# Computation expressions and monads 

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## Agenda

- Computation expressions, or monads
- Sequence expressions as computation exprs.
- Monad laws
- Simple expressions and evaluators
- Return int, standard evaluation
- Return int option, evaluation may fail
- Return int Set, evaluation may produce a set
- Return int trace, evaluation traces operators
- Express all this uniformly using comp. expr.
- Next week: async as computation expression


## What is a computation expression?

- A computation expression such as seq $\left\{\right.$ for $\mathbf{x}$ in [1 .. 3] do yield $\mathbf{x}^{*} \times$ \}
- is syntactic sugar for a standard functional expression, such as

Seq.collect (fun $x$-> Seq.singleton( $x^{*}$ x)) [1 .. 3]

- This gives a systematic way to combine operations or propagate "background" data
- Computation expressions are sometimes called monads (in mathematics, Haskell, ...)


## Computation expressions in F\#

- seq $\{\ldots\}$ is a computation expression
- async $\{\ldots\}$ is a comp. expr. (next week)
- We can define our own computation exprs:
- Optional result (None or Some)
- Set of results
- Trace of operations


## Sequences as computation expressions (H\&R p 281)

- This sequence expression

```
seq { for i in [1 .. 3] do
                for ch in ['a' .. 'd'] do
    yield (i,ch) }
```

- is syntactic sugar for this expression

Seq.collect
(fun i ->
Seq.collect (fun ch ->

Seq. singleton (i, ch))
['a' .. 'd'])
[1 .. 3]

## Transformation of seq \{...\}

The compiler rewrites for and yield keywords to normal function calls:

| Seq construct C | Transformation T(C) |
| :--- | :--- |
| for $\mathbf{x}$ in $e$ do ce | For (e, fun $\mathbf{x ~}->T(c e))$ |
| yield $e$ | Yield(e) |

The For and Yield functions must be defined:
For : seq<'a> * ('a $->$ seq<'b>) $->$ seq<'b> Yield : 'a -> seq<'a>

```
For(xs, f) = Seq.collect f xs
Yield a = Seq.singleton a
```


## How to define our own mySeq \{...\}?

- Define a MySeqBuilder class with For, Yield: type MySeqBuilder() = member this.For (xs, f) $=$ Seq.collect $f$ xs member this.Yield $x=$ Seq. singleton $x$
- Make an object of that class:
let mySeq $=$ new MySeqBuilder ()
- The object can now indicate a comp expr:
mySeq \{ for i in [1 .. 3] do
for ch in ['a' .. 'd'] do
yield (i, ch) \};;

Homemade
sequence
expression

## Understanding the H\&R example

```
seq { for i in [1 .. 3] do
    for ch in ['a' .. 'd'] do
    yield (i,ch) }
```

- Lift out inner for as a function:

```
let inner i =
    seq { for ch in ['a' .. 'd'] do yield (i,ch) }
```

let inner i =
Seq. collect (fun ch $\rightarrow$ Seq.singleton (i, ch)) ['a'.. 'd']

- Outer for is just this:

```
Seq.collect (fun i -> inner i) [1 .. 3]
```

- So in total

Seq. collect (fun i ->
Seq. collect (fun ch -> Seq.singleton (i, ch)) ['a' .. 'd']) [1 .. 3]

## List and array expressions

- The F\# list expression:
[ for $\mathbf{x}$ in [1..3] do yield $\mathbf{x * x}^{*}$ ]
- is syntactic sugar for the seq expression

Seq.toList(seq \{for $x$ in [1..3] do yield $x * x$ \})

- Similarly for F\# array expressions:
[। for $\mathbf{x}$ in [1..3] do yield $\mathbf{x}^{*}$ (]
- See F\# Specification §6.3.13 and §6.3.14
- F\# has no "list computation expression", it boils down to seq computation expressions


## Questions

- In seq \{...\} expressions one can use "if"

```
let sift a xs =
seq { for n in xs do
    if n % a <> O then
    yield n };;
```

- Q1: What function to add to MySeqBuilder to support the if operator? (H\&R Table 12.2)
- Q2: What member of the Seq module should be used to define it? (H\&R Table 11.1)


