot f_2)$.
- $\mathbf{selftran}^{\Pi, \Lambda, \rho, A \setminus \{r\}, R[\rho \mapsto r]} w', w''$.
- $A' = (A \setminus \{r\}) \cap \text{dom}(w'') \cup (\text{dom}(w'') \setminus \text{dom}(w'))$.
- $h''_1 = g'_1 \cdot h'_1 \cdot f_1$.
- $\mu \in \mathbf{effs}_{\varepsilon, h''_1}^{A', R[\rho \mapsto r]} w''$.
- $(g'_1, h'_1, f_1, g'_2, h'_2, f_2) \in \mathbf{splits}^{\Pi, \Lambda, \rho, A', R[\rho \mapsto r]} w''$.
- $(j - n, g_1, g_2, g'_1, g'_2) \in \mathbf{Q}_\varepsilon^{\Pi, R[\rho \mapsto r]} w', w''$.
- $(j - n, h_1^\circ, h_2^\circ, e'_1, e'_2, h'_1, h'_2) \in \mathbf{safe}_{\tau, \varepsilon}^{\Pi, \Lambda, \rho, A', R[\rho \mapsto r]} w^\circ, w''$.

Now we must have $r \in \text{dom}(w'')$ and letting $A'' = (A \cap \text{dom}(w'')) \cup (\text{dom}(w'') \setminus \text{dom}(w'))$ gives $r \in A''$ as well as $A'' \setminus \{r\} = A'$. With A'' as the choice of new anonymous regions, the obligations in proving the termination branch of (5) are easily met; we just remark that we rely on the induction hypothesis to establish

$$(j - n, h_1^\circ, h_2^\circ, e'_1, e'_2, h'_1, h'_2) \in \mathbf{safe}_{\tau, \varepsilon - \rho}^{\Pi, \Lambda, A'', R} w^\circ, w''.$$

■

There are many details to the proof, but the idea is simple: we simultaneously view the computation in two ways, both with ρ as a private region and with ρ masked out. From the safety of former, we then get safety of the latter. In the masked out case, ρ is no longer a private region variable and the region name r associated with ρ joins the anonymous regions. In most cases, however, (region names associated with) private regions and anonymous regions are treated the same, so there is work to do only when anonymous regions are considered in isolation.

Lemma 6. $\cdot \mid \Pi, \Lambda \mid \Gamma \models e_1 \preceq e_2 : \tau, \varepsilon$ implies $\Pi \mid \Lambda \mid \Gamma \models \text{atomic } e_1 \preceq \text{atomic } e_2 : \tau, \varepsilon$ if $\text{als } \varepsilon \subseteq \text{rds } \varepsilon \cap \text{wrs } \varepsilon$.

Proof: Let $a \in \mathbb{N}$ be the required number of anonymous regions from the assumption. To prove the desired, we unroll the definition of the logical relation: choose a ,

pick arbitrary $k^\circ \in \mathbb{N}$, $w^\circ \in \mathbf{W}$, $R : \Pi \cup \Lambda \hookrightarrow |w^\circ|$, $A \subseteq \text{dom}(w^\circ)$, $\gamma_1, \gamma_2 \in \mathcal{V}^{|\Gamma|}$ and $h_1^\circ, h_2^\circ, h_1^{\circ\prime}, h_2^{\circ\prime} \in \mathcal{H}$. Assume that $R(\text{FRV}(\varepsilon)) \subseteq \text{dom}(w^\circ)$, $A \# R(\Pi \cup \Lambda)$, $|A| \geq a$, $\forall r \in A. w^\circ(r) = \emptyset$, $(k^\circ, \gamma_1, \gamma_2) \in \llbracket \Gamma \rrbracket^{R, w^\circ}$, $h_1^{\circ\prime} \subseteq h_1^\circ$, $h_2^{\circ\prime} \subseteq h_2^\circ$, and $(k^\circ, h_1^{\circ\prime}, h_2^{\circ\prime}) \in \mathbf{P}_{\varepsilon}^{\Lambda, R, w^\circ}$. We must show that $(k^\circ, h_1^{\circ\prime}, h_2^{\circ\prime}, \text{atomic } e_1[\gamma_1/\Gamma], \text{atomic } e_2[\gamma_2/\Gamma], h_1^\circ, h_2^\circ) \in \text{safe}_{\tau, \varepsilon}^{\Pi, \Lambda, A, R, w^\circ, w^\circ}$.

We need to generalize a bit to handle the loopy behavior of atomic: we prove that for any $k \leq k^\circ$ and any $w \in \mathbf{W}$ with $\text{envtran}^{\Pi, \Lambda, A, R} w^\circ, w$ we have

$$(k, h_1^{\circ\prime}, h_2^{\circ\prime}, \text{atomic } e_1[\gamma_1/\Gamma], \text{atomic } e_2[\gamma_2/\Gamma], h_1^\circ, h_2^\circ) \in \text{safe}_{\tau, \varepsilon}^{\Pi, \Lambda, A, R} w^\circ, w.$$

This we do by well-founded induction on k . Unroll the definition of safety. Pick arbitrary $j \leq k$, $w' \in \mathbf{W}$ and $g_1, g_2, f_1, f_2 \in \mathcal{H}$. Assume that the prerequisites hold, i.e., that we have $\text{envtran}^{\Pi, \Lambda, A, R} w, w'$, $(j, g_1, g_2) \in \mathbf{P}_{\varepsilon}^{\Pi, R} w'$ and $(g_1, h_1, f_1, g_2, h_2, f_2) \in \text{splits}^{\Pi, \Lambda, A, R} w'$. Atomic commands always have the possibility of looping; hence we need not consider the termination branch. So we get to the progress branch: we assume that there are e_1', h_1^\dagger , μ and $n \leq j - 1$ such that $(\text{atomic } e_1[\gamma_1/\Gamma] | g_1 \cdot h_1 \cdot f_1) \xrightarrow{\mu^{n+1}} (e_1' | h_1^\dagger)$ and we must match this.

Observe first that $\text{envtran}^{\Pi, \Lambda, A, R} w^\circ, w'$ holds. If, now, the configuration loops to itself, we take no steps on the right hand side and conclude with the induction hypothesis. So we are left to consider the case where $(e_1[\gamma_1/\Gamma] | g_1 \cdot h_1 \cdot f_1) \xrightarrow{\mu^n} (e_1' | h_1^\dagger)$ and $e_1' \in \mathcal{V}$. Let $m \in \mathbb{N}$ be the number of reduction steps, we name the configurations $(e_1^0 | h_1^0), (e_1^1 | h_1^1), \dots, (e_1^m | h_1^m)$ in order and pick $n_1, n_2, \dots, n_m \in \mathbb{N}$ non-zero and $\mu_1, \mu_2, \dots, \mu_m$ such that

$$(e_1[\gamma_1/\Gamma] | g_1 \cdot h_1 \cdot f_1) = (e_1^0 | h_1^0) \xrightarrow{n_1} (e_1^1 | h_1^1) \xrightarrow{n_2} \dots \xrightarrow{n_m} (e_1^m | h_1^m) = (e_1' | h_1^\dagger)$$

with $n = n_1 + n_2 + \dots + n_m$ and $\mu = \mu_1 \cup \mu_2 \cup \dots \cup \mu_m$. For convenience, we furthermore write $N_j = \sum_{i=j+1}^m n_i$ for each $0 \leq j \leq m$.

It is time for the crux: for any $0 \leq j \leq m$ there are $w'', A' \subseteq \text{dom}(w'')$, h_1', e_2', h_2' such that

- $(e_2[\gamma_2/\Gamma] | g_2 \cdot h_2 \cdot f_2) \xrightarrow{*} (e_2' | h_2' \cdot f_2)$
- $\text{selftran}^{\Pi, \Lambda, A, R} w', w''$
- $h_1' \cdot f = h_1^j$
- $\bigcup_{i=1}^j \mu_i \setminus (\text{dom}_1^{R(\Pi \cup \Lambda)}(w'') \setminus \text{dom}_1^{R(\Pi \cup \Lambda)}(w')) \in \text{effs}_{\varepsilon, h_1'}^{A', R} w''$
- $(\llbracket, h_1', f_1, \llbracket, h_2', f_2) \in \text{splits}^{\Pi, \Lambda, A', R} w''$
- $(N_j + j - n, g_1 \cdot h_1, g_2 \cdot h_2, e_1^j, e_2', h_1', h_2') \in \text{safe}_{\tau, \varepsilon}^{\Pi, \Lambda, A', R} w', w''$.

