Programs as Data
Backwards optimizing compilation of micro-C

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Today

- Husk kursusevaluering denne uge!
- Spørgetime, 2.-5. januar; hvornår?
  (Eksamen ligger pt. mandag 9. januar)
- Prøveeksamen mandag 28. november??
- Genopfriskning af typeinferens-træer

- Deficiencies in the old micro-C compiler
- A backwards, continuation-based, compiler
- Optimizing code on the fly

Micro-C compiler shortcomings

Last slide from lecture 8

- The compiler often generates **inefficient code**

```
GETBP
CSTI 0
ADD
LDI
```

could be

```
GETBP
LDI
```

```
INCSP -1
INCSP -1
```

could be

```
INCSP -2
```

- The **compiler itself is inefficient**, using (@) a lot:

```
| If(e, stmt1, stmt2) ->
  let labelse = newLabel()
  let labend  = newLabel()
  in cExpr e varEnv funEnv @ [IFZERO labelse]
     @ cStmt stmt1 varEnv funEnv @ [GOTO labend]
     @ [Label labelse] @ cStmt stmt2 varEnv funEnv
     @ [Label labend]
```

- **Tail calls** are not executed in constant space

Example, if-statement (ex19.c)

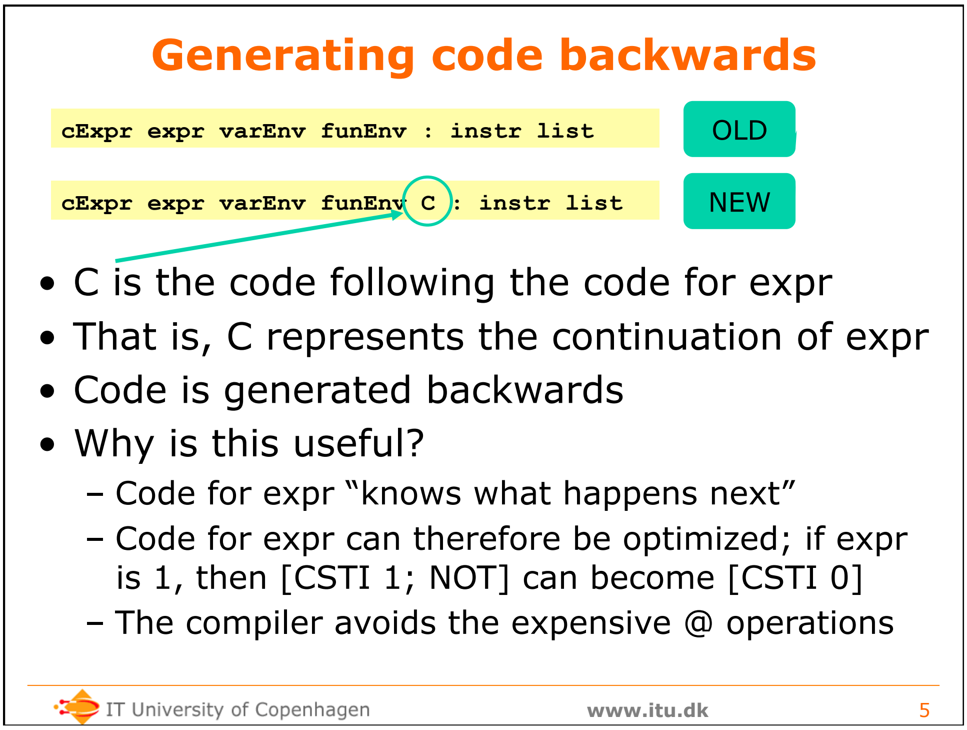
```
void main(int x) {
  if (x == 0) print 33; else print 44;
}
```

- The old micro-C compiler produces this:

```
    GETBP; CSTI 0; ADD; LDI; CSTI 0; EQ; IFZERO L2;
    CSTI 33; PRINTI; INCSP -1; GOTO L3;
L2: CSTI 44; PRINTI; INCSP -1;
L3: INCSP 0; RET 0
```

- We want it produce this better code:

```
    GETBP; LDI; IFNZRO L2;
    CSTI 33; PRINTI; RET 1;
L2: CSTI 44; PRINTI; RET 1
```