## For=Bind=let!, Yield=Return

- For and Yield are special names that make sense in seq\{...\} expressions
- Normal names are Bind/let! and Return
- One could define sillySeq using Bind/Return:

| ```type SillySeqBuilder() = member this.Bind(xs, f) = Seq.collect f xs member this.Return x = Seq.singleton x let sillySeq = new SillySeqBuilder()``` |  |
| :---: | :---: |
|  |  |
| ```sillySeq { let! i = [1 .. 3] let! ch = ['a' .. 'd'] return (i, ch) }``` | ```seq { for i in [1 .. 3] do for ch in ['a' .. 'd'] do yield (i,ch) }``` |

## Kært barn har mange navne

- Function Bind in computation expression is
- for and For in seq \{...\}
- List.collect on F\# lists
- Seq. collect on F\# sequences
- flatMap on Scala and Haskell lists, sequences, ...
- SelectMany in Microsoft Linq (eg. C\#P p. 205)
- bind in monads
- Function Return in computation expressions
- yield in seq \{...\}
- (fun x -> [x]) on F\# lists
- Seq. singleton on F\# sequences
- unit or return in monads


## The option \{...\} computation expr. A form of error propagation

```
optM { let x = 56
    let! y = Some 78
        return x+y };
optM { let x = 56
    let! y = Some 78
    let! z = None
    return x+y };;
```

let $\mathrm{x}=56$
optM. Bind (Some 78, fun $y$->
let $x=56$
optM. Bind(Some 78,
fun $y$->
optM. Bind (None, fun $z$-> optM.Return( $x+y$ )))
type OptionBuilder() = member this. $\operatorname{Bind}(x, f)=$ let! match $x$ with
| None -> None
| Some v -> f v
member this.Return (x) $=$ Some $x$
return

## Question: What's happening here

```
type OptionBuilder() =
    member this.Bind(x, f) =
    printfn "this.Bind: %A" x
        match x with
        | None -> None
        | Some v -> f v
    member this.Return(x) = Some x
let optM = OptionBuilder()
optM { let x = 56
    let! y = Some 78
    let! z = None
    let! v = Some 42
    return x+y+v };;
``` here

\section*{Monad laws}
- For(Yield \(a, f)=f(a)\) collect (singleton a) \(f=f(a)\)
- For(xs, Yield) = xs collect xs singleton \(=x s\)
- \(\operatorname{For}(\operatorname{For}(x s, f), g)=\operatorname{For}(x s, f u n x->g(f(x)))\) collect (collect xs f) \(g=\) collect xs (fun \(x->g(f(x))\) )
- The laws are the same for
- Bind instead of For, and Return instead of Yield
- Let's check them for the maybe monad

\section*{A standard simple evaluator}
- Very simple expressions like \(7+9\) * 10
type expr =
```

| CstI of int
| Prim of string * expr * expr

```

Prim("+",
CstI (7) ,
Prim("*", CstI (9), CstI (10)))
- A simple evaluator:
let rec evalle : int = match e with
| CstI i -> i
| Prim(op, e1, e2) -> let v 1 = evall e1 let v 2 = evall e2 match op with | "+" -> v1 + v2 | "*" -> v1 * v2 | "/" -> v1 / v2
let opEval op v1 v2 : int = match op with
| "+" -> v1 + v2
| "*" -> v1 * v2
| "/" -> v1 / v2
let rec eval2 e : int \(=\) match e with
| CstI i -> i
Prim(op, e1, e2) -> let v1 = eval2 e1 let v2 = eval2 e2 opEval op v1 v2

\section*{An evaluator that may fail (w. None)}
let opEvalopt op v1 v2 : int option = match op with
| "+" -> Some(v1 + v2)
| "*" -> Some(v1 * v2)
| "/" -> if v2 = 0 then None else Some(v1 / v2)
let rec optionEval2 e: int option = match e with
| CstI i -> Some i
| Prim(op, e1, e2) ->
match optionEval2 e1 with
| None -> None
| Some v1 ->
match optionEval2 e2 with
| None -> None
| Some v2 -> opEvalopt op v1 v2

\section*{An evaluator giving a set of results}
let opEvalSet op v1 v2 : int Set =
match op with
| "+" -> Set [v1 + v2]
| "*" -> Set [v1 * v2]
| "/" -> if v2 = 0 then Set.empty else Set [v1 / v2]
| "choose" -> Set [v1; v2]
```

let rec setEvall e : int Set=
match e with
CstI i -> Set [i]
| Prim(op, e1, e2) ->
let s1 = setEval1 e1
let yss =
Set.map (fun v1 ->