Notice the abuse of notation in item four: we do not really remove locations, we remove any actual effects tagged with the listed locations. The base case follows from the assumption of the lemma; the induction simply by unrolling the safety. In the end, we are able to use the properties in the case $j = m$ to produce a right hand side reduction in the overall proof. There

are many details to this; here we just remark that it would be more more pleasant to have

$$\bigcup_{i=1}^j \mu_i \in \text{effs}_{\varepsilon, h_1'}^{A', R} w''$$

as item four. Unfortunately, that would break in the inductive step, and we are stuck with the more complex version, which again forces the side condition on the rule. ■

Lemma 7. $\Pi, \Lambda | \cdot | \Gamma \models e_1 \preceq e_2 : \tau, \varepsilon$ and $\Pi, \Lambda | \cdot | \Gamma \models e_1^\dagger \preceq e_2^\dagger : \tau^\dagger, \varepsilon^\dagger$ together implies $\Pi | \Lambda | \Gamma \models \text{par } e_1$ and $e_1^\dagger \preceq \text{par } e_2$ and $e_2^\dagger : \tau \times \tau^\dagger, \varepsilon \cup \varepsilon^\dagger$.

Proof: Let $a_1, a_2 \in \mathbb{N}$ be the required numbers of anonymous regions from the respective assumptions. To prove the desired, we unroll the definition of the logical relation: choose $a_1 + a_2$, pick arbitrary $k^\circ \in \mathbb{N}$, $w^\circ \in \mathbf{W}$, $R : \Pi \cup \Lambda \hookrightarrow |w^\circ|$, $A^\circ \subseteq \text{dom}(w^\circ)$, $\gamma_1, \gamma_2 \in \mathcal{V}^{|\Gamma|}$ and $h_1^\circ, h_2^\circ, h_1^{\circ\prime}, h_2^{\circ\prime} \in \mathcal{H}$. Assume that $R(\text{FRV}(\varepsilon - \rho)) \subseteq \text{dom}(w^\circ)$, $A^\circ \# R(\Pi \cup \Lambda)$, $|A^\circ| \geq a_1 + a_2$, $\forall r \in A^\circ. w^\circ(r) = \emptyset$, $(k^\circ, \gamma_1, \gamma_2) \in \llbracket \Gamma \rrbracket^{R, w^\circ}$, $h_i^{\circ\prime} \subseteq h_i^\circ$, and $(k^\circ, h_1^{\circ\prime}, h_2^{\circ\prime}) \in \mathbf{P}_{\varepsilon_1 \cup \varepsilon_2}^{\Lambda, R, w^\circ}$. We must show that $(k^\circ, h_1^{\circ\prime}, h_2^{\circ\prime}, \text{par } e_1 \text{ and } e_1^\dagger[\gamma_1/\Gamma], \text{par } e_2 \text{ and } e_2^\dagger[\gamma_2/\Gamma], h_1^\circ, h_2^\circ) \in \text{safe}_{\tau_1 \times \tau_2, \varepsilon_1 \cup \varepsilon_2}^{\Pi, \Lambda, A^\circ, R} w^\circ, w^\circ$.

We need, obviously, to make use of the assumptions of the lemma. To this end, we claim that to have

$$(k, h_1^{\circ\prime}, h_2^{\circ\prime}, \text{par } e_1^1 \text{ and } e_1^2, \text{par } e_2^1 \text{ and } e_2^2, h_1 \cdot h_1^1 \cdot h_1^2, h_2 \cdot h_2^1 \cdot h_2^2) \in \text{safe}_{\tau_1 \times \tau_2, \varepsilon_1 \cup \varepsilon_2}^{\Pi, \Lambda, A_1 \cup A_2, R} w^\circ, w, \quad (7)$$

it suffices to know that

$$(k, h_1, h_2) \in \mathbf{P}_{\varepsilon_1 \cup \varepsilon_2}^{\Lambda, R} w \quad (8)$$

$$(h, h_1^{\circ\prime}, h_2^{\circ\prime}, h_1, h_2) \in \mathbf{Q}_{\varepsilon_1 \cup \varepsilon_2}^{\Lambda, R} w^\circ, w \quad (9)$$

$$(k, \emptyset, \emptyset, e_1^1, e_2^1, h_1^1, h_2^1) \in \text{safe}_{\tau_1, \varepsilon_1}^{\Pi \cup \Lambda, A_1, R} w^\circ, w \quad (10)$$

$$(k, \emptyset, \emptyset, e_1^2, e_2^2, h_1^2, h_2^2) \in \text{safe}_{\tau_2, \varepsilon_2}^{\Pi \cup \Lambda, A_2, R} w^\circ, w, \quad (11)$$

whenever $A_1 \# A_2$ and $\text{dom}_i^{A_j}(w) \subseteq \text{dom}(h_i^j)$ holds. Observe, that if we can prove this, then we are done by splitting A° into two parts of appropriate sizes followed by an easy application of the assumptions of the lemma. The remaining part is to prove the claim, this we do by complete induction on k .

Assume the claim holds for all naturals strictly less than $k \in \mathbb{N}$; we will now prove it also holds for k . To prove 7 we unroll the definition of safety. Pick arbitrary $j \leq k$, $w' \in \mathbf{W}$ and $g_1, g_2, f_1, f_2 \in \mathcal{H}$. Assume the prerequisites hold, i.e., that we have $\text{envtran}^{\Pi, \Lambda, A_1 \cup A_2, R} w, w'$, $(j, g_1, g_2) \in \mathbf{P}_{\varepsilon_1 \cup \varepsilon_2}^{\Pi, R} w'$ and $(g_1, h_1 \cdot h_1^1 \cdot h_1^2, f_1, g_2, h_2 \cdot h_2^1 \cdot h_2^2, f_2) \in \text{splits}^{\Pi, \Lambda, A_1 \cup A_2, R} w'$. Observe that we can instantiate assumptions 10 and 11 by taking $g_i \cdot h_i$ as the public part of the heap, and adding the spare private part to the frame.

We follow the termination branch first, and assume that $\text{irr}(\text{par } e_1^1 \text{ and } e_1^2 | g_1 \cdot h_1 \cdot h_1^1 \cdot h_1^2)$. This means each of the subexpressions is also irreducible so from safety assumptions we learn, among other things, that $e_1^1, e_1^2 \in \mathcal{V}$. However, this means that $(\text{par } e_1^1 \text{ and } e_1^2 | g_1 \cdot h_1 \cdot h_1^1 \cdot h_1^2) \mapsto ((e_1^1, e_1^2) | g_1 \cdot h_1 \cdot h_1^1 \cdot h_1^2)$, which contradicts the irreducibility assumption.

For the progress branch, we are left with three possibilities of reduction: either both e_1^1 and e_1^2 are values, or the reduction happens inside one of them. We omit the case for reduction inside e_1^2 , since it is completely symmetric to the one for e_1^1 ; the other two cases follow.

