4 Implementing an Interpreter in ML

The purpose of this lecture is to show a worked example of program development using ML modules. We shall tackle the problem of implementing a small ML system. The system is of course going to considerable simplified compared to a real ML implementation.

We implement only a few of the language constructs found in real ML. The user of our system will not get the ability to declare new types and data types; however, there will be arithmetic on the build-in integers, if...then...else expressions, and indeed lists, higher order functions and recursion, so it is far from a trivial language. We shall refer to this language as Mini ML.

Moreover, the system will be an interpreter rather than a compiler. It still has a typechecker, indeed we shall see how one can implement a restricted form of polymorphism.

The system is actually running and you can modify and extend it provided you have access to an implementation and to the files listed in Appendix B. To make life easier for you, we provide a parse functor which can parse a string (the Mini ML source expression) into an ABSTRACT SYNTAX TREE, the shape of which will be defined below. The rest of the interpreter works on abstract syntax trees.

The interpreter uses a TYPECHECKER to check the validity of input expressions and an EVALUATOR to evaluate them. Initially, the typechecker and evaluator handle only a tiny subset of Mini ML. In this lecture I shall show how one in successive steps can extend the typechecker to handle polymorphic lists, variables and let expressions. In the practical sessions you can extend the evaluator in the same manner (it is easier than extending the typechecker).

The typechecker and the evaluator can be developed independently as long as you do not change the signatures we provide. This will allow you to take the typechecker functors I have written and plug into your own system as you improve the power of your evaluator. Alternatively, you might want to modify or extend my typechecker functors, and take over evaluator functors that other people write.

The source of the bare interpreter is in Appendix A. An overview of how to run the systems is provided in Appendix B.

The development of the typechecker and the evaluator need not be in step. You can disable either by assigning false to one of the variables tc and eval.
signature INTERPRETER=
  sig
  val interpret: string -> string
  val eval: bool ref
  and tc : bool ref
  end;  

The syntax of the language is as follows

\[
\begin{align*}
\text{exp} & \ ::= \text{exp} + \text{exp} \\
& \quad \text{exp} - \text{exp} \\
& \quad \text{exp} \times \text{exp} \\
& \quad \text{true} \\
& \quad \text{false} \\
& \quad \text{exp} = \text{exp} \\
& \quad \text{if} \ \text{exp} \ \text{then} \ \text{exp} \ \text{else} \ \text{exp} \\
& \quad \text{exp} :: \exp \\
& \quad [ \ \text{exp}_1 , \ldots , \text{exp}_n \ ] \quad (n \geq 0) \\
& \quad \text{let} \ x = \text{exp} \ \text{in} \ \text{exp} \\
& \quad \text{let rec} \ x = \text{exp} \ \text{in} \ \text{exp} \\
& \quad \text{x} \\
& \quad \text{fn} \ x \ \Rightarrow \ \text{exp} \\
& \quad \text{exp} \ (\ \text{exp} \ ) \quad \text{(function application)} \\
& \quad \text{n} \quad \text{(natural numbers)} \\
& \quad (\ \text{exp} \ )
\end{align*}
\]

The abstract syntax of Mini ML is defined as a datatype in the signature EXPRESSION.

**Exercise 1** Find this signature. What is the constructor corresponding to let expressions?

We program with signatures and functors only. After the signatures, which we shall not yet study, the first functor is the interpreter itself.

**Exercise 2** Find this functor. Find the application of Ty.prType. Find its type. What do you think Ty.prType is supposed to do? What is the type of abstsyn? What do you think the evaluator is supposed to do when asked to evaluate something which has not yet been implemented?

We shall now describe Version 1, the bare typechecker, and then proceed to the extensions.
4.1 VERSION 1: The bare Typechecker (Appendix A)

The first version is just able to type check integer constants and \(+\). As signature \texttt{TYPE}
reveals, the type \texttt{Type} of types is abstract, but there are functions we can use to build basic
types and decompose them. \texttt{unTypeInt} is one of the latter; it is supposed to raise \texttt{Type}
if applied to any Mini ML type different from the int (however the type \texttt{int} is represented).
This is a common way of hiding implementation details, and it might be helpful to look
at how functor \texttt{Type} produces a structure which matches the signature \texttt{Type}.

