

# Modeling, Simulation, Verification & Code Generation with IAR visualSTATE

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# Outline

- Introducing the modeling language (air conditioner example).
- Tool demo (modeling, simulation, verification and code generation).
- Discussion of generated code.
- Usage contexts (specialization, validation, monitored execution).

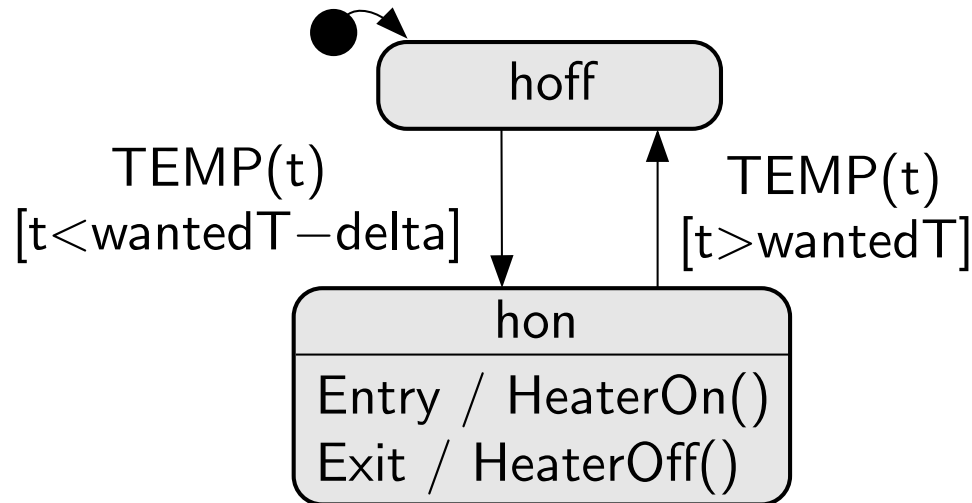
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# Trivialized Air Conditioning System

- Components:
  - heater
  - cooler
  - fan
  - user interface
  - goal temperature display
  - current fan speed display
- User can set a goal temperature in the room.
- System uses the heater and cooler to achieve the temperature.
- Efficiency of conditioning is controlled by the speed of fan.
  - manual mode: user controls the speed
  - automatic: a fixed, factory predefined speed
  - fast: fan at maximum speed.