Assume that $(e_1^1 | g_1 \cdot h_1 \cdot h_1^1 \cdot h_1^2 \cdot f_1) \rightarrow_{\mu}^n (e_1^1 | h_1^{\dagger})$ for some $n \leq j$. From 10 we get $e_2^1, w'', h_1', h_2', g_1', g_2', h_1^1$ and h_2^1 , along with the following properties:

- $(e_2^1 | g_2 \cdot h_2 \cdot h_2^1 \cdot h_2^2 \cdot f_2) \mapsto^* (e_2^1 | g_2' \cdot h_2' \cdot h_2^1 \cdot h_2^2 \cdot f_2)$,
- **selftran** $_{\Pi \cup \Lambda, \cdot, A_1, R} w', w''$,
- $A_1' = (A_1 \cap \text{dom}(w'')) \cup (\text{dom}(w'') \setminus \text{dom}(w'))$,
- $h_1^{\dagger} = g_1' \cdot h_1' \cdot h_1^1 \cdot h_1^2 \cdot f_1$,
- $\mu \in \text{effs}_{\varepsilon_2, h_1^1}^{A_1, R} w''$
- $(g_1' \cdot h_1', h_1^1, f_1 \cdot h_1^2, g_2' \cdot h_2', h_2^1, f_2 \cdot h_2^2) \in \text{splits}_{\Pi \cup \Lambda, \cdot, A_2, R} w''$,
- $(j - n, g_1 \cdot h_1, g_2 \cdot h_2, g_1' \cdot h_1', g_2' \cdot h_2') \in \mathbf{Q}_{\varepsilon_1}^{\Pi \cup \Lambda, R} w', w''$,
- $(j - n, \emptyset, \emptyset, e_1^1, e_2^1, h_1^1, h_1^2) \in \text{safe}_{\tau_1, \varepsilon_1}^{\Pi \cup \Lambda, \cdot, A_2, R} w^{\circ}, w''$.

Obviously, we can now replay the reduction sequence for the whole right-hand-side expression. By postcondition separation and postcondition weakening lemmas, we can also split up the postcondition into separate parts for Π and Λ , the first of which is needed by progress branch. The other obligations of the progress branch are simple, leaving us with the final safety requirement: $(j - n, h_1^{\circ}, h_2^{\circ}, \text{par } e_1^1 \text{ and } e_2^1, \text{par } e_2^1 \text{ and } e_2^2, h_1^1 \cdot h_1^2 \cdot h_2^1 \cdot h_2^2) \in \text{safe}_{\tau_1 \times \tau_2, \varepsilon_1 \cup \varepsilon_2}^{\Pi, \Lambda, A_1 \cup A_2, R} w^{\circ}, w''$. This follows by induction hypothesis, since $n > 0$, if we can show the four assumptions. 8 and 9 hold due to downwards-closure and composition lemmas for pre- and post-condition, and the Λ part of the assumption above, and 10 we have verbatim as an assumption. This leaves us only 11 to show. To show it, notice that safety is both downwards closed in the step index, and closed under environment transitions and observe that we have **envtran** $_{\Pi \cup \Lambda, \cdot, A_2, R} w, w''$ by composing the two world transitions above. This, in conjunction with the original assumption, suffices to finish this part of the proof.

The final case we consider is when $e_1^1, e_2^1 \in \mathcal{V}$ and $(\text{par } e_1^1 \text{ and } e_2^1 | g_1 \cdot h_1 \cdot h_1^1 \cdot h_1^2) \rightarrow_{\emptyset} (\langle e_1^1, e_2^1 \rangle | g_1 \cdot h_1 \cdot h_1^1 \cdot h_1^2)$. In this case, we start with instantiating 10 and from the termination branch obtain:

- $(e_2^1 | g_2 \cdot h_2 \cdot h_2^1 \cdot h_2^2 \cdot f_2) \mapsto^* (e_2^1 | g_2' \cdot h_2' \cdot h_2^1 \cdot h_2^2 \cdot f_2)$,
- **selftran** $_{\Pi \cup \Lambda, \cdot, A_1, R} w', w''$,
- $\emptyset = (A_1 \cap \text{dom}(w'')) \cup (\text{dom}(w'') \setminus \text{dom}(w'))$,
- $g_1 \cdot h_1 \cdot h_1^1 = g_1' \cdot h_1' \cdot h_1^1$,
- $(j, e_1^1, e_2^1) \in \llbracket \tau_1 \rrbracket^R (w'')$,
- $(j, g_1 \cdot h_1, g_2 \cdot h_2, g_1' \cdot h_1', g_2' \cdot h_2') \in \mathbf{Q}_{\varepsilon_1}^{\Pi \cup \Lambda, R} w', w''$.

Now, with the new heaps, we can instantiate 11, and get:

- $(e_2^2 | g_2' \cdot h_2' \cdot h_2^1 \cdot h_2^2 \cdot f_2) \mapsto^* (e_2^2 | g_2'' \cdot h_2'' \cdot h_2^1 \cdot h_2^2 \cdot f_2)$,
- **selftran** $_{\Pi \cup \Lambda, \cdot, A_2, R} w'', w'''$,
- $\emptyset = (A_2 \cap \text{dom}(w''')) \cup (\text{dom}(w''') \setminus \text{dom}(w''))$,
- $g_1' \cdot h_1' \cdot h_1^2 = g_1'' \cdot h_1'' \cdot h_1^2$,
- $(j, e_2^1, e_2^2) \in \llbracket \tau_2 \rrbracket^R (w''')$,
- $(j, g_1' \cdot h_1', g_2' \cdot h_2', g_1'' \cdot h_1'', g_2'' \cdot h_2'') \in \mathbf{Q}_{\varepsilon_2}^{\Pi \cup \Lambda, R} w'', w'''$.

Now we can build the complete reduction for the right-hand side: $(\text{par } e_2^1 \text{ and } e_2^2 | g_2 \cdot h_2 \cdot h_2^1 \cdot h_2^2 \cdot f_2) \mapsto^*$

$((e_2^1, e_2^2) | g_2'' \cdot h_2'' \cdot h_2^1 \cdot h_2^2 \cdot f_2)$, and prove the rest of the required obligations: self-transition out of the composition of the two given, postcondition of public heap fragments by postcondition separation, composition and weakening, and the rest by simple manipulations. This leaves the final safety of two pairs to be proven. We unroll the definition again, taking $i \leq j - 1$, this time only considering the termination branch, since the left-hand side is a value, and do not do any reductions. This leaves us with proving two interesting obligations. The first one is $(i, \langle e_1^1, e_2^1 \rangle, \langle e_2^1, e_2^2 \rangle) \in \llbracket \tau_1 \times \tau_2 \rrbracket^R (w^{\dagger})$. However, since $w^{\dagger} \sqsupseteq w'''$, we can easily show this by definition of the denotation of product types, downwards- and future-world-closure, and the value assumptions obtained above. The other interesting obligation is showing the postcondition for the local, Λ part of the heap. This is done by using postcondition separation, composition and weakening on the postcondition properties obtained from assumptions and by assumption 9 itself. ■

E. Proof of the Parallelization Theorem

We start by stating and proving the catch-up lemma mentioned in the explanation in the paper, that covers most of the proof of the theorem.

Lemma 17 (Catch-up). *Assuming that*

- 1) $(k, h_1, h_2) \in \mathbf{P}_{\varepsilon_1}^{\Lambda, R} w$,
- 2) $(k, \emptyset, e_1, \tilde{e}_1, h_{11}, h_{12}) \in \text{safe}_{\tau_1, \varepsilon_1}^{\Pi \cup \Lambda, \cdot, A_1, R} w^{\circ}, w$ and
- 3) *for any* $g_1, g_2, h_1^{\dagger}, h_2^{\dagger}, f_1, f_2, w', j \leq k$ *such that*

- **envtran** $_{\Pi \cup \Lambda, \cdot, A_2, R} w, w'$
- $(g_1 \cdot h_1^{\dagger}, h_{21}, f_1, g_2 \cdot h_2^{\dagger}, \emptyset, f_2) \in \text{splits}_{\Pi \cup \Lambda, \cdot, A_2, R} w'$
- $(j, h_1, h_2, h_1^{\dagger}, h_2^{\dagger}) \in \mathbf{Q}_{\varepsilon_1}^{\Lambda, R} w, w'$

there exist $h_1', h_2', g_1', g_2', h_{22}, h_{21}', \tilde{e}_1', w''$ *such that*