As revealed by \texttt{signature TYPECHECKER}, the typechecker is going to depend on
the abstract syntax and a \texttt{Type} structure. However, as you can see from the declaration of
functor \texttt{TypeChecker}, all the typechecker knows about the implementation of types is
what is specified by the signature \texttt{TYPE}. This allows us to experiment with the implementa-
tion of types to obtain greater efficiency without changing the typechecker, as we shall
see in the later stages. As you see from \texttt{functor TypeChecker}, all the typechecker is
capable of handling is integer constants and \(+\).

\textbf{Exercise 3}  Modify the typechecker to handle \texttt{true}, \texttt{false}, and multiplication of integers.

Given the signature and functor declarations in Appendix A, one can build the system. First we import the parser

\begin{verbatim}
use "parser.sml";
\end{verbatim}

and then we build the system by the following declarations (which can be read from file \texttt{build1.sml}).

\begin{verbatim}
structure Expression = Expression();

structure Parser = Parser(Expression);

structure Value = Value();

structure Evaluator =
   Evaluator(structure Expression = Expression
                          structure Value = Value);

structure Ty = Type();

structure TyCh =
   TypeChecker(structure Ex = Expression
                           structure Ty = Ty);

structure Interpreter =
   Interpreter(structure Ty = Ty

\end{verbatim}

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structure Value = Value
structure Parser = Parser
structure TyCh = TyCh
structure Evaluator = Evaluator;

open Interpreter;

4.2 VERSION 2: Adding lists and polymorphism

The first extension is to implement the type checking of lists. In Version 1 the type of an expression could be inferred either directly (as in the case of true and false, or from the type of the subexpressions (as in the case of the arithmetic operations). When we introduce list, this is no longer the case. Consider for example the expression

\[
\text{if ([]} = [9]) \text{ then 5 else 7}
\]

Suppose we want to type check \([\_] = [9]\) by first type checking the left subexpression \([\_]\), then the right subexpression \([9]\) and finally checking that the left and right-hand sides are of the same type before returning the type bool. The problem now is that when we try to type check \([\_]\) we cannot know that this empty list is supposed to be an integer list. The typechecker therefore just ascribes the type 'a list to \([\_]\), where 'a is a TYPE VARIABLE. The \([9]\) of course turns out to be an int list. The typechecker now “compares” the two types 'a list and int list and discovers that they can be made the same by applying the substitution that maps 'a to int. Hence the type of the expression \([\_]\) depends not just on the expression itself, but also on the context of the expression. The context can force the type inferred for the expression to become more specific.

This “comparison” of types performed by the typechecker is called UNIFICATION and is an algebraic operation of great importance in symbolic computing. Indeed, whole-pro-gramming languages have evolved around the idea of unification (PROLOG, for example). Here is a couple of examples to illustrate how unifications works in the special case of interest, that of unifying types.

\[
[\_, [\_]]
\]

This expression is well-typed! The point is that the \([\_]\) can be regarded as an int list list. Let us see how the typechecker manages to infer the type int list list for (1). The typechecker first rewrites the expression to the equivalent:

\[
[\_] :: ((5 :: [\_]) :: [\_]) :: [\_]
\]

Checking the first argument of the topmost :: yields:

\[
[\_] : 'a1 list
\]
To check \(((5 : : \) []) : : []\)), we first check the left-hand \((5 : : []) : : []\). To check this, we first check the left-hand \(5\), for which the typechecker wisely infer the type int. Continuing to the right-hand part of \((5 : : []\) : : []\), \(\) gets the type \(\text{int} \text{ list}\). To check the :: of \((5 : : []\) : : []\), we now unify \text{int} \text{ list} and \a2 \text{ list}, which results in the substitution

\[ S_1(\a2) = \text{int}. \]

Thus the type of \((5 : : []\) : : []\) is \text{int} \text{ list}.

Returning to \(((5 : : []\) : : []\), the right-hand \(\) first gets type \a3 \text{ list} which by unification with \text{int} \text{ list} \text{ list} yields the substitution

\[ S_2(\a3) = \text{int} \text{ list}. \]

Thus the type of \(((5 : : []\) : : []\) is \text{int} \text{ list} \text{ list}.