```
        let \(s 2\) = setEval1 e2
        let xss \(=\) Set.map (fun v2 -> opEvalSet op v1 v2) s2
        Set.unionMany xss)
        s1
    Set.unionMany yss

\section*{An evaluator tracing the operators}
```

type 'a trace = string list * 'a
let opEvalTrace op v1 v2 : int trace =
match op with
| "+" -> (["+"], v1 + v2)
| "*" -> (["*"], v1 * v2)
| "/" -> (["/"], v1 / v2)
let rec traceEvall e: int trace=
match e with
| CstI i -> ([], i)
| Prim(op, e1, e2) ->
let (trace1, v1) = traceEval1 e1
let (trace2, v2) = traceEval1 e2
let (trace3, res) = opEvalTrace op v1 v2
(trace1 @ trace2 @ trace3, res)

```

\section*{A mess; comp expr to the rescue}
- The four evaluators look very different
- ... and very complicated
- By defining the combining operations as computation expressions,
- the evaluators all get to look the same
- the evaluators look much simpler

\section*{An evaluator that may fail, NEW}
```

type OptionBuilder() =
member this.Bind(x, f) =
match x with
| None -> None
| Some v -> f v
member this.Return x = Some x
member this.ReturnFrom x = x
let optionM = OptionBuilder();;
let rec optionEval3 e : int option =
match e with
| CstI i -> optionM { return i }
| Prim(op, e1, e2) ->
optionM { let! v1 = optionEval3 e1
let! v2 = optionEval3 e2
return! opEvalOpt op v1 v2 }

```

\section*{An evaluator ... set of results, NEW}
type SetBuilder () =
member this.Bind (x, f) =
Set. unionMany (Set.map \(f \mathbf{x}\) )
member this.Return \(x=\) Set [x]
member this.ReturnFrom \(\mathbf{x}=\mathbf{x}\)
let setM = SetBuilder(); ;
let rec setEval3 e : int Set = match e with
| CstI i -> setM \{ return i \}
| Prim(op, e1, e2) ->
setM \{ let! v1 = setEval3 e1
let! v2 = setEval3 e2
return! opEvalSet op v1 v2 \}

\section*{An evaluator ... trace operators, NEW}
```

type TraceBuilder() =
member this.Bind(x, f) =
let (trace1, v) = x
let (trace2, res) = f v
(trace1 @ trace2, res)
member this.Return x = ([], x)
member this.ReturnFrom x = x

```
let traceM = TraceBuilder(); ;
let rec traceEval3e : int trace \(=\)
    match e with
    | CstI i -> traceM \{ return i \}
    | Prim(op, e1, e2) ->
        traceM \{ let! v1 = traceEval3 e1
        let! v2 = traceEval3 e2
        return! opEvalTrace op v1 v2 \}

\section*{A standard evaluator, NEW}
type IdentityBuilder() = member this.Bind ( \(x, f\) ) \(=f \mathbf{x}\)
member this.Return \(\mathbf{x}=\mathbf{x}\)
member this.ReturnFrom \(\mathbf{x}=\mathbf{x}\)
let identM \(=\) new IdentityBuilder(); ;
let rec eval3 e : int = match e with
| CstI i -> identM \{ return i \}
| Prim(op, e1, e2) ->
identM \{ let! v1 = eval3 e1
let! v2 = eval3 e2
return! opEval op v1 v2 \}

\section*{Reflections on}

\section*{computation expressions}
- They reveal similarities
- between different kinds of computations
- between different kinds of data: list, seq, option, ...
- They clarify the structure of the evaluators
- Unfortunately, in F\# a computation expression builder (optionM, setM, traceM, identM) cannot be a parameter to a function
- Hence one cannot have a single "super-eval"
- ... but in Scala we can, you'll see in November

\section*{References}
- F\# computation expressions
- Hansen and Rischel chapter 12
- http://en.wikibooks.org/wiki/F_Sharp_Programming/ Computation_Expressions
- http://msdn.microsoft.com/en-us/library/dd233182.aspx```