- **selftran** $_{\Pi \cup \Lambda, \cdot, A_2, R} w', w''$
- $(\tilde{e}_2 | h_2^{\dagger} \cdot g_2 \cdot f_2) \mapsto^* (\tilde{e}_2' | h_2' \cdot g_2' \cdot h_{22} \cdot f_2)$
- $h_1^{\dagger} \cdot g_1 \cdot h_{21} = h_1' \cdot g_1' \cdot h_{21}'$
- $(k, g_1, g_2, g_1', g_2') \in \mathbf{Q}_{\varepsilon_2}^{\Pi, R} w', w''$
- $(k, h_1', h_2') \in \mathbf{P}_{\varepsilon_2}^{\Lambda, R} w''$
- $(j, h_1^{\circ}, h_2^{\circ}, h_1', h_2') \in \mathbf{Q}_{\varepsilon_1 \cup \varepsilon_2}^{\Lambda, R} w^{\circ}, w''$
- $(k, \emptyset, \emptyset, e_2, e_2', h_{21}', h_{22}) \in \text{safe}_{\tau_2, \varepsilon_2}^{\Pi \cup \Lambda, \cdot, A_2', R} w^{\circ}, w''$,

we have $(k, h_1^{\circ}, h_2^{\circ}, \text{par } e_1 \text{ and } e_2, \langle \tilde{e}_1, \tilde{e}_2 \rangle, h_1 \cdot h_{11} \cdot h_{21}, h_2 \cdot h_{12}) \in \text{safe}_{\tau_1 \times \tau_2, \varepsilon_1 \cup \varepsilon_2}^{\Pi, \Lambda, A_1 + A_2, R} w^{\circ}, w$, *as long as* $\text{rds } \varepsilon_1 \cup \text{wrs } \varepsilon_1 \cup \text{rds } \varepsilon_2 \cup \text{wrs } \varepsilon_2 \subseteq \Lambda$, $\text{rds } \varepsilon_1 \# \text{wrs } \varepsilon_2$, $\text{rds } \varepsilon_2 \# (\text{wrs } \varepsilon_1 \cup \text{als } \varepsilon_1)$, $\text{wrs } \varepsilon_1 \# \text{wrs } \varepsilon_2$, $\text{dom}(h_{i1}) \supseteq \text{dom}^{A_1}(w)$, $\text{dom}(h_{i2}) \supseteq \text{dom}^{A_2}(w)$ *and* $\text{dom}^{A_1}(w) \text{ and } \text{dom}^{A_2}(w) = \emptyset$.

Proof: The proof proceeds by well-founded induction on k . Take any $j \leq k, g_1, g_2, f_1, f_2, w_1$ such that

- **envtran** $_{\Pi, \Lambda, A_1 + A_2, R} w, w_1$
- $(j, g_1, g_2) \in \mathbf{P}_{\varepsilon_1 \cup \varepsilon_2}^{\Pi, R} w_1$
- $(g_1, h_1 \cdot h_{11} \cdot h_{21}, f_1, g_2, h_2 \cdot h_{12}, f_2) \in \text{splits}_{\Pi, \Lambda, A_1 + A_2, R} w_1$

There are two branches to consider. However, $\text{par } e_1$ and e_2 can always reduce, so the termination branch is trivial. In the progress case, we proceed by case analysis on the reduction to get three subcases.

1) $(e_1 | g_1 \cdot h_1 \cdot h_{11} \cdot h_{21} \cdot f_1) \rightarrow_\mu^n (e'_1 | h'_1)$

In this case we can use the assumption (2), which we instantiate with $j \leq k, g_1 \cdot h_1, g_2 \cdot h_2, f_1 \cdot h_{21}, f_2, w_1$. We need to show the prerequisites:

- $\mathbf{envtran}^{\Pi\cup\Lambda, \cdot, A_1, R} w, w_1$, which holds by env-transition weakening
- $(j+1, g_1 \cdot h_1, g_2 \cdot h_2) \in \mathbf{P}_{\varepsilon_1}^{\Pi\cup\Lambda, R} w_1$, which holds by assumptions, precondition composition, precondition weakening and future-world closure of precondition
- $(g_1 \cdot h_1, h_{11}, f_1 \cdot h_{21}, g_2 \cdot h_2, h_{12}, f_2) \in \mathbf{splits}^{\Pi\cup\Lambda, \cdot, A_1, R} w_1$, which holds by massaging the assumption about splits.

From the progress branch we obtain the following properties:

- $(\tilde{e}_1 | g_2 \cdot h_2 \cdot h_{12} \cdot f_2) \mapsto^* (\tilde{e}'_1 | g'_2 \cdot h'_2 \cdot h'_{12} \cdot f_2)$
- $\mathbf{selftran}^{\Pi\cup\Lambda, \cdot, A_1, R} w_1, w_2$
- $A'_1 = (A_1 \cap \text{dom}(w_2)) \cup (\text{dom}(w_2) \setminus \text{dom}(w_1))$
- $h'_1 \dagger = g'_1 \cdot h'_1 \cdot h'_{11} \cdot f_1 \cdot h_{21}$
- $\mu \in \mathbf{effs}_{\varepsilon_1, h'_1}^{A'_1, R} w_2$
- $(g'_1 \cdot h'_1, h'_{11}, f_1 \cdot h_{21}, g'_2 \cdot h'_2, h'_{12}, f_2) \in \mathbf{splits}^{\Pi\cup\Lambda, \cdot, A_1, R} w_2$
- $(j-n, g_1 \cdot h_1, g_2 \cdot h_2, g'_1 \cdot h'_1, g'_2 \cdot h'_2) \in \mathbf{Q}_{\varepsilon_1}^{\Pi\cup\Lambda, R} w_1, w_2$
- $(j-n, \emptyset, \emptyset, e'_1, \tilde{e}'_1, h'_{11}, h'_{12}) \in \mathbf{safe}_{\tau_1, \varepsilon_1}^{\Pi\cup\Lambda, \cdot, A'_1, R} w^\circ, w_2$

From these we can conclude that:

- $(\langle \tilde{e}_1, \tilde{e}_2 \rangle | g_2 \cdot h_2 \cdot h_{12} \cdot f_2) \mapsto^* (\langle \tilde{e}'_1, \tilde{e}'_2 \rangle | g'_2 \cdot h'_2 \cdot h'_{12} \cdot f_2)$ (recall pairs of expressions is syntactic sugar)
- $\mathbf{selftran}^{\Pi, \Lambda, A_1+A_2, R} w_1, w_2$, by self-transition strengthening
- $A'_1 + A_2 = ((A_1 + A_2) \cap \text{dom}(w_2)) \cup (\text{dom}(w_2) \setminus \text{dom}(w_1))$, by self-transition restrictions
- $(g'_1, h'_1 \cdot h'_{11} \cdot h_{21}, f_1, g'_2, h'_2 \cdot h'_{12}, f_2) \in \mathbf{splits}^{\Pi, \Lambda, A'_1+A_2, R} w_2$, by massaging the splits assumption
- $\mu \in \mathbf{effs}_{\varepsilon_1 \cup \varepsilon_2, h'_1 \cdot h'_{11} \cdot h_{21}}^{A'_1+A_2, R} w_2$, by effs strengthening
- $(j-n, g_1, g_2, g'_1, g'_2) \in \mathbf{Q}_{\varepsilon_1 \cup \varepsilon_2}^{\Pi, R} w_1, w_2$, by postcondition separation and postcondition weakening
- $(j-n, h'_1, h'_2) \in \mathbf{P}_{\varepsilon_1}^{\Lambda, R} w_2$, by postcondition separation and precondition composition
- $\mathbf{envtran}^{\Pi\cup\Lambda, \cdot, A_2, R} w, w_2$, by envtransition composition and relation of env-transition and self-transition
- $(j-n, h_1, h_2, h'_1, h'_2) \in \mathbf{Q}_{\varepsilon_1}^{\Lambda, R} w, w_2$, by postcondition separation

Since $0 < n \leq j$, we can now use the induction hypothesis to discharge the final safety obligation: we have already obtained preconditions (1) and (2), while (3) is closed under the world transition that we have made (see the final two properties above).