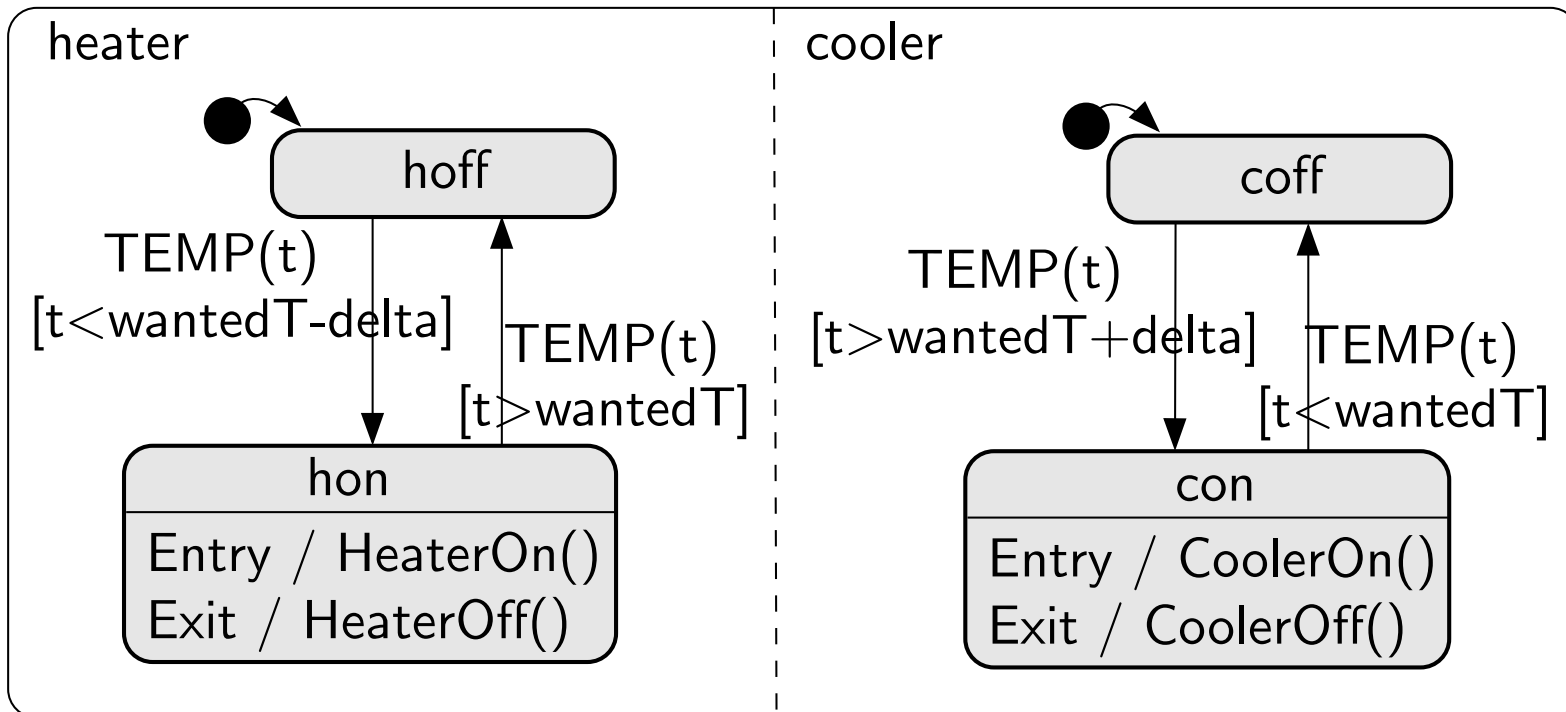
# Heater



- Heater is initially in *hoff*.
- A periodic process supplies  $TEMP(t)$  regularly.
- Value  $t$  is the current temperature returned by the sensor.
- Variable `int wantedT` stores current goal temperature.
- Constant `int delta` gives the acceptable error.
- Heater is activated (*HeaterOn*), whenever *hon* is entered.
- Heater is deactivated (*HeaterOff*) on exit from *hon*.

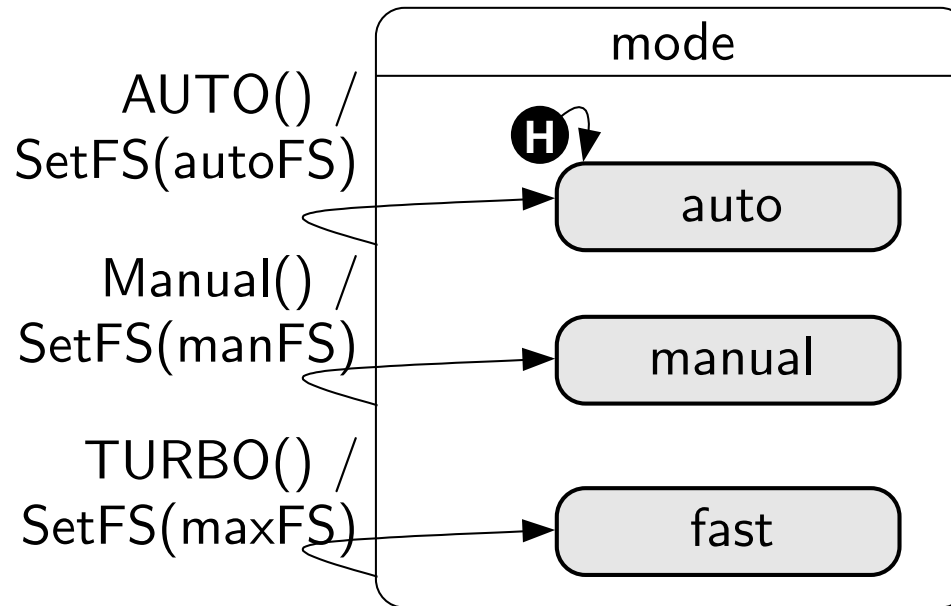
# Heater & Cooler

Cooler and heater are analogous and independent.



Compose them together.

