This exam consists of 5 problems with a total of 15 questions. The weight of each problem is stated. You have 4 hours to answer all questions. The complete assignment consists of 6 numbered pages (including this page).

If you cannot give a complete answer to a question, try to give a partial answer. You may choose to write your answer in Danish or English. Write only on the front of sheets, and remember to write your CPR-number on each page. Please start your answer to each question at the top of a new page. Please order and number the pages before handing in.


All written aids are allowed / Alle skriftlige hjælpemidler er tilladt.
1 SQL DDL and normalization (20%)

The national railroad company in the Republic of Delalaya maintains a relational database of information on (physical) trains, train personnel, and manning. The database stores key numbers for each train, and information on which employees work on a given train departure. To simplify administrative procedures, employees always work on the same physical train, and always work all the way from the departure station to the final destination. The database has the following schema:

\[
\text{Train}(\text{Tid}, \text{type}, \text{productionYear}, \text{capacity}) \\
\text{Personnel}(\text{Pid}, \text{worksOnTrain}, \text{worksAs}, \text{hiredDate}, \text{salary}) \\
\text{Manning}(\text{Tid}, \text{Pid}, \text{Sid}, \text{onDate})
\]

The attribute \text{type} of \text{Train} refers to the train manufacturer’s code for a specific kind of trains. All trains with the same code are identical. The attribute \text{worksOnTrain} of \text{Personnel} is a foreign key reference to \text{Train}(\text{Tid}). This reflects the fact that each person always works on the same train (but of course not all the time). Any tuple in \text{Personnel} must contain a valid reference to a tuple in \text{Train}. The \text{Sid} attribute of \text{Manning} is a reference to a code for a specific departure in the time table. However, the time table is not part of the database. The attribute pair (\text{Pid}, \text{Tid}) in \text{Manning} is a foreign key reference to \text{Personnel}(\text{Pid}, \text{worksOnTrain}).

a) Write SQL statements that create the above relations, including key and foreign key constraints. You can make any reasonable assumption about the data types.

Answer:

CREATE TABLE Train (  
  Tid INT PRIMARY KEY,  
  type INT,  
  productionYear INT,  
  capacity INT  
);

CREATE TABLE Personnel (  
  Pid INT PRIMARY KEY,  
  worksOnTrain INT NOT NULL REFERENCES Train(Tid),  
  worksAs VARCHAR(20),  
  hiredDate DATE,  
  salary INT,  
  UNIQUE (Pid, worksOnTrain)  
);

CREATE TABLE Manning (  
  Tid INT,  
  Pid INT,  
  Sid INT,  
  onDate DATE
);
The unique constraint on (Pid, worksOnTrain) is needed at least in some DBMSs, even though it is implied by the primary key constraint. The "NOT NULL" constraint on worksOnTrain is needed to ensure referential integrity, as required.

b) Identify all functional dependencies in the relations that do not have a superkey on the left hand side (i.e., are “avoidable”). Use the dependencies to decompose the schema into BCNF, and state the resulting database schema. Briefly discuss the quality of the new schema (ignoring efficiency issues).

Answer:

We have the following avoidable FDs:

- type → capacity
- Pid → Tid

Decomposition according to these leads to the schema:

- Train(Tid, type, productionYear)
- TrainType(type, capacity)
- Personnel(Pid, worksOnTrain, worksAs, hiredDate, salary)
- Manning(Pid, Sid, onDate)
- Manning2(Pid, Tid)

We can notice that Manning2 is redundant because it is equal to the projection of Personnel onto Pid, worksOnTrain.

2 Data modeling (30%)

The transportation authority in the Republic of Delalaya has decided to implement a database to keep statistics on public transportation (fuel-driven buses and electricity-driven trains), with emphasis on keeping track of delays and the number of travelers. In particular, it should be used to identify particular weaknesses in the transportation systems (e.g., stretches of railroad that often cause delays, or employees who have trouble keeping the schedule).