2) $(e_2 | g_1 \cdot h_1 \cdot h_{11} \cdot h_{21} \cdot f_1) \rightarrow_\mu^n (e'_2 | h'_1)$

First, we need to establish that $h'_1 \dagger = g_1 \cdot h'_1 \cdot h'_{21} \cdot f_1 \cdot h_{11}$, $\text{dom}(h'_1) = \text{dom}(h_1)$, $\mu \in \mathbf{effs}_{\varepsilon_2, h'_{21}}^{A, R} w$ and $(j-n, h'_1, h_2) \in \mathbf{P}_{\varepsilon_1}^{\Lambda, R} w_1$. This can be obtained by using the assumption (3), as follows. Take $g_1, g_2, h_1, h_2, f_1 \cdot h_{11}, f_2 \cdot h_{12}, w_1, k \leq k$ as the universally quantified variables in (3). We need to show:

$n, h'_1, h_2) \in \mathbf{P}_{\varepsilon_1}^{\Lambda, R} w_1$. This can be obtained by using the assumption (3), as follows. Take $g_1, g_2, h_1, h_2, f_1 \cdot h_{11}, f_2 \cdot h_{12}, w_1, k \leq k$ as the universally quantified variables in (3). We need to show:

- $\mathbf{envtran}^{\Pi\cup\Lambda, \cdot, A_2, R} w_1, w_1$, which holds by reflexivity
- $(g_1 \cdot h_1, h_{21}, f_1 \cdot h_{11}, g_2 \cdot h_2, \emptyset, f_2 \cdot h_{12}) \in \mathbf{splits}^{\Pi\cup\Lambda, \cdot, A_2, R} w_1$, by massaging the splits we have assumed earlier
- $(k, h_1, h_2, h_1, h_2) \in \mathbf{Q}_{\varepsilon_1}^{\Lambda, R} w_1, w_1$, again, by reflexivity.

From this we obtain a bunch of properties, by unfolding safety. However, we are only interested in a few, namely that $h'_1 \dagger = g_1 \cdot h'_1 \cdot h'_{21} \cdot f_1 \cdot h_{11}$, and $\mu \in \mathbf{effs}_{\varepsilon_2, h'_1}^{A, R} w^\dagger$ and $\mathbf{selftran}^{\Pi\cup\Lambda, \cdot, A_2, R} w_1, w^\dagger$. Let $h'_1 \dagger = g_1 \cdot h'_1 \cdot h'_{21}$.

Now, as witnesses for the progress case, we take $e_2, w_1, A_1 + A_2, h'_1 \cdot h_{11} \cdot h'_{21}, h_2 \cdot h_{12}, g_1, g_2$, where $h'_1 = h'_1 \dagger \upharpoonright_{\text{dom}_1^R(w_1)}$ and $h'_{21} = h'_1 \dagger \upharpoonright_{\text{dom}_1(w_1) \setminus \text{dom}_1^R(w_1)}$. It suffices to show:

- $(\tilde{e}_2 | g_2 \cdot h_2 \cdot h_{12} \cdot f_2) \mapsto^* (\tilde{e}'_2 | g_2 \cdot h_2 \cdot h_{12} \cdot f_2)$ by reflexivity
- $\mathbf{selftran}^{\Pi, \Lambda, A_1+A_2, R} w_1, w_1$ by reflexivity
- $h'_1 \dagger = g_1 \cdot h'_1 \cdot h_{11} \cdot h'_{21} \cdot f_1$ by the restriction on actual effects (μ) we know g_1 didn't change, and we have an assumption to finish this off
- $(g_1, h'_1 \cdot h_{11} \cdot h'_{21}, f_1, g_2, h_2 \cdot h_{12}, f_2) \in \mathbf{splits}^{\Pi, \Lambda, A_1+A_2, R} w_1$ by assumption
- $\mu \in \mathbf{effs}_{\varepsilon_2, h'_{21}}^{A_2, R} w_1$, since any region killed between w_1 and w^\dagger has to be in A_2 , $\text{dom}(h'_1) \subseteq \text{dom}(h'_{21})$, and we have an assumption for the other effects inside the world
- $(j-n, g_1, g_2, g_1, g_2) \in \mathbf{Q}_{\varepsilon_1 \cup \varepsilon_2}^{\Pi, R} w_1, w_1$ by reflexivity,

so it remains to show that $(j-n, h^\circ_1, h^\circ_2, \text{par } e_1 \text{ and } e'_2, \langle \tilde{e}_1, \tilde{e}_2 \rangle, h'_1 \cdot h_{11} \cdot h'_{21}, h_2 \cdot h_{12}) \in \mathbf{safe}_{\tau_1 \times \tau_2, \varepsilon_1 \cup \varepsilon_2}^{\Pi, \Lambda, A_1+A_2, R} w^\circ, w_1$ holds. We can use induction hypothesis to prove it, provided we show the three assumptions. The first one $((j-n, h'_1, h_2) \in \mathbf{P}_{\varepsilon_1}^{\Lambda, R} w_1)$ is simple, since μ is confined to ε_2 , the second one $((j-n, \emptyset, \emptyset, e_1, \tilde{e}_1, h_{11}, h_{12}) \in \mathbf{safe}_{\tau_1, \varepsilon_1}^{\Pi\cup\Lambda, \cdot, A_1, R} w^\circ, w_1)$ holds by env-transition closure and downwards closure of \mathbf{safe} . This leaves us the last precondition to show. To this end, we take $h'_1 \dagger, h'_2 \dagger, g_1 \dagger, g_2 \dagger, f_1 \dagger, f_2 \dagger, i \leq j-n$ and w_2 , such that

- $\mathbf{envtran}^{\Pi\cup\Lambda, \cdot, A_2, R} w_1, w_2$
- $(i, h'_1 \dagger, h_2 \dagger, h'_1 \dagger, h'_2 \dagger) \in \mathbf{Q}_{\varepsilon_1}^{\Lambda, R} w_1, w_2$
- $(g_1 \dagger \cdot h'_1 \dagger, h'_{21} \dagger, f_1 \dagger, g_2 \dagger \cdot h'_2 \dagger, \emptyset, f_2 \dagger) \in \mathbf{splits}^{\Pi\cup\Lambda, \cdot, A_2, R} w_2$

Now we can instantiate the original assumption (3) with $h_1[h'_1 \dagger / h_1], h_2 \dagger, g_1 \dagger, g_2 \dagger, f_1 \dagger, f_2 \dagger, i \leq k$ and w_2 , where the heap substitution notation means that any change (allocation or update) from $h'_1 \dagger$ to $h_1 \dagger$ should be replayed on h_1 (this is well-specified, since $\text{dom}(h_1) = \text{dom}(h'_1)$). We need to prove the following:

- $\mathbf{envtran}^{\Pi\cup\Lambda, \cdot, A_2, R} w, w_2$, by envtran-comp and

envtran-weaken

- $(i, h_1, h_2, h_1[h_1^\dagger/h_1'], h_2^\dagger) \in \mathbf{Q}_{\varepsilon_1}^{\Lambda, R} w, w_2$, by definition of heap substitution, and assumption
- $(g_1^\dagger \cdot h_1[h_1^\dagger/h_1'], h_{21}, f_1^\dagger, g_2^\dagger \cdot h_2^\dagger, \emptyset, f_2^\dagger) \in \mathbf{splits}_{\Pi\cup\Lambda, \cdot, A_2, R} w_2$ by assumption, given that $\text{dom}(h_1^\dagger) = \text{dom}(h_1[h_1^\dagger/h_1'])$

By proving these properties, we obtain the following:

- $(k, h_1^\dagger, h_2^\dagger) \in \mathbf{P}_{\varepsilon_2}^{\Lambda, R} w_3$
- $\mathbf{selftran}_{\Pi\cup\Lambda, \cdot, A_2, R} w_2, w_3$
- $(\tilde{e}_2|h_2^\dagger \cdot g_2^\dagger \cdot f_2^\dagger) \mapsto^* (e_2'|h_2^\dagger \cdot g_2^\dagger \cdot h_{22} \cdot f_2^\dagger)$
- $h_1[h_1^\dagger/h_1'] \cdot g_1^\dagger \cdot h_{21} = h_1^\dagger \cdot g_1^\dagger \cdot h_{21}^\dagger$
- $(k, g_1^\dagger, g_2^\dagger, g_1^\dagger, g_2^\dagger) \in \mathbf{Q}_{\varepsilon_2}^{\Pi, R} w_2, w_3$
- $(i, h_1^\circ, h_2^\circ, h_1^\dagger, h_2^\dagger) \in \mathbf{Q}_{\varepsilon_1 \cup \varepsilon_2}^{\Lambda, R} w^\circ, w_3$
- $(g_1^\dagger \cdot h_1^\dagger, h_{21}, f_1^\dagger, g_2^\dagger \cdot h_2^\dagger, h_{22}, f_2^\dagger) \in \mathbf{splits}_{\Pi\cup\Lambda, \cdot, A_2, R} w_3$
- $(k, \emptyset, \emptyset, e_2, e_2', h_{21}, h_{22}) \in \mathbf{safe}_{\tau_2, \varepsilon_2}^{\Pi\cup\Lambda, \cdot, A_2, R} w^\circ, w_3$

Recall that $\mu \in \mathbf{effs}_{\varepsilon_2, h_{21}}^{A_2, R} w_1$, and $(e_2 | g_1 \cdot h_1 \cdot h_{11} \cdot h_{21} \cdot f_1) \rightarrow_\mu^n (e_2' | g_1 \cdot h_1' \cdot h_{21}' \cdot f_1 \cdot h_{11})$. Hence, we can use Lemma 2 to get

$$(e_2 | g_1^\dagger \cdot h_1[h_1^\dagger/h_1'] \cdot h_{21} \cdot f_1^\dagger) \rightarrow_\mu^n (e_2' | g_1^\dagger \cdot h_1^\dagger \cdot h_{21}' \cdot f_1^\dagger),$$

which allows us to instantiate the safety predicate above. Note that this argument is precisely the commutation of α and β in the explanation in Section IV-A (cf. Figure 10). We need to show:

- $\mathbf{envtran}_{\Pi\cup\Lambda, \cdot, A_2, R} w_3, w_3$, by reflexivity
- $(j, g_1^\dagger, g_2^\dagger) \in \mathbf{P}_{\varepsilon_2}^{\Pi, R} w_3$, trivial, since $\text{rds } \varepsilon_2 \# \Pi$
- $(j, h_1^\dagger, h_2^\dagger) \in \mathbf{P}_{\varepsilon_2}^{\Lambda, R} w_3$, by assumption
- $(g_1^\dagger \cdot h_1^\dagger, h_{21}, f_1 \cdot h_{11} \cdot f_1^\dagger, g_2^\dagger \cdot h_2^\dagger, h_{22}, f_2^\dagger) \in \mathbf{splits}_{\Pi\cup\Lambda, \cdot, A_2, R} w_3$, by massaging the assumption.

These facts, along with the reduction step, give us the following:

- $\mathbf{selftran}_{\Pi\cup\Lambda, \cdot, A_2, R} w_3, w_4$
- $(\tilde{e}_2|h_2^\dagger \cdot g_2^\dagger \cdot h_{22} \cdot f_2^\dagger) \mapsto^* (e_2''|g_2^f \cdot h_2^f \cdot h_{22}^f \cdot f_2^\dagger)$
- $g_1^\dagger \cdot h_1^\dagger \cdot h_{21} = g_1^f \cdot h_1^f \cdot h_{21}^f$
- $(g_1^f \cdot h_1^f, h_{21}, f_1, g_2^f \cdot h_2^f, h_{22}, f_2^\dagger) \in \mathbf{splits}_{\Pi\cup\Lambda, \cdot, A_2, R} w_4$
- $(j - n, g_1^\dagger \cdot h_1^\dagger, g_2^\dagger \cdot h_2^\dagger, g_1^f \cdot h_1^f, g_2^f \cdot h_2^f) \in \mathbf{Q}_{\varepsilon_2}^{\Pi\cup\Lambda, R} w_3, w_4$
- $(j - n, \emptyset, \emptyset, e_2', e_2'', h_{21}, h_{22}) \in \mathbf{safe}_{\tau_2, \varepsilon_2}^{\Pi\cup\Lambda, \cdot, A_2, R} w^\circ, w_4$

Now we can show the remaining goals:

- $\mathbf{selftran}_{\Pi\cup\Lambda, \cdot, A_2, R} w_2, w_4$, by self-transition composition
- $(\tilde{e}_2|g_2^\dagger \cdot h_2^\dagger \cdot f_2^\dagger) \mapsto^* (e_2''|g_2^f \cdot h_2^f \cdot h_{22}^f \cdot f_2^\dagger)$, by transitivity
- $h_1^\dagger \cdot g_1^\dagger \cdot h_{21} = g_1^f \cdot h_1^f \cdot h_{21}^f$, by assumption
- $(j - n, g_1^\dagger, g_2^\dagger, g_1^f, g_2^f) \in \mathbf{Q}_{\varepsilon_2}^{\Pi, R} w_2, w_4$, by postcondition composition
- $(j - n, h_1^f, h_2^f) \in \mathbf{P}_{\varepsilon_2}^{\Lambda, R} w_4$, by precondition composition
- $(i, h_1^\circ, h_2^\circ, h_1^f, h_2^f) \in \mathbf{Q}_{\varepsilon_1 \cup \varepsilon_2}^{\Lambda, R} w^\circ, w_4$, by postcondition composition

- $(j - n, g_1^\dagger \cdot h_1^\dagger, g_2^\dagger \cdot h_2^\dagger, g_1^f \cdot h_1^f, g_2^f \cdot h_2^f) \in \mathbf{Q}_{\varepsilon_2}^{\Pi\cup\Lambda, R} w_3, w_4$, by assumption
- $(j - n, \emptyset, \emptyset, e_2', e_2'', h_{21}, h_{22}) \in \mathbf{safe}_{\tau_2, \varepsilon_2}^{\Pi\cup\Lambda, \cdot, A_2, R} w^\circ, w_4$, by assumption,

which ends this case.

- 3) $(\text{par } e_1 \text{ and } e_2 | g_1 \cdot h_1 \cdot h_{11} \cdot h_{21} \cdot f_1) \rightarrow_{\emptyset}^1 ((e_1, e_2) | g_1 \cdot h_1 \cdot h_{11} \cdot h_{21} \cdot f_1)$

This means in particular that $\text{irr}(e_i | g_1 \cdot h_1 \cdot h_{11} \cdot h_{21} \cdot f_1)$. We only need to consider the case where $j \geq 1$. We start by using the assumption (2). We need to provide:

- $\mathbf{envtran}_{\Pi\cup\Lambda, \cdot, A_1, R} w, w_1$, by env-transition weakening
- $(j, g_1 \cdot h_1, g_2 \cdot h_2) \in \mathbf{P}_{\varepsilon_1}^{\Pi\cup\Lambda, R} w_1$, by precondition separation, assumption and (1)
- $(g_1 \cdot h_1, h_{11}, f_1 \cdot h_{21}, g_2 \cdot h_2, h_{12}, f_2) \in \mathbf{splits}_{\Pi\cup\Lambda, \cdot, A_1, R} w_1$, by massaging the assumption about **splits**.