Generating code backwards

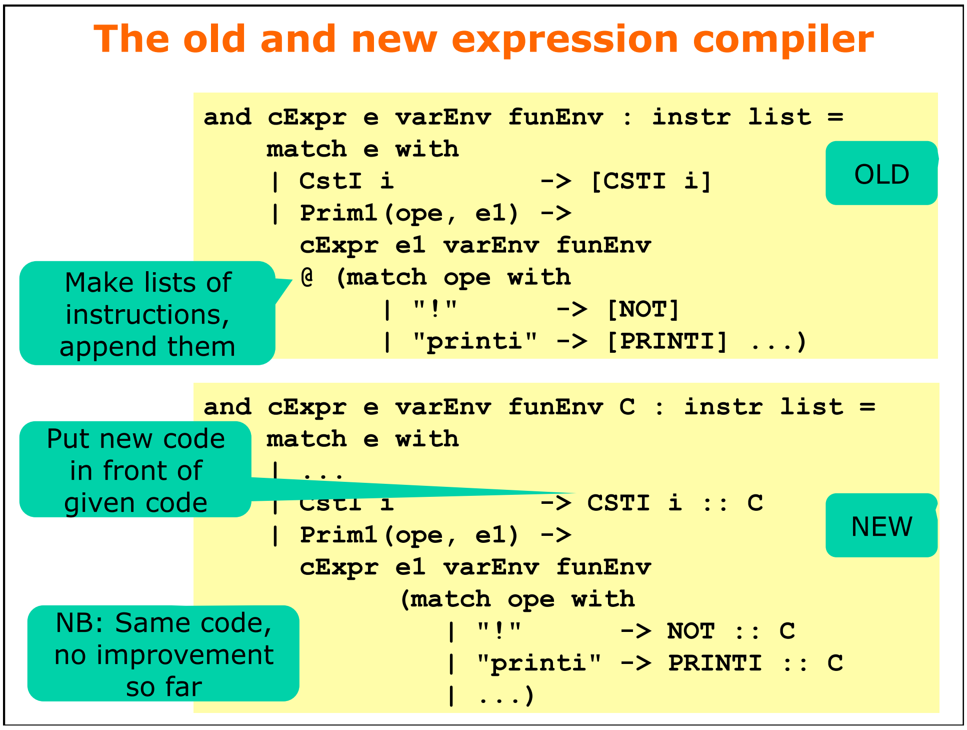
```
cExpr expr varEnv funEnv : instr list
```

OLD

```
cExpr expr varEnv funEnv C : instr list
```

NEW

- C is the code following the code for expr
- That is, C represents the continuation of expr
- Code is generated backwards
- Why is this useful?
  - Code for expr “knows what happens next”
  - Code for expr can therefore be optimized; if expr is 1, then [CSTI 1; NOT] can become [CSTI 0]
  - The compiler avoids the expensive @ operations



The old and new expression compiler

```
and cExpr e varEnv funEnv : instr list =
    match e with
    | CstI i          -> [CSTI i]
    | Prim1(ope, e1) ->
      cExpr e1 varEnv funEnv
      @ (match ope with
           | "!"       -> [NOT]
           | "printi" -> [PRINTI] ...)
```

OLD

Make lists of instructions, append them

```
and cExpr e varEnv funEnv C : instr list =
    match e with
    | ...
    | CstI i          -> CSTI i :: C
    | Prim1(ope, e1) ->
      cExpr e1 varEnv funEnv
            (match ope with
               | "!"       -> NOT :: C
               | "printi" -> PRINTI :: C
               | ...)
```

NEW

Put new code in front of given code

NB: Same code, no improvement so far

Code improvement rules based on bytecode equivalences

Code	Better equivalent code	When
0; ADD	<no code>	
0; SUB	<no code>	
0; NOT	1	
n; NOT	0	n ≠ 0
1; MUL	<no code>	
1; DIV	<no code>	
0; EQ	NOT	
INCSP 0	<no code>	
INCSP m; INCSP n	INCSP (m+n)	
0; IFZERO a	GOTO a	
n; IFZERO a	<no code>	n ≠ 0
0; IFNZRO a	<no code>	
n; IFNZRO a	GOTO a	n ≠ 0

Joint exercise 1 (code for ex13.c)

source

```
void main(int n) {
  int y;
  y = 1889;
  while (y < n) {
    y = y + 1;
    if (y % 4 == 0 && y % 100 != 0 || y % 400 == 0)
      print y;
} }
```

generated by old compiler

```
CSTI 0; GETBP; CSTI 1; ADD; CSTI 1889; STI; INCSP -1; GOTO L3; L2:
GETBP; CSTI 1; ADD; GETBP; CSTI 1; ADD; LDI; CSTI 1; ADD; STI;
INCSP -1; GETBP; CSTI 1; ADD; LDI; CSTI 4; MOD; CSTI 0; EQ; IFZERO
L9; GETBP; CSTI 1; ADD; LDI; CSTI 100; MOD; CSTI 0; EQ; NOT; GOTO
L8; L9:; CSTI 0; L8: IFNZRO L7; GETBP; CSTI 1; ADD; LDI; CSTI 400;
MOD; CSTI 0; EQ; GOTO L6; L7: CSTI 1; L6: IFZERO L4; GETBP; CSTI 1;
ADD; LDI; PRINTI; INCSP -1; GOTO L5; L4: INCSP 0; L5: INCSP 0; L3:
GETBP; CSTI 1; ADD; LDI; GETBP; CSTI 0; ADD; LDI; LT; IFNZRO L2;
INCSP -1; RET 0
```

- Simplify the bytecode, eg. by deleting superfluous instructions

Optimizing code while generating it

```
and cExpr e varEnv funEnv C : instr list =
    match e with
    | CstI i           -> addCST i C
```

