# Data storage
# Tree indexes

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# Access paths

- For many database queries and updates, only a small fraction of the data needs to be accessed.
- Extreme examples are looking or updating the single tuple with a given key value.
- General question: *”How do we access the relevant data, and not (much) more?”*
- Term.: Need efficient **access paths**.

# Choices influencing access paths

- How are the tuples in the relations arranged? *(this lecture)*
- What kinds of indexes exist? *(this lecture and later in course)*
- Have we pre-computed (parts of) the information needed? *(later in course)*

# Data storage in sequential files

- Storing meta-data: See textbook.
- A *relation* may be stored as a string, concatenating the tuples:
  (3,"Kurt",27),(11,"Jim",27),(89,"Elvis",42)
- The tuples may be sorted according to some attribute(s), or unordered.
- *Reading* data is very efficient.
- Problem:
  - How to accommodate changes to data: Insertions, deletions, modifications.

# From your toolbox: Linked lists

- Linked (or doubly linked) lists give great flexibility:
  - Insert a new tuple
  - Delete a tuple
  - Change the size of the data in a tuple
- Problem:
  - One ”random” memory access per item when reading the list.

# Memory access cost

- Recall from last week:
  - The speed of the CPU is 2-3 orders of magnitude larger than the speed of a RAM access (non-sequential).
  - The speed of RAM access is 5-6 orders of magnitude faster than disk access (non-sequential).
- RAM vs disk analogy: Go to Australia to borrow a cup of sugar if your neighbor is not home!

# Analysis and comparison

- *I/O model*, i.e., we count the number of block reads/writes.
- For now, assume that tuples have fixed length. Let B denote the number of tuples per block, and N total #tuples.

| *Storage method* | *Tablescan I/O* | *Update I/O* |
|---|---|---|
| Sequential file | N/B | 2N/B |
| Linked list | N | O(1)* |

(sorted)

* If update location known

# Problem session

- Think of a way of storing *sorted* relations on disk such that:
  - Reading all data in the relation is efficient (close to the best possible efficiency).
  - Inserting and deleting data is efficient (assume the location of update is known).
- Do the analysis in the *I/O model.*
- As before, assume that tuples have fixed length, and let B denote the number of tuples per block.

# Amortized analysis

- Instead of analyzing the worst-case cost of a single operation, *amortized analysis* looks at the average cost of a sequence of operations.
- Example: An *unordered* sequential file
  - must use 1 I/O for some insertions, but
  - can do B insertions in 1 I/O using an internal memory buffer
- Conclusion: The amortized cost of an insertion is 1/B I/Os. Tiny!
- Same thing for **your** proposal?

# Aside: Main memory DBMSs

- It is increasingly common to store all (or all the most frequently accessed parts) in main memory.
- Still, non-volatile memory (hard drive, flash,...) is needed to provide durability in case of system crashes.
- Time-space trade-off persists:
  Slower memories are cheaper, i.e., can be used to store much more data.
- Also recall that even main memory, like disks, is *blocked*.

# Exercise from hand-out

- *Representation of relations*
- Questions a), b), and c).

# Searching sorted relations

- Suppose a sequential file of N tuples is sorted by an attribute, A.
- It can be searched for an A-value using binary search, in $\log_2(N/B)$ I/Os.
- A (sorted) linked list may require a full traversal to locate an A-value.

| *Storage method* | *Tablescan* | *Update* | *Search* |
|---|---|---|---|
| Sequential file | N/B | 2N/B | $\log_2(N/B)$ |
| Linked list (blocked) | O(N/B) | O(1) | O(N/B) |

(sorted)

# Adding a directory

- A sorted linked list may be searched more quickly if we have a *directory*:
  - A sorted list with a *representative key* from each block (e.g., smallest key).
  - A pointer to the block of each key.
- But how do we search the directory?
  - Not so important if it fits in RAM.
  - Otherwise, it seems that this problem is *the same* as the original problem...
  - Is there any progress in this case? (Discussion on board.)

# Coping with updates

- Easy solutions:
  - Insertions: If room in the relevant block, ok. Otherwise insert in an *overflow chain*. (The *ISAM* solution.)
  - Deletions: Just mark deleted, don't try to reuse space.
- More robust solution based on the lesson from linked lists:

  *Introduce some "slack" to allow efficient updates.*

# $B^+$-tree invariants on nodes

- Suppose a node (stored in a block) has space for B-1 keys and B pointers.
- Don't want blocks to be too empty: Should be at least half full.
- Let’s see how this works out! (Board.)
- Only possible with an exception: The root may have as little as 1 key and 2 non-null pointers.

http://www.youtube.com/watch?v=coRJrcIYbF4&NR=1

# $B^+$-tree properties

- Search and update cost O(depth) I/Os.
- A $B^+$-tree of depth d holds at least $2\ (B/2)^{d-1}$ keys
- Rewriting, this means that $d < \log_{B/2}(N/2) = O(\log_B N)$
- Often the top level(s) will reside in internal memory. Then the operation time becomes $O(\log_B(N/M))$.

# Aside: Sorting using $B^+$-trees

- In internal memory, sorting can be done in O(N log N) time by inserting the N keys into a balanced search tree.
- The number of I/Os for sorting by inserting into a B-tree is $O(N \log_B N)$.
- This is more than a factor B slower (!) than multiway mergesort (Feb 26 lecture).
- Moral: What works well on internal memory may fail miserably on disk.