Returning to \(((5 : : []\) : : []\), the right-hand \(\) gets the type \a4 \text{ list} which by unification with \text{int} \text{ list} \text{ list} \text{ list} yields the substitution

\[ S_3(\a4) = \text{int} \text{ list} \text{ list}. \]

Finally, returning to (2) and (3), we get to unify \a1 \text{ list} with \text{int} \text{ list} \text{ list} \text{ list}, yielding the substitution

\[ S_4(\a1) = \text{int} \text{ list} \text{ list}. \]

The type of (2), and therefore the type of (1), is thus found to be \text{int} \text{ list} \text{ list} \text{ list}.

Note that

\[ [4, [[5]]] \]

is NOT well-typed. In an attempt to compute \(S_4\), we would now be unifying \text{int} \text{ list} \text{ list} \text{ list} \text{ list} and that gives a unification error.

To implement all this, we first extend the \text{TYPE} signature and introduce a new signature, \text{UNIFY}:

\begin{verbatim}
signature TYPE =
  sig
    eqtype tyvar
    val freshTyvar: unit -> tyvar
    ...
    val mkTypeTyvar: tyvar -> Type
    and unTypeTyvar: Type -> tyvar

    val mkTypeList: Type -> Type
    and unTypeList: Type -> Type

end
\end{verbatim}
type subst
val Id: subst
   (* the identify substitution; *)
val mkSubst: tyvar*Type -> subst
   (* make singleton substitution; *)
val on : subst * Type -> Type
   (* application; *)

val prType: Type->string  (* printing *)
end

signature UNIFY=
   sig
   structure Ty: TYPE
   exception NotImplemented of string
   exception Unify
   val unify: Type.Type * Type.Type -> Type.subst
end;

functor TypeChecker
   (...  
   structure Ty: TYPE
   structure Unify: UNIFY  
   sharing Unify.Type = Ty
   )= 

struct
   infix on
   val (op on) = Ty.on
   ...

   fun tc (exp: Ex.Expression): Ty.Type =
      (case exp of
         ...  
         | Ex.LISTexpr [] =>
               let val new = Ty.freshTyvar ()
               in Ty.mkTypeList(Ty.mkTypeTyvar new)
               end
         | Ex.CONSexpr(e1,e2) =>
            ...
let val t1 = tc e1
val t2 = tc e2
val new = Ty.freshTyvar ()
val newt= Ty.mkTypeTyvar new
val t2' = Ty.mkTypeList newt
val S1 = Unify.unify(t2, t2')
    handle Unify.Unify=>
    raise Type Error(e2,"expected list type")

val S2 = Unify.unify(S1 on newt,S1 on t1)
    handle Unify.Unify=>
    raise Type Error(exp,
        "element and list have different types")
    in S2 on (S1 on t2)
end
| ... 

)handle Unify.NotImplemented msg => raise NotImplemented msg

end; (*TypeChecker*)

We also have to extend the Type functor to meet the enriched TYPE signature. The easiest way of doing this is

fun mkTypeTyvar tv = TYVAR tv
and unTypeTyvar(TYVAR tv) = tv
    | unTypeTyvar _ = raise Type

fun mkTypeList(t)=LIST t
and unTypeList(LIST t)= t
    | unTypeList(_)= raise Type

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type subst = Type -> Type

fun Id x = x

fun mkSubst(tv,ty)=
  let fun su(TYVAR tv')= if tv=tv' then ty else TYVAR tv'
       | su(INT) = INT
       | su(BOOL)= BOOL
       | su(LIST ty') = LIST (su ty')
  in su
end

fun on(S,t)= S(t)

fun prType ...
  | prType (LIST ty) = "(" ^ prType ty ^ ")list"
  | prType (TYVAR tv) = "a" ^ makestring tv
end;

Exercise 4  Extend Version 2 to handle equality. All you have to do is to fill in the relevant case in the definition of the function tc. (See appendix B about how you get the source of Version 2).

4.3 VERSION 3: A different implementation of types

Version 3 arises from Version 2 by replacing the Type functor by a different implementation of types. The idea is that instead of having substitutions as functions, we can implement type variables by references (pointers) and then do substitutions directly by assignments.