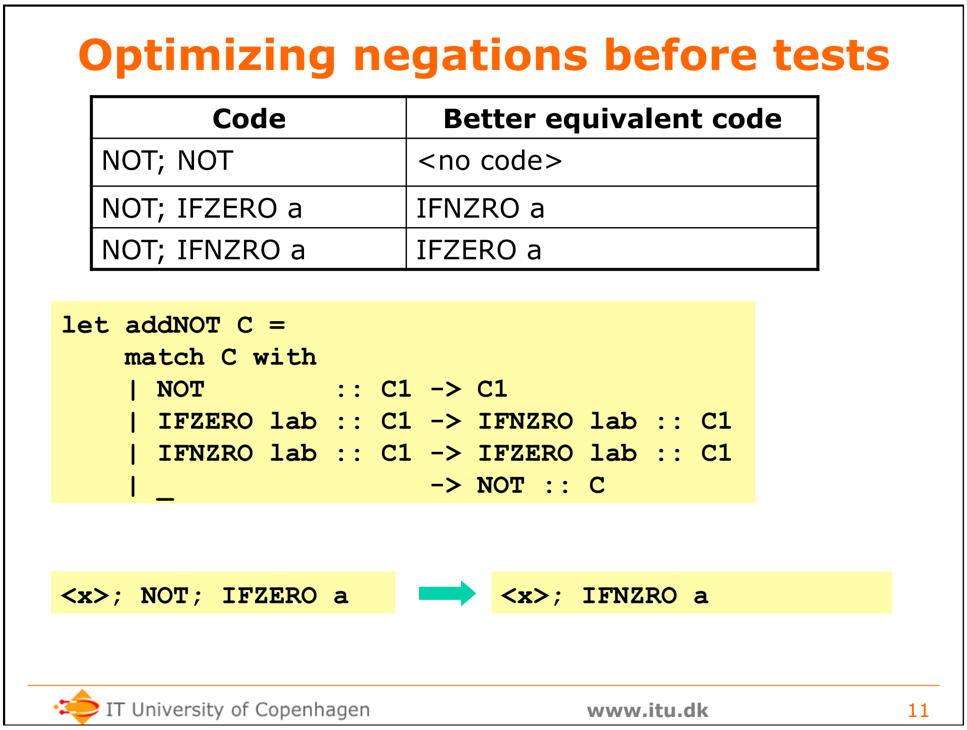
```
let rec addCST i C =
    match (i, C) with
    | (0, ADD         :: C1) -> C1
    | (0, SUB         :: C1) -> C1
    | (0, NOT         :: C1) -> addCST 1 C1
    | (_, NOT         :: C1) -> addCST 0 C1
    | (1, MUL         :: C1) -> C1
    | (1, DIV         :: C1) -> C1
    | (0, EQ          :: C1) -> addNOT C1
    | (0, IFZERO lab :: C1) -> addGOTO lab C1
    | (_, IFZERO lab :: C1) -> C1
    | ...
    | _                      -> CSTI i :: C
```

If no optimization possible, generate CSTI instruction

Some examples

`GETBP; 0; ADD; LDI` → `GETBP; LDI`

`<x>; 0; EQ; IFZERO a` → `<x>; NOT; IFZERO a`

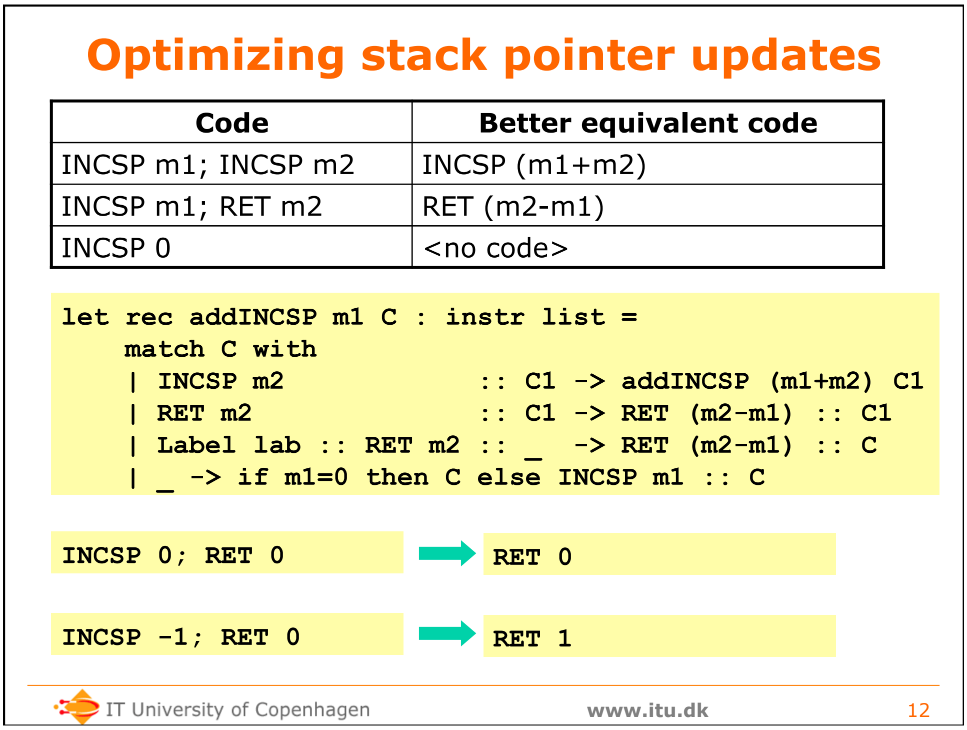


Optimizing negations before tests

Code	Better equivalent code
NOT; NOT	<no code>
NOT; IFZERO a	IFNZRO a
NOT; IFNZRO a	IFZERO a

```
let addNOT C =
    match C with
    | NOT         :: C1 -> C1
    | IFZERO lab :: C1 -> IFNZRO lab :: C1
    | IFNZRO lab :: C1 -> IFZERO lab :: C1
    | _                 -> NOT :: C
```