From the termination case we can now get:

- $(\tilde{e}_1|g_2 \cdot h_2 \cdot h_{12} \cdot f_2) \mapsto^* (e_1'|g_2' \cdot h_2' \cdot h_{12}' \cdot f_2)$
- $\mathbf{selftran}_{\Pi\cup\Lambda, \cdot, A_2, R} w_1, w_2$
- $g_1 \cdot h_1 \cdot h_{11} = g_1' \cdot h_1' \cdot h_{11}'$
- $(g_1' \cdot h_1', h_{11}', f_1 \cdot h_{21}, g_2' \cdot h_2', h_{12}', f_2) \in \mathbf{splits}_{\Pi\cup\Lambda, \cdot, \emptyset, R} w_2$
- $(j, g_1, g_2, g_1', g_2') \in \mathbf{Q}_{\varepsilon_1}^{\Pi, R} w_1, w_2$
- $(j, h_1, h_2, h_1', h_2') \in \mathbf{Q}_{\varepsilon_1}^{\Lambda, R} w_1, w_2$
- $(j, e_1, e_1') \in \llbracket \tau_1 \rrbracket^R w_2$

Note that \tilde{e}_1' is irreducible, as it's a value. At this point we can use assumption (3). We need to show:

- $\mathbf{envtran}_{\Pi\cup\Lambda, \cdot, A_2, R} w, w_2$, by relation between self-transition and env-transition and env-transition weakening
- $(j, h_1, h_2, h_1', h_2') \in \mathbf{Q}_{\varepsilon_1}^{\Lambda, R} w, w_2$, by assumption
- $(g_1' \cdot h_1', h_{21}, f_1 \cdot h_{11}, g_2' \cdot h_2', \emptyset, f_2 \cdot h_{12}') \in \mathbf{splits}_{\Pi\cup\Lambda, \cdot, A_2, R} w_2$, by **splits** above.

By these facts we are able to obtain the following:

- $\mathbf{selftran}_{\Pi\cup\Lambda, \cdot, A_2, R} w_2, w_3$
- $(\tilde{e}_2|g_2' \cdot h_2' \cdot f_2 \cdot h_{12}') \mapsto^* (e_2''|g_2'' \cdot h_2'' \cdot h_{22} \cdot f_2 \cdot h_{12}')$
- $h_1' \cdot g_1' \cdot h_{21} = h_1'' \cdot g_1'' \cdot h_{21}'$
- $(k, g_1', g_2', g_1'', g_2'') \in \mathbf{Q}_{\varepsilon_2}^{\Pi, R} w_2, w_3$
- $(k, h_1', h_2') \in \mathbf{P}_{\varepsilon_2}^{\Lambda, R} w_3$
- $(j, h_1^\circ, h_2^\circ, h_1'', h_2'') \in \mathbf{Q}_{\varepsilon_1 \cup \varepsilon_2}^{\Lambda, R} w^\circ, w_3$
- $(k, \emptyset, \emptyset, e_2, e_2', h_{21}, h_{22}) \in \mathbf{safe}_{\tau_2, \varepsilon_2}^{\Pi\cup\Lambda, \cdot, A_2, R} w^\circ, w_3$
- $(g_1'' \cdot h_1'', h_{21}, f_1 \cdot h_{11}, g_2'' \cdot h_2'', h_{22}, f_2 \cdot h_{12}') \in \mathbf{splits}_{\Pi\cup\Lambda, \cdot, A_2, R} w_3$.

Finally, we can instantiate the safety predicate obtained above. We need to show:

- $\mathbf{envtran}_{\Pi\cup\Lambda, \cdot, A_2, R} w_3, w_3$, by reflexivity
- $(j, h_1'', h_2'') \in \mathbf{P}_{\varepsilon_2}^{\Lambda, R} w_3$, by assumption
- $(j, g_1'', g_2'') \in \mathbf{P}_{\varepsilon_2}^{\Pi, R} w_3$, trivial, since $\text{rds } \varepsilon_2 \# \Pi$
- $(g_1'' \cdot h_1'', h_{21}, f_1 \cdot h_{11}, g_2'' \cdot h_2'', h_{22}, f_2 \cdot h_{12}') \in \mathbf{splits}_{\Pi\cup\Lambda, \cdot, A_2, R} w_3$, by assumption.

From the termination branch (by irreducibility of e_2), we get:

- $\text{selftran}^{\Pi \cup \Lambda, \cdot, A_2, R} w_3, w_4$
- $(\tilde{e}'_2 | g''_2 \cdot h''_2 \cdot h_{22} \cdot f_2 \cdot h'_{12}) \mapsto^* (\tilde{e}'_2 | g'''_2 \cdot h'''_2 \cdot h'''_2 \cdot f_2 \cdot h'_{12})$
- $g''_1 \cdot h''_1 \cdot h_{21} = g'''_1 \cdot h'''_1 \cdot h'''_{21}$
- $(j, g''_1, g''_2, g'''_1, g'''_2) \in \mathbf{Q}_{\varepsilon_2}^{\Pi, R} w_3, w_4$
- $(j, h''_1, h''_2, h'''_1, h'''_2) \in \mathbf{Q}_{\varepsilon_2}^{\Lambda, R} w_3, w_4$
- $(j, e_2, \tilde{e}'_2) \in \llbracket \tau_2 \rrbracket^R w_4$
- $(g'''_1 \cdot h'''_1, h'''_{21}, f_1 \cdot h'_{11}, g'''_2 \cdot h'''_2, h'''_{22}, f_2 \cdot h'_{12}) \in \text{splits}^{\Pi \cup \Lambda, \cdot, \emptyset, R} w_4$

At this point we can finally provide witnesses and prove the remaining properties:

- $\text{selftran}^{\Pi, \Lambda, A_1 + A_2, R} w_1, w_4$, by weakening and composition of previous self-transitions
- $(\tilde{e}'_1, \tilde{e}'_2) | g_2 \cdot h_2 \cdot h_{12} \cdot f_2 \mapsto^* (\tilde{e}'_1, \tilde{e}'_2) | g'''_2 \cdot h'''_2 \cdot h'''_{22} \cdot h'_{12} \cdot f_2$, by transitivity and irreducibility of \tilde{e}'_1
- $\emptyset \in \text{effs}_{\varepsilon_1 \cup \varepsilon_2, h'''_1 \cdot h'_{11}, h'''_{21}} w_4$, which is trivially true
- $g_1 \cdot h_1 \cdot h_{11} \cdot h_{21} = g'''_1 \cdot h'''_1 \cdot h'_{11} \cdot h'_{21}$, by congruence closure of equalities
- $(g'''_1 \cdot h'''_1, h'''_{21}, f_1, g'''_2, h'''_2 \cdot h'_{22}, h'_{12}, f_2) \in \text{splits}^{\Pi, \Lambda, \emptyset, R} w_4$, by massaging the assumption above
- $(j - 1, g_1, g_2, g'''_1, g'''_2) \in \mathbf{Q}_{\varepsilon_1 \cup \varepsilon_2}^{\Pi, R} w_1, w_4$, by weakening and composition of postcondition assumptions

This leaves us with the final obligation to prove: $(j - 1, h_1^\circ, h_2^\circ, \langle e_1, e_2 \rangle, \langle \tilde{e}'_1, \tilde{e}'_2 \rangle, h'''_1 \cdot h'_{11} \cdot h'_{21}, h'''_2 \cdot h'_{12} \cdot h'_{22}) \in \text{safe}_{\tau_1 \times \tau_2, \varepsilon_1 \cup \varepsilon_2}^{\Pi, \Lambda, \emptyset, R} w^\circ, w_4$. To do this, we assume:

- $\text{envtran}^{\Pi, \Lambda, \emptyset, R} w_4, w_5$
- $(i, g_1^\dagger, g_2^\dagger) \in \mathbf{P}_{\varepsilon_1 \cup \varepsilon_2}^{\Pi, R} w_5$
- $(g_1^\dagger, h'''_1 \cdot h'_{11} \cdot h'_{21}, h'''_2, f_1^\dagger, g_2^\dagger, h'''_2 \cdot h'_{12} \cdot h'_{22}, f_2^\dagger) \in \text{splits}^{\Pi, \Lambda, \emptyset, R} w_5$

Obviously, both expressions are irreducible, so we take the same values as witnesses. Most obligations are proved by reflexivity or assumption, the ones left are:

- $(i, \langle e_1, e_2 \rangle, \langle \tilde{e}'_1, \tilde{e}'_2 \rangle) \in \llbracket \tau_1 \times \tau_2 \rrbracket^R w_5$, which holds by downwards- and future-world-closure of type denotation, and definition of product types, and
- $(i, h_1^\circ, h_2^\circ, h'''_1, h'''_2) \in \mathbf{Q}_{\varepsilon_1 \cup \varepsilon_2}^{\Lambda, R} w^\circ, w_5$, which holds by weakening, composition and downwards closure of previous postcondition assumptions.