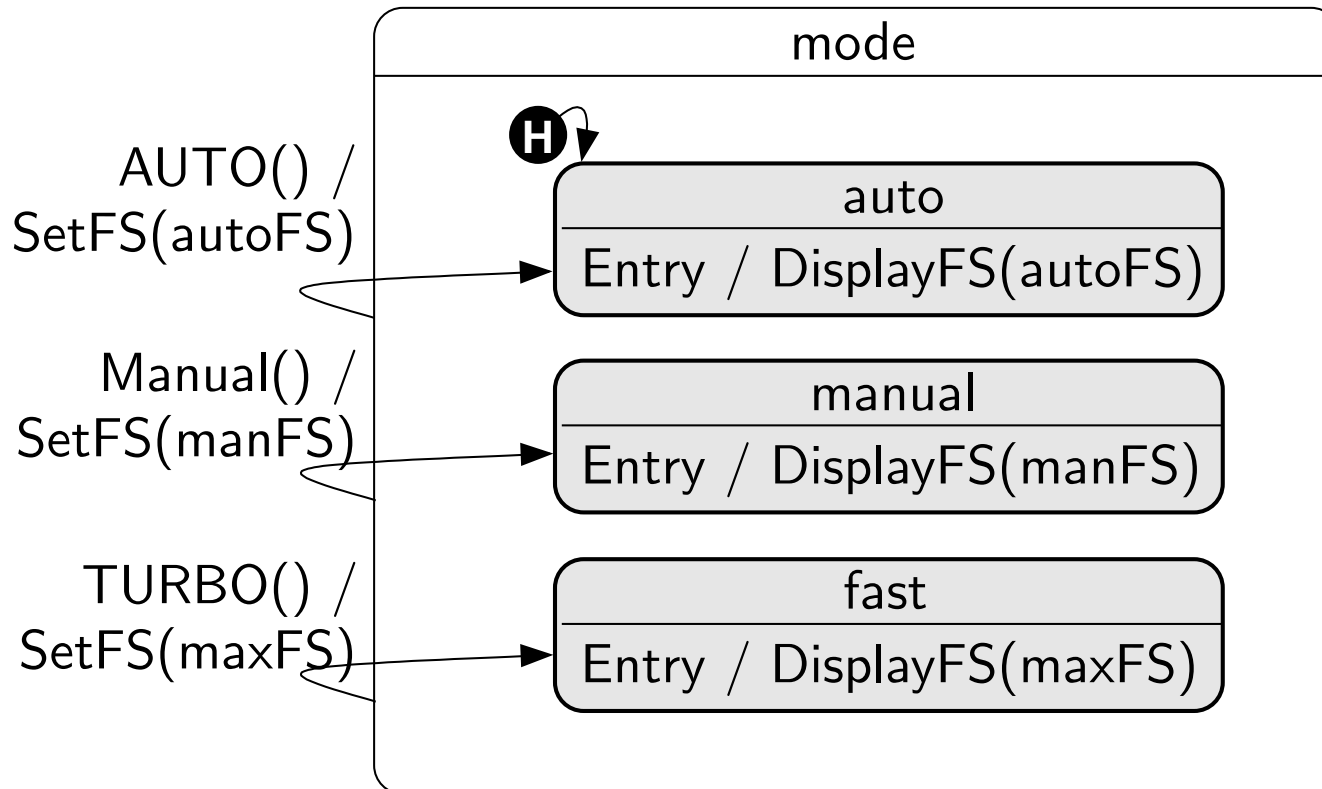
# User Interface



- Buttons: AUTO, TURBO, TINC, TDEC, FSINC, FSDEC.
- Button group: Manual = { FSINC, FSDEC }.
- Variable manFS stores fan speed manually adjusted by user.
- Constant autoFS specifies factory-set automatic mode speed.
- Constant maxFS specifies maximum possible fan speed.
- History state – remember value across executions.

# User Interface

(II)



Interface has got 2 displays: current fan speed and goal temperature. *DisplayFS* updates the value shown for fan speed. Generally, state actions help to guarantee state invariants.



# User Interface

(III)

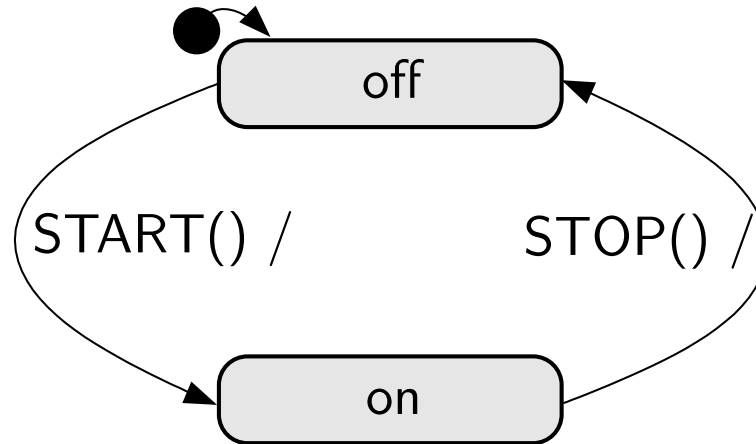
- Internal rules for adjusting the goal temperature and fan speed:

TINC  $[wantedT < maxT] / [wantedT = wantedT + 1]$  DisplayT( $wantedT$ )  
TDEC  $[wantedT > minT] / [wantedT = wantedT - 1]$  DisplayT( $wantedT$ )  
FSINC  $[manFS < maxFS] / [manFS = manFS + 1]$   
FSDEC  $[manFS < minFS] / [manFS = manFS - 1]$

- These rules use variables instead of states.
- As with normal transitions they fire, whenever guards are satisfied.
- They should be active whenever the user interface is active.
- Variable  $wantedT$  stores the goal temperature set by user.
- Variable  $manFS$  stores the desired fan speed.
- Maximum&minimum goal temperature: constants  $maxT, minT$ .
- Maximum&minimum fan speed: constants  $maxFS, minFS$ .

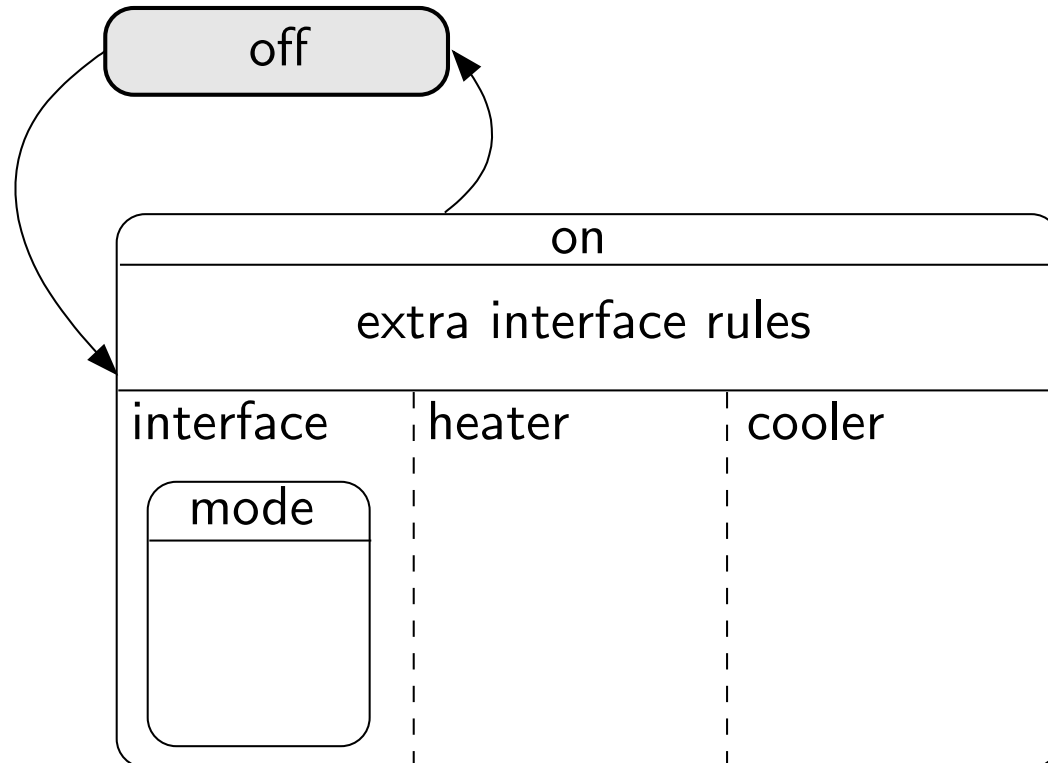
# Top Level

START and STOP buttons turn the machinery on and off.