In case you have not seen the reserved word withtype before, withtype is used to declare a type abbreviation locally within a datatype declaration.

functor ImpType():TYPE =
struct
  datatype 'a option = NONE | SOME of 'a

  datatype Type = INT
               | BOOL
               | LIST of Type
               | TYVAR of tyvar

  withtype tyvar = Type option ref
type tyvar = Type option ref

fun freshTyvar() = ref (NONE)

exception Type

fun mkTypeInt() = INT
and unTypeInt(INT)=()
| ...
| unTypeInt(TYVAR(ref (SOME t)))= unTypeInt t
| unTypeInt _ = raise Type

...

type subst = unit

val Id = ();

exception MkSubst;

fun mkSubst(tv,ty)=
    case tv of
    ref(NONE) => tv:= (SOME ty)
    | ref(SOME t) => raise MkSubst

fun on(S,t)= t

fun prType ...
| prType (TYVAR (ref NONE)) = "a?"
| prType (TYVAR (ref (SOME t))) = prType t
end;

We can now build two systems at the same time and compare the efficiency of the two implementations. The nice thing is that we do not have to modify the typechecker functor at all, nor do we even have to modify the unification functor; we can just extend the final sequence of structure declarations to use both implementations of types.

Exercise 5 When I did this, I found (to my surprise), that the functional version in some cases was twice as fast, and never slower than the imperative variant. The relative performance of the two vary greatly from expression to expression. Can you find an expression for which the imperative version really is faster? (See Appendix B for how to get hold of the source of Version 3). Be careful with generating very demanding tasks for the ML system; you can make it crash!
ML implementors normally opt for the imperative version. In all fairness, the above comparison ignores that composing substitutions is much easier in the imperative version than it is in the applicative version; in the fragment of Mini ML considered so far, we have not had to compose substitutions.

One should not be too concerned with performance issues at too early a stage. It can be surprisingly difficult to predict where efficiency is most needed, and it is much more important, at first, to get the overall structure of the system right. It was important, for example, that we did NOT make the constructors of the datatype Type visible in the signature TYPE, and that we wrote the unification algorithm in a way which does not use the internal structure of Type. Had we not done this, we would not have been able to switch from one implementation to another that easily, and therefore chances are that we would chosen the imperative one, assuming that it was the more efficient one, without ever trying the “obvious” applicative implementation.

4.4 VERSION 4: Introducing variables and let

We now extend Version 3 by implementing the type checking of let expressions and of identifiers.

The typechecker function tc now has to take TWO arguments,

\[ tc(TE, e) \]

where \( e \) is an expression and \( TE \) is a TYPE ENVIRONMENT, which maps variables occurring free in \( e \) to TYPE SCHEMES. The definition of what a type scheme is will be given below; for now it suffices to know that every type can be regarded as a type scheme.

To take an example, if \( TE \) maps \( x \) to \( \text{int} \) and \( y \) to \( \text{int} \), then \( tc \) will deduce the type \( \text{int} \) for the expression \( x + y \). (However, if \( TE \) mapped \( y \) to \( \text{bool} \), there would be a type error.)

The fact that we can bind variables to expressions whose types have been inferred to contain type variables means that we get type variables in the type environment. For instance, to type check

\[
\text{let } x = [] \text{ in } 4 :: x \text{ end}
\]

we first check \( [] \) yielding the type \( 'a1 \text{ list} \), say. Then we bind \( x \) to the type scheme \( \forall 'a1. 'a1 \text{ list} \). Here the binding \( \forall 'a1 \) of \( 'a1 \) indicates that when we look up the the type of \( x \) in the type environment, we return a type obtained from the type scheme \( \forall 'a1. 'a1 \text{ list} \) by instantiating the bound variables (here just \( 'a1 \)) by fresh type variables. In our example, when we look up \( x \) in the type environment during the checking of \( 4 :: x \), we instantiate \( 'a1 \) to a fresh type variable \( 'a2 \), say, yielding the type \( 'a2 \text{ list} \) for \( x \). Thus we get to unify \( \text{int list} \) against \( 'a2 \text{ list} \), yielding the substitution of \( \text{int} \) for \( 'a2 \).

Throughout the body of the let, \( x \) will be bound to \( \forall 'a1. 'a1 \text{ list} \) in the type environment. Since we take a fresh instance of this type scheme each time we look up \( x \), we can use \( x \) both as an \( \text{int list} \) and as an \( \text{int list list} \), say:

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let x = [] in (4::x)::x end

Exercise 6  Assuming that you instantiate the bound 'a1 to 'a3 when you meet the last occurrence of x, what two types should be unified, and what is the resulting substitution on 'a3?