`<x>; NOT; IFZERO a` → `<x>; IFNZRO a`



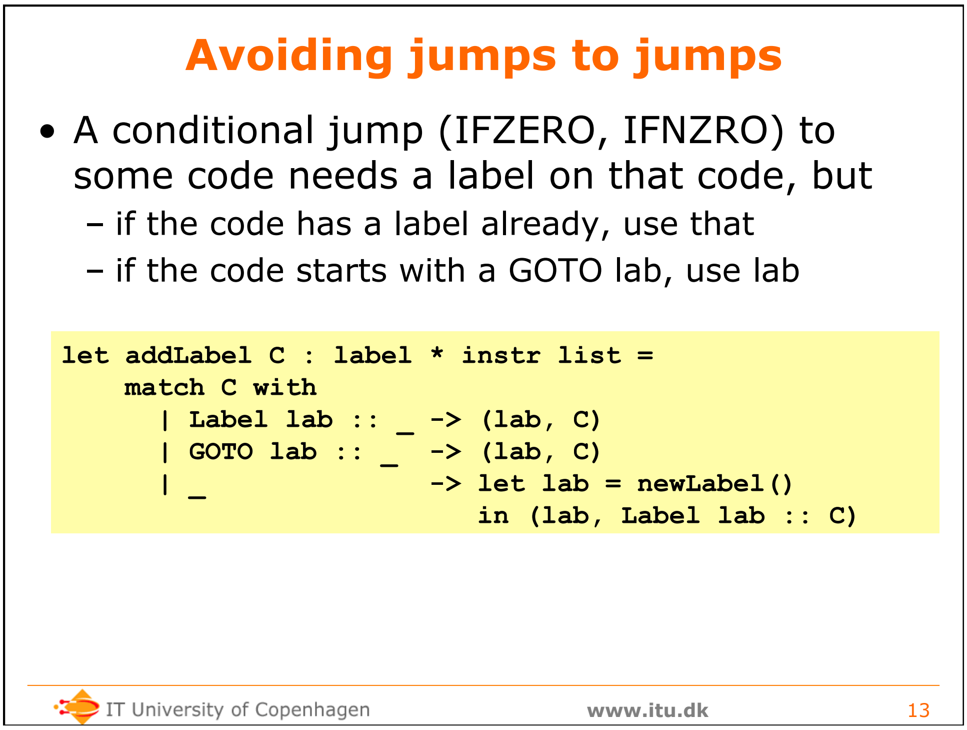
Optimizing stack pointer updates

Code	Better equivalent code
INCSP m1; INCSP m2	INCSP (m1+m2)
INCSP m1; RET m2	RET (m2-m1)
INCSP 0	<no code>

```
let rec addINCSP m1 C : instr list =
    match C with
    | INCSP m2              :: C1 -> addINCSP (m1+m2) C1
    | RET m2                :: C1 -> RET (m2-m1) :: C1
    | Label lab :: RET m2 :: _   -> RET (m2-m1) :: C
    | _ -> if m1=0 then C else INCSP m1 :: C
```