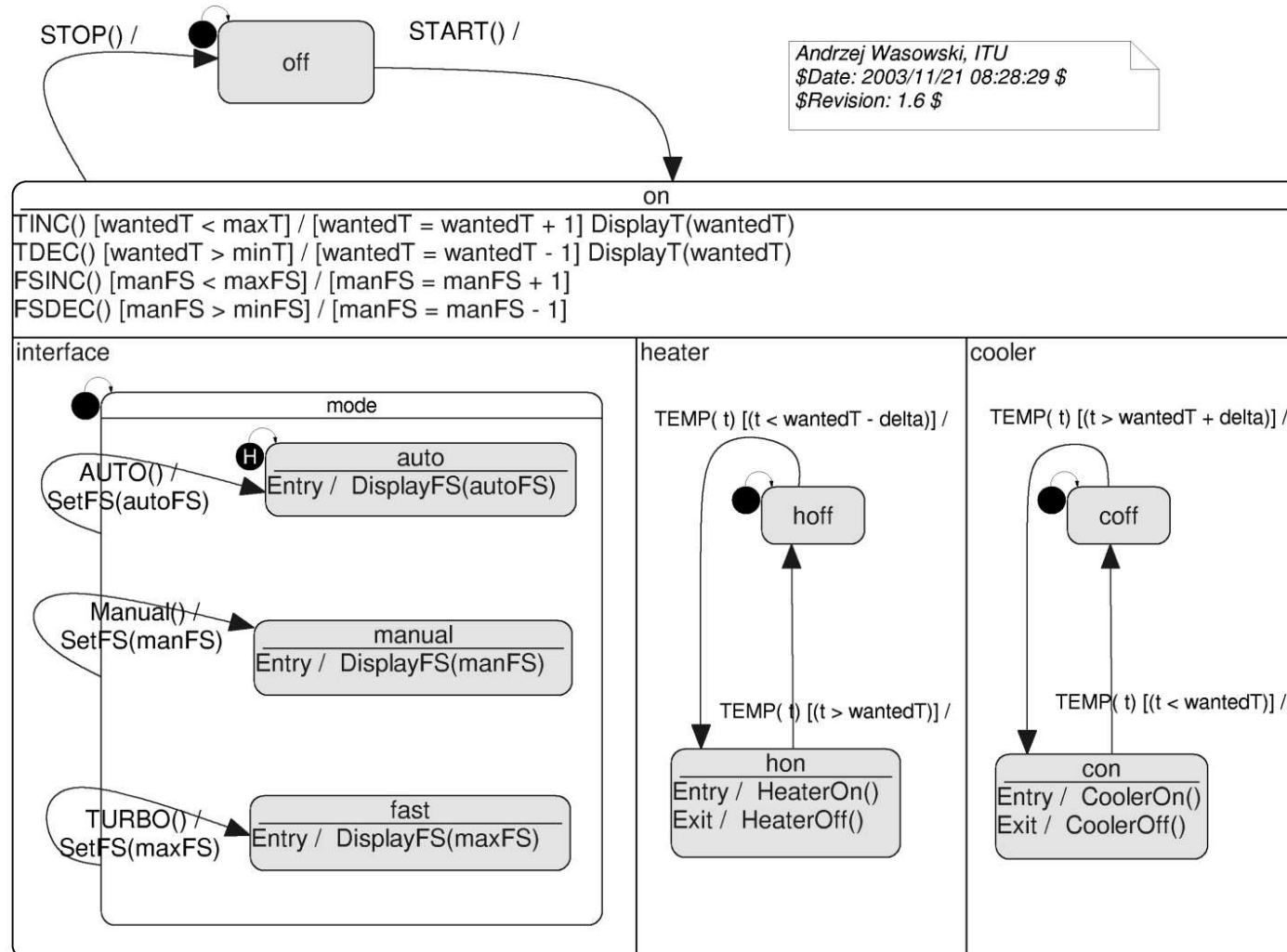
Some entry and exit actions should be added.

# System Overview



# Complete Model of Air Conditioner

The complete visualSTATE model (tool printout):



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# IAR visualSTATE Demo

- Designer
- Simulator
- Verifier
- Code Generator

# visualSTATE model checker

Model checker automatically verifies if following hold in the model:

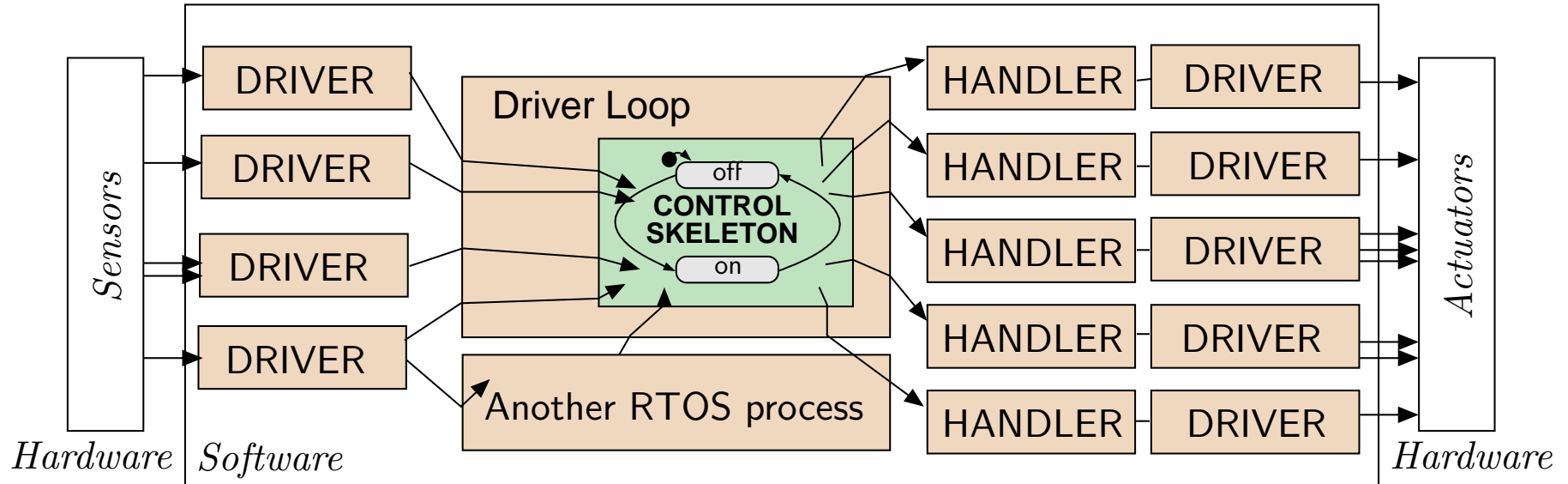
- No unused components [states, variables]
- No unreachable guards. It must be possible to enable all of the guards in the system. This means that there must exist a reachable state for each guard  $g$  that enables this guard. Unreachable guards mean dead code (dead transitions).
- No conflicting transitions.
- No deadlocks.
- No illegal operations. Arithmetic operations should be checked for overflow and illegal operations such as division by zero.
- No divergent behavior. If the signal queue is used then the macrostep should always be finite.
- No overflow of the signal queue.

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# Application Components



- Generated control skeleton (green).
- Brown parts hand coded, but fortunately small and easier.
- Sometimes multiple processes are avoided in favour of the loop.
- In some cases it is even possible to give up the RTOS entirely.
- Discuss the cost of automatic generation of the skeleton.

# Executable Size [Control Algorithm]

The table presents executable sizes for the air conditioner model (only control code, no RTOS, no action functions, interfaces to sensors, etc):

platform	compiler	optimizations	cod. gen.	size [b]
i386/Linux	gcc 3.2	-Os + strip	IAR VS 4.3	4 428
i386/Linux	gcc 3.2	-Os + strip	SCOPE	3 732
h8300	gcc 2.95	-O2 + strip	IAR VS 4.3	8 528
h8300	gcc 2.95	-O2 + strip	SCOPE	7 922
h8300	gcc 3.3	-O2 + strip	IAR VS 4.3	<b>2 388</b>
h8300	gcc 3.3	-O2 + strip	SCOPE	<b>1 822</b>

The gcc 3.3 reported is an experimental version and executables were not tested. You know better if commercial compilers can be expected to generate more efficiently.

# Memory Consumption ctd.

- Following executable sizes given for gcc 3.3 on H8/300:
- The visualSTATE kernel (compiled with dummy model) takes 1.5k
- Complex model of coffee machine (200 transitions) is below 7k
- RAM usage in SCOPE [quick generous estimate, assuming 8bit word, 32bit addressing]

	[bytes]
current event (global)	3
state representation (global)	8
stack	30
model variables (global)	4
<b>TOTAL</b> [bytes]	<b>41+4</b>

- More expensive if signal communication is used.
- VisualSTATE has similar performance.
- If this is not sufficient we can try targeting assembler directly.

# Code Excerpts [SCOPE]

Most of the code take up read-only tables:

```
/* and-state projection of hierarchy */
const anatomycell anatomy[10] = {
    /* 0 */ STMRK, MCHN 3, MCHN 3, MCHN 0, MCHN 2,
    /* 5 */ MCHN 2, MCHN 1, MCHN 1, MCHN 1, MCHN 0,
};

/* transitions array */
const transcell trans[TRANS_MAX] = {
    /* 0 */ PCNC(2) 1, 0, STATE 3, ACGD(2) 4, 1,
    /* 5 */ STATE 8, STMRK, PCNC(2) 1, 0, STATE 3,
    /* 10 */ ACGD(2) 6, 1, STATE 7, STMRK, PCNC(2) 1,
    /* 15 */ 0, STATE 3, ACGD(2) 6, 1, STATE 7,
    /* 20 */ STMRK, PCNC(2) 1, 0, STATE 9, ACGD(2) 1,
    /* 25 */ 1, STATE 3, STATE 2, STATE 5, STMRK,
    /* 30 */ PCNC(2) 2, 0, STATE 3, STATE 1, ACGD(2) 3,
    ...
}
```