The variable x is an example of POLYMORPHISM: after x has been declared, an occurrence of x can potentially be given infinitely many types: int list, bool list, int list list, and so on, all captured by the type scheme ∀'a1.'a1 list. In ML, a TYPE SCHEME always takes the form ∀α₁···αₙ.τ, (n ≥ 0), where α₁, ..., αₙ are type variables and τ is a type. In the fragment of Mini ML considered so far, it will always be the case that any type variable occurring in τ is amongst the α₁, ..., αₙ, but when one introduces functions and application, this no longer is the case.

Here is how we implement variables and let. We first extend the TYPE signature:

|signature TYPE =
|  sig
|    ...
|    type TypeScheme
|  val instance: TypeScheme -> Type
|  val close: Type -> Type
|end

Version 1 (Appendix A) already contains a signature for environments (find it). It was actually intended for the practical where you need it to extend the evaluator, but we can make use of it to implement type environments. The signature of the typechecker can be left unchanged, but we need to change the functor that builds the typechecker by including the environment management among the formal parameters:

|functor TypeChecker
|  (structure Ex: EXPRESSION
|   structure Ty: TYPE
|   structure Unify: UNIFY
|       sharing Unify.Type = Ty
|   structure TE: ENVIRONMENT
|)
|struct
|  infix on
|  val (op on) = Ty.on
|  structure Exp = Ex
|  structure Type = Ty
fun tc (TE: Ty.TypeScheme TE.Environment, exp: Ex.Expression): Ty.Type =
  (case exp of
    Ex.BOOLexpr b => Ty.mkTypeBool()
  | Ex.NUMBERexpr _ => Ty.mkTypeInt()
  | Ex.SUMexpr(e1,e2) => checkIntBin(TE,e1,e2)
  | Ex.DIFFexpr(e1,e2) => checkIntBin(TE,e1,e2)
  | Ex.PRODexpr(e1,e2) => checkIntBin(TE,e1,e2)
  | Ex.LISTexpr [] =>
    let val new = Ty.freshTyvar ()
    in Ty.mkTypeList(Ty.mkTypeTyvar new)
    end
  | Ex.LISTexpr(e::es) =>
    let val t1 = tc(TE, e1)
    val t2 = tc(TE, e2)
    val new = Ty.freshTyvar ()
    val newt= Ty.mkTypeTyvar new
    val t2' = Ty.mkTypeList newt
    val S1 = Unify.unify(t2, t2')
    handle Unify.Unify=>
      raise TypeError(e2,"expected list type")
    end
  | Ex.EQexpr _ => raise Not Implemented "(equality)"
  | Ex.CONDexpr _ => raise Not Implemented "(conditional)"
  | Ex.DECLexpr(x,e1,e2) =>
    let val t1 = tc(TE,e1);
    val typeScheme = Ty.close(t1)
    in tc(TE.declare(x,typeScheme,TE), e2)
    end
  | Ex.RECDECLexpr _ => raise Not Implemented "(rec decl)"
  | Ex.IDENTexpr x =>
    (Ty.instance(TE.retrieve(x,TE))
    handle TE.Retrieve _=>
      raise TypeError(exp,"identifier " ^ x ^ " not declared")
    end
  | Ex.LAMBDAPERMS _ => raise Not Implemented "(function)"

exception NotImplementedException of string
exception TypeError of Ex.Expression * string
| Ex.APPLexpr _ => raise NotImplementedException  "(application)"

)handle Unify.NotImplemented msg => raise Unimplemented msg

and checkIntBin(TE,e1,e2) =
let val t1 = tc(TE,e1)
  val _  = Ty.unTypeInt t1
        handle Ty.Type => raise TypeError(e1,"expected int")
val t2 = tc(TE,e2)
val _  = Ty.unTypeInt t2
        handle Ty.Type => raise TypeError(e2,"expected int")
in Ty.mkTypeInt()
end;

fun typecheck(e) = tc(TE.emptyEnv,e)
end; (*TypeChecker*)