First of all, data from the railroad company, as described in Problem 1, should be integrated into the system, but not necessarily using the schema stated there.
a) Draw an ER diagram corresponding to the data described in Problem 1, including, if possible, all integrity constraints stated there. You should follow general rules for making a good ER design — your design does not need to translate to the three relations stated in Problem 1.

Answer:

In addition to the above, the following information is needed:

- **Information on buses.** Similar to the information on trains, but in addition the *range* of each bus (the number of kilometers it will drive on a full tank) should be recorded.

- **Information on bus drivers.** Similar to the information on train personnel, with the following changes: All bus drivers have the same work (driving the bus). A bus driver does not drive a specific bus, but may drive any bus.

- **Route information.** The sequence of stops on each train/bus route. Each route has a unique *route number*. 
• **Vehicle usage.** For each route and departure time, on each day, record which physical train/bus was used. (In Delalaya, this is always a single bus or train – several trains can’t be coupled together.)

• **Timetable information.** Information on *planned* arrival and departure times for each route and every stop.

• **Timetable statistics.** Information on *actual* arrival and departure times for every train/bus and every stop, on every day.

• **Traveler statistics.** Periodically, surveys are being made that record the destinations of all travelers in a given bus/train at a given time (between two stops).

• **Manning.** Who has worked on a particular vehicle at every time. It should be taken into account that manning may change at any time during a route.

b) Draw an ER diagram that integrates the additional data described above with the ER diagram of Problem 2 a). You should follow general rules for making a good ER design. If you need to make any assumptions, state them.

Answer:

Many different designs are possible, with different trade-offs between various aspects. We give here one example of a design. It introduces the concept of a “DriveLog” which is implicit in the description above. It makes the assumption that there are not two stops on a route with the same name (or that we distinguish between going to a given stop first and second time).
c) Suppose that we desire the database to evolve over time (e.g. with new time tables), but we also want to be able to store and query historical data. Outline how the ER diagram of Problem 2 b) could be changed to achieve this.

Answer:

We could use the tuple versioning approach to support temporal “transaction time” data. (Bi-temporal data is not needed in this application.) This would introduce two new attributes of every entity set and many-many relationship. However, for those entities and relationships where we will never want to delete a tuple we need only to record the time when the tuple was inserted. In fact, this information (or information with a similar effect such as onDate) is already present in some entities and could be used instead.
3 SQL (25 %)

Consider a database of flight departures and airplanes, with the following schema:

```
Departure(departureID, airplaneID, destination, departureTime, bookedSeats)
Airplane(airplaneID, modelID, fabricationYear)
Model(ModelID, name, capacity)
```

a) Write an SQL query that lists the \texttt{airplaneID} of all airplanes made before 1960.

Answer:

```
SELECT airplaneID FROM Airplane WHERE fabricationYear < 1960;
```

b) Write an SQL query that lists the \texttt{departureID} for all departures bounded for destinations starting with the letter “D”.

Answer:

```
SELECT departureID FROM Departure where destination LIKE 'D%';
```

c) Write an SQL query that lists the average \texttt{capacity} of airplanes fabricated 1970 or later.

Answer:

```
SELECT AVG(capacity) FROM Airplane NATURAL JOIN Model WHERE fabricationYear >= 1970
```

d) Write an SQL query that lists the \texttt{departureID} for every overbooked departure (i.e. where the number of bookings exceed the capacity of the plane).

Answer:

```
SELECT departureID FROM Departure NATURAL JOIN Airplane NATURAL JOIN Model
WHERE bookedSeats > capacity;
```

e) Write a query in \textbf{relational algebra} that lists the model-\texttt{name} of every airplane that was fabricated in 1970.

Answer:

```
\pi_{name}(\sigma_{fabricationYear=1970}(Airplane \bowtie Model))
```

f) Write an SQL query that lists the \texttt{fabricationYear} of the oldest and second oldest plane. You may assume that the two oldest planes have different values on the attribute \texttt{fabricationYear}.

(Note: Some versions of SQL have a special syntax for this kind of query – however, you are required to write the query using only SQL features found in RG.)