# Code Excerpts [SCOPE]

(II)

```
/* guards dispatcher */
int eval ( const guardref g ) {
    switch (g) {
        case 1: return ((CurrEvent.fields._E_TEMP.f0)
                        <(wantedT));
        case 2: return ((CurrEvent.fields._E_TEMP.f0)
                        >((wantedT)+(delta)));
        case 3: return ((CurrEvent.fields._E_TEMP.f0)
                        >(wantedT));
        case 4: return ((CurrEvent.fields._E_TEMP.f0)
                        <((wantedT)-(delta)));
        case 5: return (1);
        ...
    }
    return 1;
}
```

# Code Excerpts [SCOPE]

(III)

```
/* Main loop */

int main (void)
{
    st_init();          // initialize the system
    while (1) {

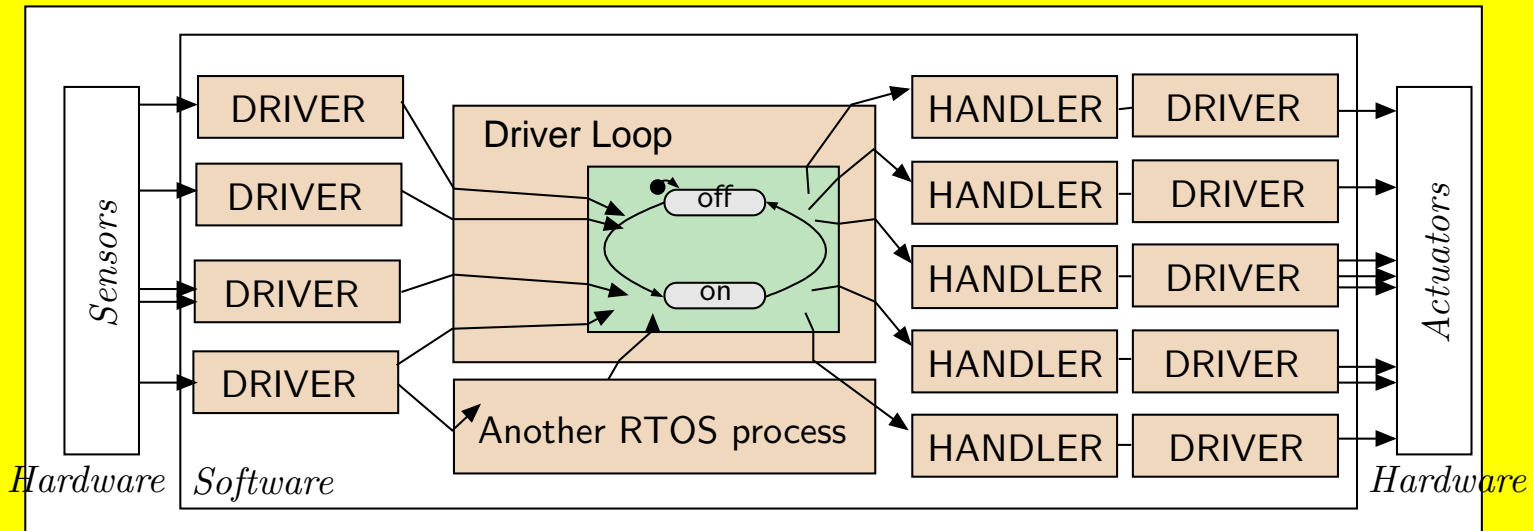
        ...            // compute the next event in i

        CurrEvent.tag = i;
        macrostep(); // call the VS kernel
    }
    return 0;
}
```

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# Usage Contexts [in progress]