`INCSP 0; RET 0` → `RET 0`

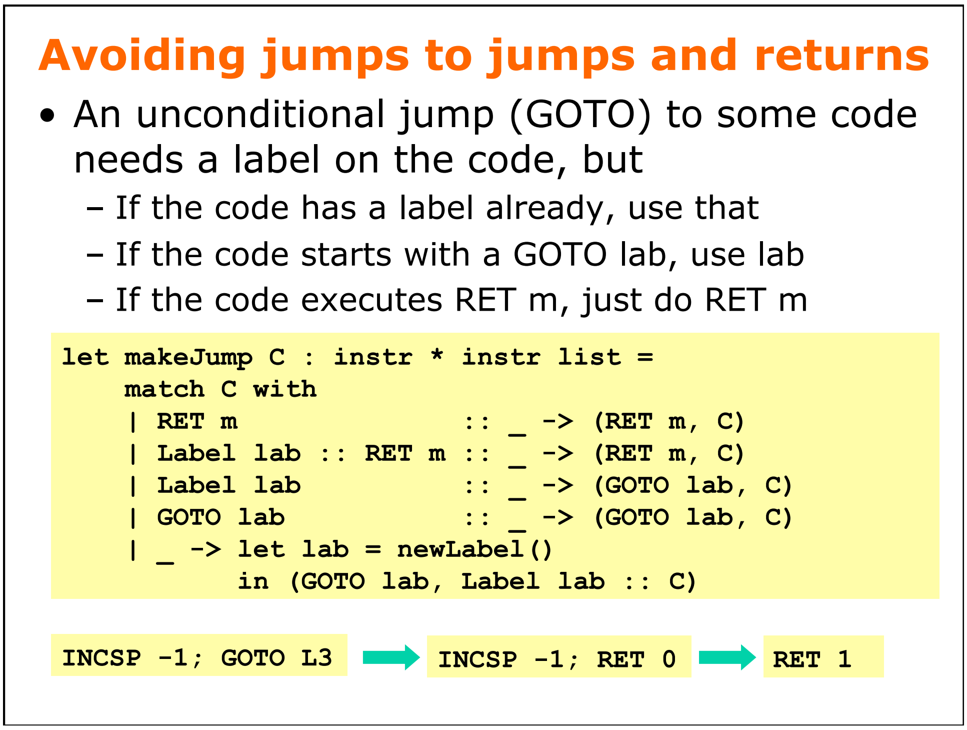
`INCSP -1; RET 0` → `RET 1`



Avoiding jumps to jumps

- A conditional jump (IFZERO, IFNZRO) to some code needs a label on that code, but
  - if the code has a label already, use that
  - if the code starts with a GOTO lab, use lab

```
let addLabel C : label * instr list =
    match C with
      | Label lab :: _ -> (lab, C)
      | GOTO lab :: _  -> (lab, C)
      | _              -> let lab = newLabel()
                             in (lab, Label lab :: C)
```

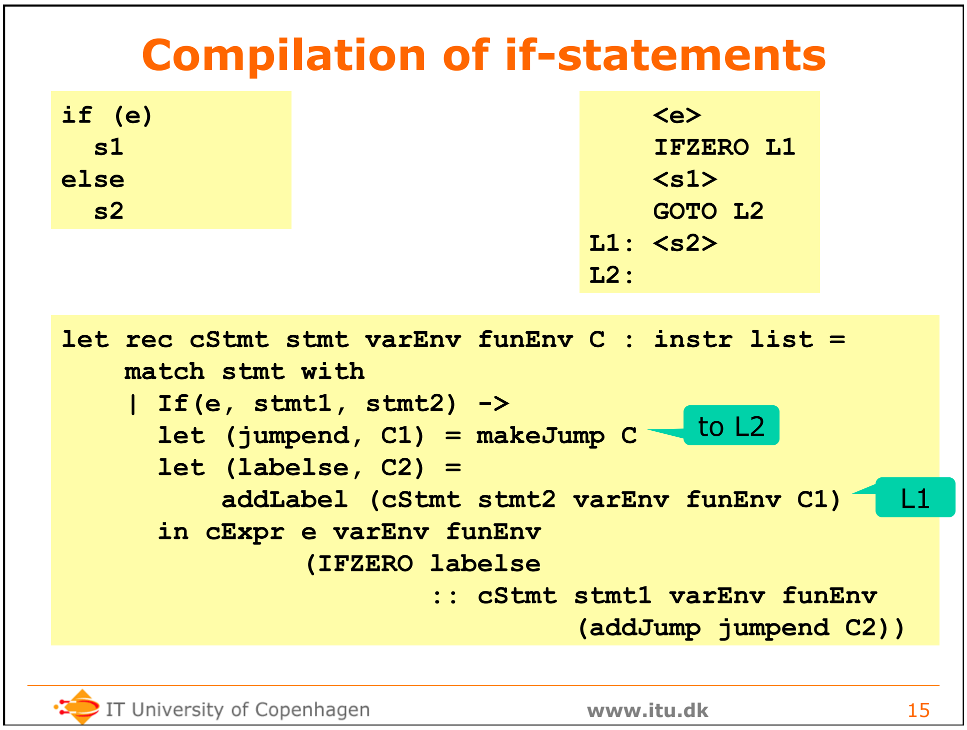


Avoiding jumps to jumps and returns

- An unconditional jump (GOTO) to some code needs a label on the code, but
  - If the code has a label already, use that
  - If the code starts with a GOTO lab, use lab
  - If the code executes RET m, just do RET m

```
let makeJump C : instr * instr list =
    match C with
    | RET m                 :: _ -> (RET m, C)
    | Label lab :: RET m :: _ -> (RET m, C)
    | Label lab             :: _ -> (GOTO lab, C)
    | GOTO lab              :: _ -> (GOTO lab, C)
    | _ -> let lab = newLabel()
           in (GOTO lab, Label lab :: C)
```