Then we extend the Type functor to match the TYPE signature:

functor Type():TYPE =
struct
  ...
  datatype TypeScheme = FORALL of tyvar list * Type

  fun instance(FORALL(tyvars,ty))=
    let val old_to_new_tyvars = map (fn tv=>(tv,freshTyvar())) tyvars
      exception Find;
      fun find(tv,[])= raise Find
         | find(tv,(tv',new_tv)::rest)=
            if tv=tv' then new_tv else find(tv,rest)
    fun ty_instance INT = INT
        | ty_instance BOOL = BOOL
        | ty_instance (LIST t) = LIST(ty_instance t)
        | ty_instance (TYVAR tv) =
          TYVAR(find(tv,old_to_new_tyvars)
             handle Find=> tv)

    in
    ty_instance ty
    end

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fun close(ty)= 
let fun fv(INT) = [] 
    | fv(BOOL) = [] 
    | fv(LIST t) = fv(t) 
    | fv(TYVAR tv) = [tv] 
  in FORALL(fv ty,ty) 
end;

Finally, the system is re-built as in Version 2, except that we have to provide and link in an Environment functor which matches ENVIRONMENT.

**Exercise 7** Extend Version 4 with if .. then .. else. (This extension has no subtle implications for the type checking.)

**Exercise 8** [For the extra keen] Extend Version 4 to cope with lambda abstraction (fn) and application. First, you have to introduce arrow types with constructors and destructors. Then you have to change the type of close so that it takes two arguments, namely a type environment and a type. It should return the type scheme that is obtained by quantifying on all the type variables that occur in the type but do not occur free in the type environment.

Then you can modify the type checker. When you type check a lambda abstraction, you just bind the formal parameter to the trivial type scheme which is just a fresh type variable (no quantified variables). Thus the type environment can now contain type schemes with free type variables.

An application tc(TE,e) now yields two arguments, namely a type t and a substitution S; the idea is that if you apply the substitution S to the type environment TE, which now can contain free type variables, the expression e has the type t. When an expression consists of more than one subexpression, the type environment gradually becomes more and more specific by applying the substitutions produced by the checking of the subexpressions one by one. Moreover, the substitution returned from the whole expression is the composition of these individual substitutions. (You have to extend the TYPE signature (and the Type functor) with composition of substitutions.)

Finally, you can extend the unification algorithm to cope with arrow types. (This will also use composition of substitutions.)

### 4.5 Acknowledgement
The parser and evaluator and all the signatures related to them are due to Nick Rothwell.
Appendix A: The bare Interpreter

(* interp1.sml : VERSION 1: the bare interpreter *)

signature INTERPRETER=
  sig
    val interpret: string -> string
    val eval: bool ref
    and tc : bool ref
  end;

  (* syntax *)

signature EXPRESSION =
  sig
    datatype Expression =
        SUMexpr of Expression * Expression |
        DIFFexpr of Expression * Expression |
        PRODexpr of Expression * Expression |
        BOOLexpr of bool |
        EQexpr of Expression * Expression |
        CONDexpr of Expression * Expression * Expression |
        CONSexpr of Expression * Expression |
        LISTexpr of Expression list |
        RECDECLexpr of string * Expression * Expression |
        IDENTexpr of string |
        LAMBDArexp of string * Expression |
        APPLexpr of Expression * Expression |
        NUMBERexpr of int
  end

  (* parsing *)

signature PARSER =
  sig
    structure E: EXPRESSION

    exception Lexical of string
    exception Syntax of string

  end

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val parse: string -> E.Expression
end

(* environments *)

signature ENVIRONMENT =
  sig
  type 'object Environment

  exception Retrieve of string

  val emptyEnv: 'object Environment
  val declare: string * 'object * 'object Environment
  -> 'object Environment
  val retrieve: string * 'object Environment -> 'object

  end

(* evaluation *)

signature VALUE =
  sig

  type Value

  exception Value

  val mkValueNumber: int -> Value
  and unValueNumber: Value -> int

  val mkValueBool: bool -> Value
  and unValueBool: Value -> bool

  val ValueNil: Value
  val mkValueCons: Value * Value -> Value
  and unValueHead: Value -> Value
  and unValueTail: Value -> Value

  val eqValue: Value * Value -> bool
  val printValue: Value -> string

  end

signature EVALUATOR =
  sig


structure Exp: EXPRESSION
structure Val: VALUE
exception Unimplemented
val evaluate: Exp.Expression -> Val.Value
end