# From your toolbox: Search trees

- You may have seen: AVL-trees, red-black trees, or (2,4)-trees.
- All have maximum degree 2 or 4, and depth O(log N).
- How do they compare to $B^+$-trees on external memory?

- Rough analysis: $\log(N) = \log_B(N)\log(B)$
- Factor around 5-10 depending mainly on the outdegree of the $B^+$-tree.

# Buffering in B-trees

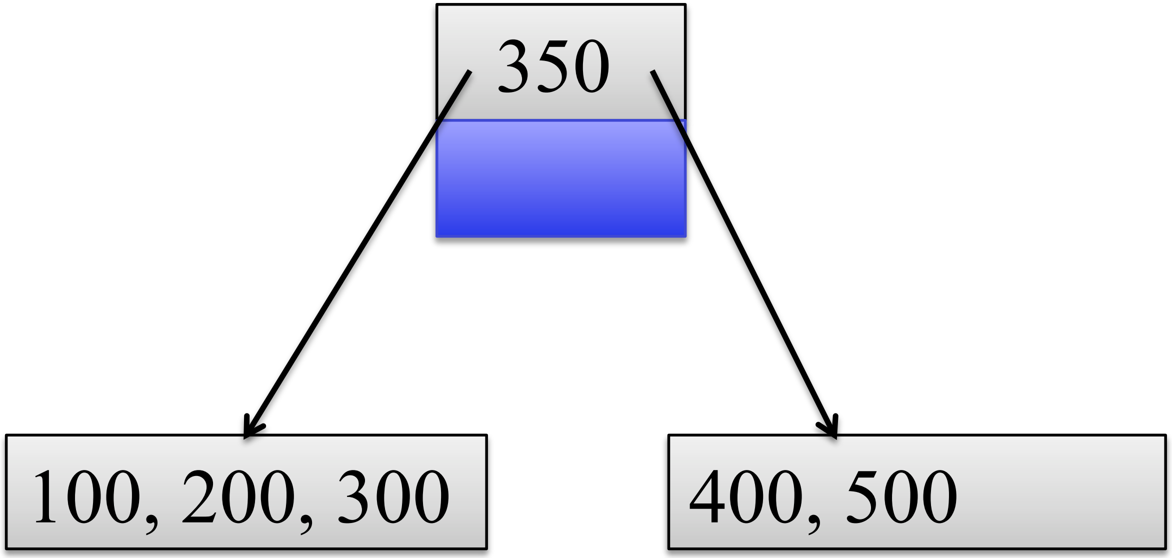
- Relatively new technique to speed up updates in external search trees.
  - Early work in [Arge '96, O'Neil et al '96, Jermaine et al '99, Buchsbaum et al 2000]
  - Explicit in [Brodal and Fagerberg 2003]
  - Newer study by [Graefe 2007]
- We will consider:
  1. A simple approach for tables that are not too much larger than main memory.
  2. An implementation based on constant degree search trees. <u>Main idea</u>: Use buffers, do many things per I/O. (Board.)

# Simple approach

- Suppose relation R is frequently updated.
- Create a small table $R_U$ that records recent updates.
  - Table and index stays in RAM.
  - A view is made available that gives access to the table with the updates performed.
- When the size of $R_U$ becomes close to the amount of RAM available, transfer all the updates to R (on disk).
  - Each I/O used to make many updates.

# Buffered B-trees

- Discussion on board…



# Buffering summary

- Comparing B+-trees and buffered trees

| *Index* | *Search I/O* | *Update I/O* |
|---|---|---|
| B+-tree | $O(\log_B N)$ | $O(\log_B N)$ |
| Buffered B-tree | $O(\log N)$ | $O(\log(N)/B)$ |

- In general, the outdegree is a parameter that can be chosen to trade off search time and update time. (See [BF03] for such a trade-off.)

# More on B-trees

- Claim in textbook: “Rebalancing on higher levels of a B-tree occurs very rarely.”
- This is true (in particular at the top levels), but a little hard to see.
- Easier seen for *weight-balanced* B-trees. Details in [Pagh03].
- Just one of many B-tree variants…

# Yet more on B-trees

- In practice, variable length (possibly long) keys must be handled.
- Later in course: The *String B-tree*, an elegant solution that is efficient even for long strings.
- Space-saving trick: Key compression. See RG for details.
- Building a B-tree: Repeated insertion is much slower than “bulk-loading” the sorted list of keys. (Feb. 26)

# Two types of indexes

- In a primary index, records are stored in an order determined by the search key.
- A relation can have at most one primary index. (Often on the primary key.)
- A secondary index cannot take advantage of any specific order. Must contain all keys and corresponding references to tuples ("dense").
- Sometimes, a secondary index provides a faster access path than a primary index! (More about that next week.)

# Summary

- We saw the basics of storing and indexing data on external memory.
- Many trade-offs (soft or hard):
  - Space usage vs update efficiency.
  - Update vs access efficiency.
  - Optimizing for one kind of access may slow down another type of access.
- We have started thinking and analyzing in terms of I/Os.

# Exercise from hand-out

- *$B^+$-trees*
- Questions a), b), and c).