`INCSP -1; GOTO L3` → `INCSP -1; RET 0` → `RET 1`



Compilation of if-statements

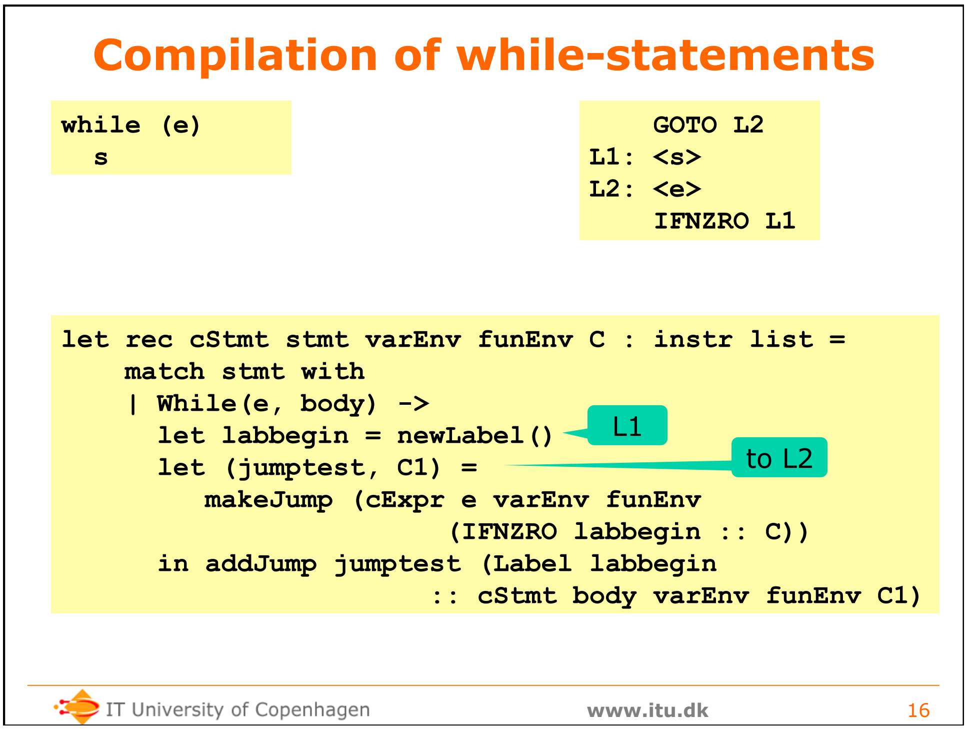
```
if (e)
  s1
else
  s2
```

```
     <e>
     IFZERO L1
     <s1>
     GOTO L2
L1:  <s2>
L2:
```

```
let rec cStmt stmt varEnv funEnv C : instr list =
    match stmt with
    | If(e, stmt1, stmt2) ->
      let (jumpend, C1) = makeJump C
      let (labelse, C2) =
          addLabel (cStmt stmt2 varEnv funEnv C1)
      in cExpr e varEnv funEnv
             (IFZERO labelse
                     :: cStmt stmt1 varEnv funEnv
                                (addJump jumpend C2))
```

to L2

L1



Compilation of while-statements

```
while (e)
  s
```

```
     GOTO L2
L1:  <s>
L2:  <e>
     IFNZRO L1
```

```
let rec cStmt stmt varEnv funEnv C : instr list =
    match stmt with
    | While(e, body) ->
      let labbegin = newLabel()
      let (jumptest, C1) =
         makeJump (cExpr e varEnv funEnv
                        (IFNZRO labbegin :: C))
      in addJump jumptest (Label labbegin
                     :: cStmt body varEnv funEnv C1)
```

L1

to L2

Compiling shortcut logical expressions

- Logical expression (`m==0 && n==0`) may
  - produce a value, as in `b = (m==0&&n==0);`
  - decide a test, as in `if (m==0&&n==0) ...`

```
     <e1>
     IFZERO L1
     <e2>
     GOTO L2
L1:  0
L2:
```

Standard code for value of `e1&&e2`

When used in if (…) …

```
     <e1>
     IFZERO L1
     <e2>
     GOTO L2
L1:  0
L2:  IFZERO L3
```

… we can optimize it

```
     <e1>
     IFZERO L3
     <e2>
     IFZERO L3
```

Compiling e1 && e2

```
and cExpr e varEnv funEnv C : instr list =
    match e with
    | Andalso(e1, e2) ->
      match C with
        | IFZERO lab :: _ ->
          cExpr e1 varEnv funEnv
             (IFZERO lab :: cExpr e2 varEnv funEnv C)
        | IFNZRO labthen :: C1 ->
          let (labelse, C2) = addLabel C1
          in cExpr e1 varEnv funEnv
             (IFZERO labelse
              :: cExpr e2 varEnv funEnv (IFNZRO labthen :: C2))
        | _ ->
          let (jumpend,  C1) = makeJump C
          let (labfalse, C2) = addLabel (addCST 0 C1)
          in cExpr e1 varEnv funEnv
              (IFZERO labfalse
               :: cExpr e2 varEnv funEnv (addJump jumpend C2))
    | Orelse(e1, e2) -> ... dual to Andalso ...
```

`if (e1&&e2)`

`while (e1&&e2)`

`b = (e1&&e2)`

Joint exercise 2 (code for ex13.c)

source

```
void main(int n) {
  int y;
  y = 1889;
  while (y < n) {
    y = y + 1;
    if (y % 4 == 0 && y % 100 != 0 || y % 400 == 0)
      print y;
} }
```

generated by optimizing compiler

```
CSTI 0; GETBP; CSTI 1; ADD; CSTI 1889; STI; INCSP -1; GOTO
L3; L2: GETBP; CSTI 1; ADD; GETBP; CSTI 1; ADD; LDI; CSTI 1;
ADD; STI; INCSP -1; GETBP; CSTI 1; ADD; LDI; CSTI 4; MOD;
IFNZRO L5; GETBP; CSTI 1; ADD; LDI; CSTI 100; MOD; IFNZRO
L4; L5: GETBP; CSTI 1; ADD; LDI; CSTI 400; MOD; IFNZRO L3;
L4: GETBP; CSTI 1; ADD; LDI; PRINTI; INCSP -1; L3: GETBP;
CSTI 1; ADD; LDI; GETBP; LDI; LT; IFNZRO L2; RET 1
```