## Execution Context (reality...)

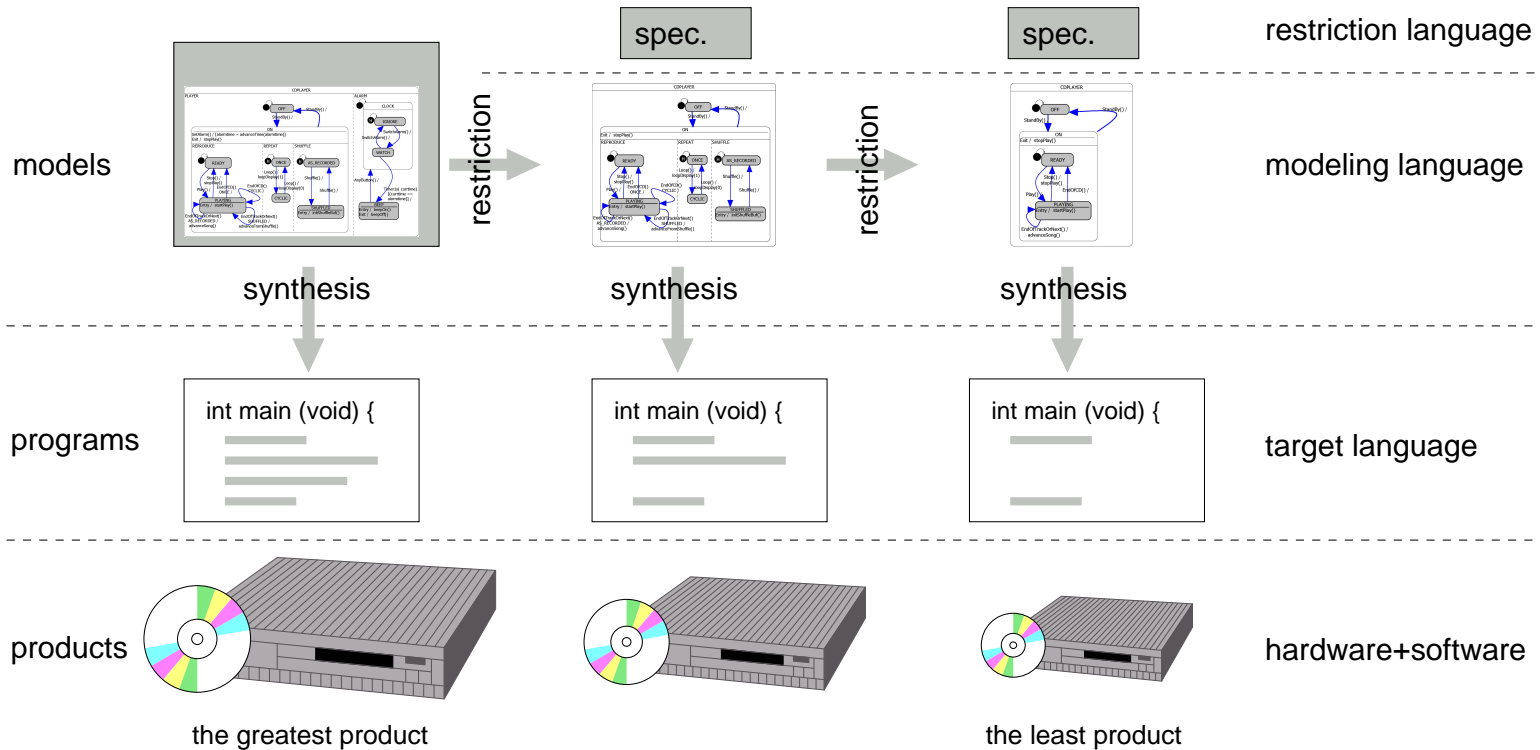
- All systems are executed in contexts which are limited in some ways.
- Context limits possible interactions of the embedded systems.
- Modeling contexts adds another level of value to the development process.



# Usage Contexts [in progress] (II)

- Test case generation is improved. Resources are not wasted on executing completely unrealistic cases. Number of false negatives is reduced.
- Number of false negatives is reduced for verification too.
- Monitored execution: alarms can be generated at runtime whenever system (or environment) violates the assumed contract.
- Specialization: possibility of generating code specifically optimized for given user context. Supports architecture of product line based on single source code.
- Air conditioner can be specialized to a heater, a cooler, a device with less than 3 modes, a device without display, etc.
- A single model for all those.
- A small context specification for each of those.
- Automatic instantiation of general model in specific context
- We need your input and examples on what contexts are.

# Fast Generation of User Variants



# Fast Generation of User Variants (II)

```
restriction WithoutAlarm {  
    impossible SetAlarm();  
    impossible SwitchAlarm();  
};  
WithoutAlarm CDPLAYER;
```

```
restriction Least restricts WithoutAlarm {  
    impossible Loop();  
    impossible Shuffle();  
};
```

```
Least CDPLAYER;
```

- Hierarchies of contexts can be build.
- More dynamic context properties can be expressed with automata. Especially useful for test generation and automatic verification.
- Keen to see what kind of properties are needed.

# Summary

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Thank you for  
Your attention.