(* type checking *)
signature TYPE =
sig
type Type

(*constructors and decstructors*)
exception Type
val mkTypeInt: unit -> Type
 and unTypeInt: Type -> unit

val mkTypeBool: unit -> Type
 and unTypeBool: Type -> unit

val prType: Type->string
end

signature TYPECHECKER =
sig
 structure Exp: EXPRESSION
 structure Type: TYPE
 exception NotImplemented of string
 exception TypeError of Exp.Expression * string
 val typecheck: Exp.Expression -> Type.Type
end;

(* the interpreter*)
functor Interpreter
(structure Ty: TYPE
 structure Value : VALUE
 structure Parser: PARSER
 structure TyCh: TYPECHECKER
 structure Evaluator:EVALUATOR
 sharing Parser.E = TyCh.Exp = Evaluator.Exp
 and TyCh.Type = Ty

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and Evaluator.Val = Value

): INTERPRETER=

struct
val eval = ref true (* toggle for evaluation *)
and tc = ref true (* toggle for type checking *)
fun interpret(str)=

let val abstsyn= Parser.parse str
val typestr= if !tc then
    Ty.prType(TyCh.typecheck abstsyn)
else "(disabled)"

val valuestr= if !eval then
    Value.printValue(Evaluator.evaluate abstsyn)
else "(disabled)"

in   valuestr ^ " : " ^ typestr
end

handle Evaluator.Unimplemented =>
    "Evaluator not fully implemented"
| TyCh.NotImplemented msg =>
    "Typechecker not fully implemented " ^ msg
| Value.Value => "Run-time error"
| Parser.Syntax msg => "Syntax Error: " ^ msg
| Parser.Lexical msg=> "Lexical Error: " ^ msg
| TyCh.TypeError(_,msg)=> "Type Error: " ^ msg
end;

(* the evaluator *)

functor Evaluator
(structure Expression: EXPRESSION
structure Value: VALUE):EVALUATOR=

struct
structure Exp= Expression
structure Val= Value
exception Unimplemented

local
    open Expression Value
    fun evaluate exp =
        case exp
          of BOOLexpr b => mkValueBool b
| NUMBERexpr i => mkValueNumber i
| SUMexpr(e1, e2) =>
|     let val e1' = evaluate e1
|         val e2' = evaluate e2
|     in
|         mkValueNumber(unValueNumber e1' +
|                         unValueNumber e2')
|     end
| DIFFexpr(e1, e2) =>
|     let val e1' = evaluate e1
|         val e2' = evaluate e2
|     in
|         mkValueNumber(unValueNumber e1' -
|                         unValueNumber e2')
|     end
| PRODexpr(e1, e2) =>
|     let val e1' = evaluate e1
|         val e2' = evaluate e2
|     in
|         mkValueNumber(unValueNumber e1' *
|                         unValueNumber e2')
|     end
| EExpr _ => raise Unimplemented
| CONDexpr _ => raise Unimplemented
| CONSexpr _ => raise Unimplemented
| LISTexpr _ => raise Unimplemented
| DECLexpr _ => raise Unimplemented
| RECDECLExpr _ => raise Unimplemented
| IDENTexpr _ => raise Unimplemented
| LAMBDAexpr _ => raise Unimplemented
| APPLexpr _ => raise Unimplemented

val evaluate = evaluate
end
end;

(* the typechecker *)