- Layout so structure is clearly visible
- Compare to code generated from Java or C# in PLCSD figure 9.9

Eliminating dead code

- Dead code is code that cannot be executed:

```
GOTO L1; 1; ADD; PRINTI; INCSP -1; L1
```

Dead code

- Code following GOTO or RET is dead unless preceded by a label:

```
let rec deadcode C =
    match C with
      | []              -> []
      | Label lab :: _  -> C
      | _         :: C1 -> deadcode C1
```

Discard dead code

Compiling tail calls

- A call is a *tail call* if it is the last action of a function, e.g. ex12.c:

```
int main(int n) {
  if (n)
    return main(n-1);
  else
    return 17;
}
```

Tail call; nothing to do after it

- In the code, a tail call is followed by RET:

```
L1: GETBP; CSTI 0; ADD; LDI; IFZERO L2;
    GETBP; CSTI 0; ADD; LDI; CSTI 1; SUB;
    CALL L1; RET 1; GOTO L3
L2: CSTI 17; RET 1;
L3: INCSP 0; RET 0
```

Tail call

Stack machine execution of TCALL

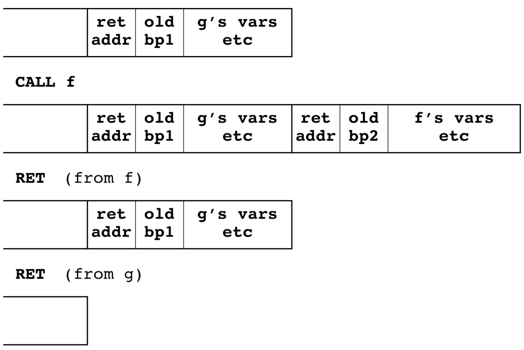
```
int main() { … g(…) … }
int g(…) { return f(…); }
int f(…) { return …; }
```

main calls **g**

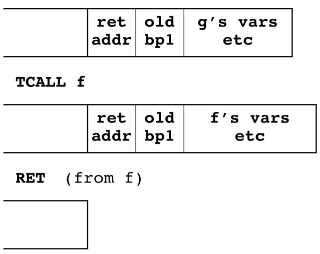
g calls **f**

f returns

Ordinary call and two returns



Tail call and one return



- Return from tail call goes directly to original caller

Micro-C machine call and return

- CALL creates a stack frame
- RET destroys a stack frame
- TCALL destroys one frame and creates another:

19	CALL m a	$s, v_1, .., v_m$	$\Rightarrow$	$s, \boxed{r, bp, v_1, .., v_m}$
20	TCALL m n a	$s, \boxed{r, b, u_1, .., u_n}, v_1, .., v_m$	$\Rightarrow$	$s, \boxed{r, b, v_1, .., v_m}$
21	RET m	$s, \boxed{r, b, v_1, .., v_m}, v$	$\Rightarrow$	s, v

- There is nothing to do after tail call, except return
- So the caller's stack frame can be discarded before the tail call
- So a sequence of tail calls should not require unbounded stack space

Recognizing tail calls in the compiler

- To call a function, compile arguments and emit a call:

```
and callfun f es varEnv funEnv C : instr list =
    ...
    cExprs es varEnv funEnv (makeCall argc labf C)
```

- A tail call is a call followed by RET
- Easy to see when generating code backwards:

```
let makeCall m lab C : instr list =
    match C with
    | RET n            :: C1 -> TCALL(m, n, lab) :: C1
    | Label _ :: RET n :: _  -> TCALL(m, n, lab) :: C
    | _                      -> CALL(m, lab) :: C
```

The effect of optimizations

New code for ex12.c

```
L1: GETBP; LDI; IFZERO L2;
    GETBP; LDI; CSTI 1; SUB; TCALL (1,1,"L1");
L2: CSTI 17; RET 1
```

- Compiling ex11.c with old and new comp.
- Finding 14200 solutions to 12-queen puzzle

	Code size (instr)	Running time (sec.)
Old direct compiler	792	9.06
New backwards compiler	701	8.00