Answer:
SELECT MIN(fabricationYear) FROM Airplane UNION SELECT MIN(fabricationYear)
FROM (SELECT fabricationYear from Airplane EXCEPT SELECT MIN(fabricationYear)
FROM Airplane);

g) Write an SQL query that lists all destinations that has more empty seats than the
average number of empty seat on all departures (we assume that the number of empty
seats is the number of booked seats subtracted from the capacity of the plane).

Answer:

SELECT destination FROM

(SELECT destination, AVG(capacity-bookedSeats) as emptySeats
FROM departure NATURAL JOIN Airplane NATURAL JOIN Model group by destination
)
WHERE emptySeats >

(SELECT AVG(capacity-bookedSeats) FROM Departure NATURAL JOIN Airplane NATURAL
JOIN Model
);

4 Transactions (10 %)

a) Consider an initially empty table T with the schema T(number) and transactions
running at isolation level READ COMMITTED

<table>
<thead>
<tr>
<th>Transaction 1</th>
<th>Transaction 2</th>
</tr>
</thead>
<tbody>
<tr>
<td>1</td>
<td>INSERT INTO t VALUES (1)</td>
</tr>
<tr>
<td>2</td>
<td>INSERT INTO t VALUES (2)</td>
</tr>
<tr>
<td>3</td>
<td>SELECT * FROM T</td>
</tr>
<tr>
<td>4</td>
<td>SELECT * FROM T</td>
</tr>
<tr>
<td>5</td>
<td>SELECT * FROM T</td>
</tr>
<tr>
<td>6</td>
<td>SELECT * FROM T</td>
</tr>
<tr>
<td>7</td>
<td>ROLLBACK</td>
</tr>
</tbody>
</table>

State what is returned from each of the SELECT queries at line 3, 4, 6, and 7.
If there are several possibilities, you may state any of them.

Answer:

SELECT at line 3: 1

SELECT at line 4: 2
b) Consider the two schedules below.

<table>
<thead>
<tr>
<th>Schedule 1:</th>
<th>Schedule 2:</th>
</tr>
</thead>
<tbody>
<tr>
<td><strong>Transaction 1</strong></td>
<td><strong>Transaction 1</strong></td>
</tr>
<tr>
<td>R(A)</td>
<td>R(B)</td>
</tr>
<tr>
<td>R(B)</td>
<td>R(A)</td>
</tr>
<tr>
<td>rollback</td>
<td>W(B)</td>
</tr>
<tr>
<td></td>
<td>W(A)</td>
</tr>
<tr>
<td></td>
<td>commit</td>
</tr>
</tbody>
</table>

State for each of the two schedules whether it is serializable or not. If the schedule is serializable write a serialization of the schedule, otherwise give a brief explanation of why the schedule is not serializable.

Answer:

- **The first schedule is not serializable:** Since transaction 1 reads B and transaction 2 writes to B before the end of transaction 1.

- **The second schedule is serializable:** A valid serialization of the schedule is: 1:R(B), 1:R(A), 1:W(A), 1:commit, 2:R(B), 2:W(B), 2:commit
5 Database efficiency (15%)

Consider the relation $T(id, name)$ with 100,000 tuples. Values of id are positive integers, and values of name are strings of length at most 30. The following queries are of interest:

1. $\text{SELECT } * \text{ FROM } t \text{ WHERE } id = 100$
2. $\text{SELECT } * \text{ FROM } t \text{ WHERE } id > 10$
3. $\text{SELECT } * \text{ FROM } T \text{ WHERE } id > 100 \text{ and } id < 9000 \text{ and } name = 'Mads'$

a) State for each of the above queries whether an index would speed it up. In the cases where the answer is “yes” you also have to specify:
- Would a Hash-index or a B-tree index be the fastest index?
- Which attribute(s) should be indexed?
- Would a clustered index make the query significantly faster than an unclustered index?

Answer:

1. A hash-index on id would speed up the query (and be faster than a B-tree index). A clustered index will not perform better than an unclusted index.

2. The returns almost all tuples (since id is a key). A full table scan is likely to be fastest.

3. A B-tree-index on id, name would speed up the query. A clustered index will perform better than an unclusted index.