This ends the proof. ■

Lemma 18 (Parallelization). *Assuming that*

- 1) $\Pi, \Lambda \parallel \Gamma \vdash e_1 : \tau_1, \varepsilon_1$,
- 2) $\Pi, \Lambda \parallel \Gamma \vdash e_2 : \tau_2, \varepsilon_2$,
- 3) $\text{rds } \varepsilon_1 \cup \text{wrs } \varepsilon_1 \cup \text{rds } \varepsilon_2 \cup \text{wrs } \varepsilon_2 \subseteq \Lambda$,
- 4) $\text{rds } \varepsilon_1 \cap \text{wrs } \varepsilon_2 = \text{rds } \varepsilon_2 \cap (\text{wrs } \varepsilon_1 \cup \text{als } \varepsilon_1) = \text{wrs } \varepsilon_1 \cap \text{wrs } \varepsilon_2 = \emptyset$,

the following property holds:

$$\Pi \parallel \Lambda \parallel \Gamma \models \langle e_1, e_2 \rangle \cong \text{par } e_1 \text{ and } e_2 : \tau_1 \times \tau_2, \varepsilon_1 \cup \varepsilon_2.$$

Proof: The proof consists of two parts, for two directions of contextual approximation. For both these direction we need

assumptions about e_1 and e_2 , that follow by the fundamental theorem of logical relations:

$$\Pi, \Lambda \parallel \Gamma \models e_1 \preceq e_1 : \tau_1, \varepsilon_1 \quad \Pi, \Lambda \parallel \Gamma \models e_2 \preceq e_2 : \tau_2, \varepsilon_2$$

- 1) To show: $\Pi \parallel \Lambda \parallel \Gamma \models \langle e_1, e_2 \rangle \preceq \text{par } e_1 \text{ and } e_2 : \tau_1 \times \tau_2, \varepsilon_1 \cup \varepsilon_2$.

We start this part, by taking a_1 and a_2 , the witnesses from the logical relations above, and setting $a = a_1 + a_2$ as the witness for the logical relation. Now we can assume the initial conditions:

- $R(\text{FRV}(\varepsilon_1 \cup \varepsilon_2)) \subseteq \text{dom}(w)$
- $A_1 + A_2 \# R(\Pi \cup \Lambda)$
- $|A_i| \geq a_i$, for $i \in \{1, 2\}$
- $\forall r \in A_1 + A_2. w(r) = \emptyset$
- $(k, \gamma_1, \gamma_2) \in \llbracket \Gamma \rrbracket^R w$
- $h_i^\circ \subseteq h_i$, for $i \in \{1, 2\}$
- $(k, h_1^\circ, h_2^\circ) \in \mathbf{P}_{\varepsilon_1 \cup \varepsilon_2}^{\Lambda, R} w$.

What remains to be shown is safety: $(k, h_1^\circ, h_2^\circ, \langle e_1^1, e_1^2 \rangle, \text{par } e_1^2 \text{ and } e_2^2, h_1, h_2) \in \text{safe}_{\tau_1 \times \tau_2, \varepsilon_1 \cup \varepsilon_2}^{\Pi, \Lambda, A_1 + A_2, R} w, w$, where $e_i^j = e_i[\gamma_j/\Gamma]$. We can also use the initial conditions above to obtain:

$$(k, \emptyset, \emptyset, e_1^1, e_2^2, \emptyset, \emptyset) \in \text{safe}_{\tau_i, \varepsilon_i}^{\Pi \cup \Lambda, \cdot, A_i, R} w, w.$$

We proceed by well-founded induction, keeping the additional assumption that $(k, h_1 \upharpoonright_{\text{dom}_1^{R(\Lambda)}(w)}, h_2 \upharpoonright_{\text{dom}_2^{R(\Lambda)}(w)}) \in \mathbf{P}_{\varepsilon_1 \cup \varepsilon_2}^{\Lambda, R} w$. After unrolling safety, there are three possibilities. Either e_i^1 are both irreducible, in which case we can recover the final reductions for e_i^2 , replay those, and reduce the parallel composition to the pair of values related to the left-hand-side pair, or at least one of e_i^1 can reduce. In the latter cases, we can use the respective assumption to obtain the matching reduction, perform it, and use the induction hypothesis (with downwards- and envtran-closure ensuring that the other safety assumption matches and the precondition assumption fulfilled due to Λ being public). The details are very similar to the compatibility of parallel composition, and so are omitted.

- 2) To show: $\Pi \parallel \Lambda \parallel \Gamma \models \text{par } e_1 \text{ and } e_2 \preceq \langle e_1, e_2 \rangle : \tau_1 \times \tau_2, \varepsilon_1 \cup \varepsilon_2$.

We proceed in the similar manner, by taking the witnesses, providing $a_1 + a_2$ as the witness for the whole computation, and obtaining the safety obligations: show that $(k, h_1^\circ, h_2^\circ, \text{par } e_1^1 \text{ and } e_2^1, \langle e_1^2, e_2^2 \rangle, h_1, h_2) \in \text{safe}_{\tau_1 \times \tau_2, \varepsilon_1 \cup \varepsilon_2}^{\Pi, \Lambda, A_1 + A_2, R} w, w$ provided that $(k, \emptyset, \emptyset, e_i^1, e_i^2, \emptyset, \emptyset) \in \text{safe}_{\tau_i, \varepsilon_i}^{\Pi \cup \Lambda, \cdot, A_i, R} w, w$, where $e_i^j = e_i[\gamma_j/\Gamma]$. However, to show it we use the catch-up lemma proved above. To do this, we need to show its three preconditions (we let $h_1^l = h_1 \upharpoonright_{\text{dom}_1^{R(\Lambda)}(w)}$, $h_{11} = h_1 \setminus h_1^l$, $h_{12} = \emptyset$, and similar for h_2):

- $(k, h_1^l, h_2^l) \in \mathbf{P}_{\varepsilon_1}^{\Lambda, R} w$, which holds by preweakening,
- $(k, \emptyset, \emptyset, e_1^1, e_2^2, h_{11}, h_{12}) \in \text{safe}_{\tau_1, \varepsilon_1}^{\Pi \cup \Lambda, \cdot, A_1, R} w, w$, which is easy to show since h_{1i} is just some random

heap outside the control of the world (regions in A_1 are empty),

- and the final catching-up precondition. To show this property, take any $g_1, g_2, h_1^\dagger, h_2^\dagger, f_1, f_2, w', j \leq k$ such that:

- $\text{envtran}^{\Pi \cup \Lambda, \cdot, A_2, R} w, w'$
- $(g_1 \cdot h_1^\dagger, \emptyset, f_1, g_2 \cdot h_2^\dagger, \emptyset, f_2) \in \text{splits}^{\Pi \cup \Lambda, \cdot, A_2, R} w'$
- $(j, h_1, h_2, h_1^\dagger, h_2^\dagger) \in \mathbf{Q}_{\varepsilon_1}^{\Lambda, R} w, w'$

We take the witnesses to be the precise same things, so most of the obligations hold by reflexivity. The ones left are:

- $(j, h_1^\circ, h_2^\circ, h_1^\dagger, h_2^\dagger) \in \mathbf{Q}_{\varepsilon_1 \cup \varepsilon_2}^{\Lambda, R} w, w'$, which by post-weaken, since $h_i^\circ = h_i^\dagger$,
- $(k, \emptyset, \emptyset, e_2^1, e_2^2, \emptyset, \emptyset) \in \text{safe}_{\tau_2, \varepsilon_2}^{\Pi \cup \Lambda, \cdot, A_2, R} w, w'$, which holds by envtran-closure of the safety assumption, and
- $(k, h_1^\dagger, h_2^\dagger) \in \mathbf{P}_{\varepsilon_2}^{\Lambda, R} w'$, which holds because of disjointness condition on ε_1 and ε_2 — when cut down to the ε_2 -read regions, $h_i^\dagger = h_i^\circ$, for which we have the right assumption.

This ends the proof. ■