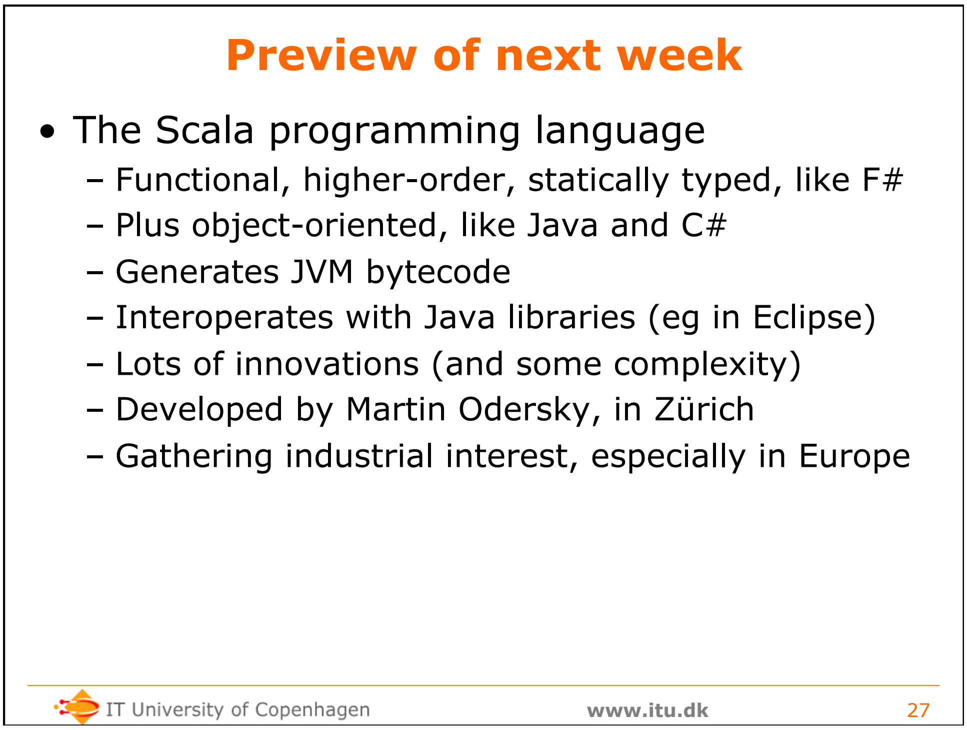
Tail call optimization

- Java does not optimize tail calls
  - And JVM does not optimize tail calls
  - The security model requires stack inspection
- C# does not optimize tail calls
  - But .NET/CLI supports tail calls

```
IL_000e:  tail.
IL_0010:  callvirt int32 MyClass::MyMethod(int32)
```

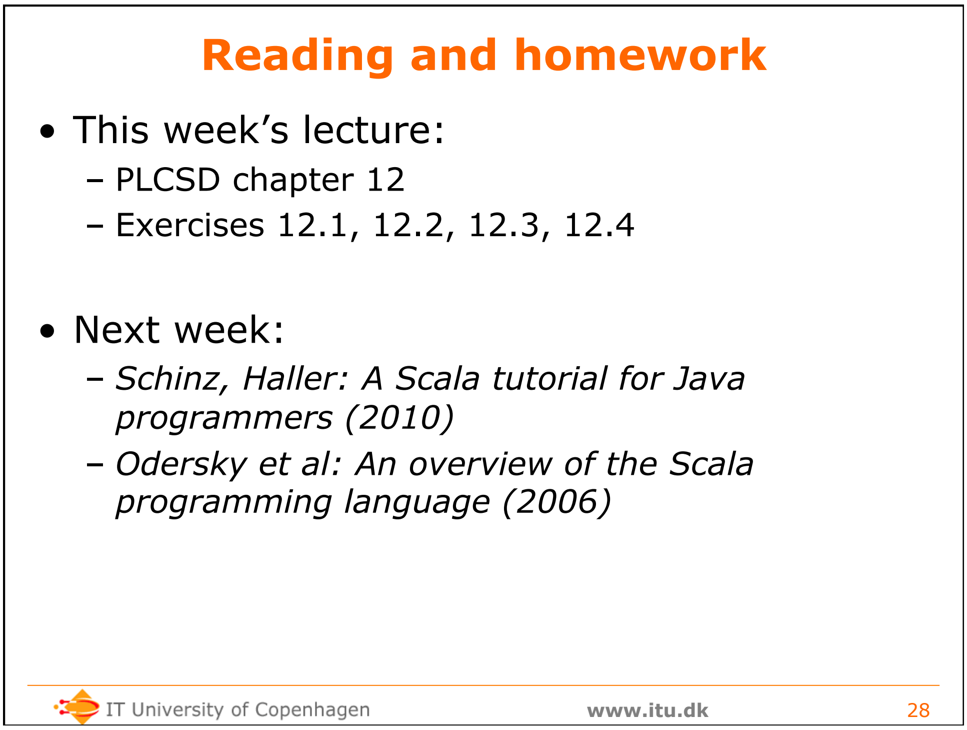
- Scheme, ML, F# … optimize tail calls
  - Needed because loops are written using tail calls

```
let rec sum xs acc =
    match xs with
      | []    -> acc
      | x::xr -> sum xr (x+acc);;
```



Preview of next week

- The Scala programming language
  - Functional, higher-order, statically typed, like F#
  - Plus object-oriented, like Java and C#
  - Generates JVM bytecode
  - Interoperates with Java libraries (eg in Eclipse)
  - Lots of innovations (and some complexity)
  - Developed by Martin Odersky, in Zürich
  - Gathering industrial interest, especially in Europe



Reading and homework

- This week's lecture:
  - PLCSD chapter 12
  - Exercises 12.1, 12.2, 12.3, 12.4

- Next week:
  - *Schinz, Haller: A Scala tutorial for Java programmers (2010)*
  - *Odersky et al: An overview of the Scala programming language (2006)*