functor TypeChecker
(structure Ex: EXPRESSION
structure Ty: TYPE)=

struct
structure Exp = Ex
structure Type = Ty

exception NotImplemented of string
exception TypeError of Ex.Expression * string

fun tc (exp: Ex.Expression): Ty.Type =
case exp of
  Ex.BOOLexpr _ => raise Not Implemented "(boolean constants)"
| Ex.NUMBERexpr _ => Ty.mkTypeInt()
| Ex.SUMexpr(e1,e2) => checkIntBin(e1,e2)
| Ex.DIFFexpr _ => raise Not Implemented "(minus)"
| Ex.PRODexpr _ => raise Not Implemented "(product)"
| Ex.LISTexpr _ => raise Not Implemented "(lists)"
| Ex.CONSEXpr _ => raise Not Implemented "(lists)"
| Ex.EQexpr _ => raise Not Implemented "(equality)"
| Ex.CONDexpr _ => raise Not Implemented "(conditional)"
| Ex.DECLexpr _ => raise Not Implemented "(declaration)"
| Ex.RECDECLexpr _ => raise Not Implemented "(rec decl)"
| Ex.IDENTexpr _ => raise Not Implemented "(identifier)"
| Ex.LAMBDAexpr _ => raise Not Implemented "(function)"
| ExAPPLExpr _ => raise Not Implemented "(application)"

and checkIntBin(e1,e2) =
  let val t1 = tc e1
  val _ = Ty.unTypeInt t1
    handle Ty.Type=>
    raise TypeError(e1,"expected int")
  val t2 = tc e2
  val _ = Ty.unTypeInt t2
    handle Ty.Type=>
    raise TypeError(e2,"expected int")
in Ty.mkTypeInt()
end;

val typecheck = tc

end; (*TypeChecker*)
functor Type():TYPE = 
  struct
    datatype Type = INT
       | BOOL

  exception Type

  fun mkTypeInt() = INT
  and unTypeInt(INT)=()
       | unTypeInt(_)= raise Type

  fun mkTypeBool() = BOOL
  and unTypeBool(BOOL)=()
       | unTypeBool(_)= raise Type

  fun prType INT = "int"
  | prType BOOL= "bool"
end;

functor Expression(): EXPRESSION = 
  struct
    type 'a pair = 'a * 'a

    datatype Expression =
        SUMexpr of Expression pair  
        | DIFFexpr of Expression pair  
        | PRODexpr of Expression pair  
        | BOOLexpr of bool
        | EQexpr of Expression pair  
        | CONDexpr of Expression * Expression * Expression  
        | CONSexpr of Expression pair  
        | LISTexpr of Expression list  
        | DECLEXPR of string * Expression * Expression  
        | RECDDECLexpr of string * Expression * Expression  
        | IDENTexpr of string  
        | LAMBDAXPR of string * Expression  
        | APPLexpr of Expression * Expression  
        | NUMBEREXPR of int
end;

functor Value(): VALUE =
  struct
    type 'a pair = 'a * 'a

datatype Value = NUMBERvalue of int |
                  BOOLvalue of bool |
                  NILvalue |
                  CONSvalue of Value pair

exception Value

val mkValueNumber = NUMBERvalue
val mkValueBool = BOOLvalue

val ValueNil = NILvalue
val mkValueCons = CONSvalue

fun unValueNumber(NUMBERvalue(i)) = i |
    unValueNumber(_) = raise Value

fun unValueBool(BOOLvalue(b)) = b |
    unValueBool(_) = raise Value

fun unValueHead(CONSvalue(c, _)) = c |
    unValueHead(_) = raise Value

fun unValueTail(CONSvalue(_, c)) = c |
    unValueTail(_) = raise Value

fun eqValue(c1, c2) = (c1 = c2)

(* Pretty-printing *)
fun printValue(NUMBERvalue(i)) = makestring(i) |
    printValue(BOOLvalue(true)) = "true" |
    printValue(BOOLvalue(false)) = "false" |
    printValue(NILvalue) = "[]" |
    printValue(CONSvalue(cons)) = "[" ^
        printValueList(cons) ^ "]

and printValueList(hd, NILvalue) = printValue(hd) |
    printValueList(hd, CONSvalue(tl)) =
        printValue(hd) ^ ", " ^ printValueList(tl) |
printValueList(_) = raise Value
end;
Appendix B: Files

The following files are available in the directory /usr/cheops/mads/course

- interp1.sml  Version 1 (as included in Appendix A).
- interp2.sml ... interp4.sml  The other versions.
- build1.sml  the structure declarations needed to build Version 1.
- build2.sml ... build4.sml  Similarly for the other versions.
- parser.sml  The parser functor.

To build Version 3, say, you type the following (assuming you have copied the files to your directory):

    use "interp3.sml";
    use "parser.sml";
    use "build3.sml";

Since the parser functor is completely closed, you don't have to include it more than once in every session, although you will probably want to build your system several times while you experiment with the extensions.