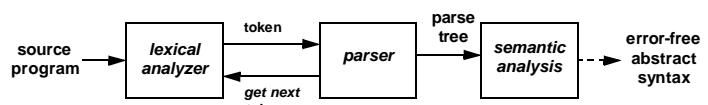


Compiler Front-End

- Almost all compilers and interpreters contain the same **front-end** --- it consists of three components:
 - Lexical Analysis** --- report lexical errors, output a list of tokens
 - Syntax Analysis** --- report syntactic errors, output a parse tree
 - Semantic Analysis** --- report semantic errors (e.g., type-errors, undefined identifiers, ...) --- generate a clean and error-free "abstract syntax tree"



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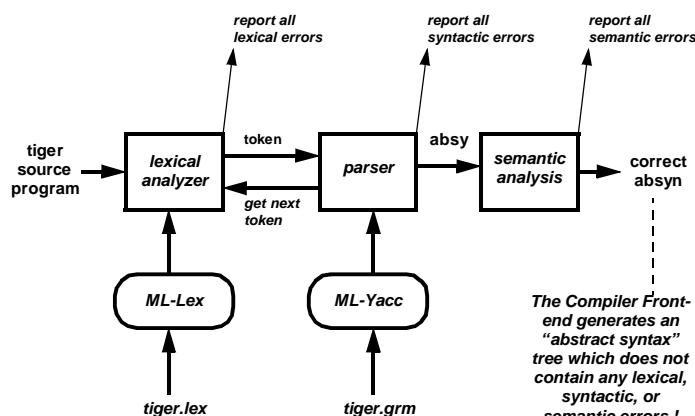
"Concrete" vs. "Abstract" Syntax

- The grammar specified in "tiger.grm" (for Yacc) is mainly used for parsing only ----- the key is to resolve all ambiguities. This grammar is called **Concrete Syntax**.
- Abstract Syntax (Absyn)** is used to characterize the essential structure of the program ----- the key is to be as simple as possible; **Absyn** may contain ambiguities.
- The grammar for **Abstract Syntax** is defined using **ML datatypes**.
- Traditional Compilers:** do semantic analysis on Concrete Syntax --- implemented as "actions" in Section 3 of "tiger.grm" file (for Yacc)
- Modern Compilers:** "tiger.grm" constructs the Abstract Syntax tree; the semantic analysis is performed on the Absyn later after parsing!

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Tiger Compiler Front End



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Tiger Program and Expression

- A *Tiger program* prog is just an expression exp
- An expression can be any of the following:

I-value	foo, foo.bar, foo[1]
Nil	nil
Integer literal	34
String literal	"Hello, World\n"
Sequencing	(exp; exp; ...; exp)
Function call	id(), id(exp{,exp})
Arithmetic expression	exp arith-op exp
Comparison expression	exp comp-op exp
Boolean operators	exp & exp, exp exp
Record creation	ty-id {id = exp, ...}, {}
Array creation	ty-id [exp ₁] of exp ₂
Assignment	lvalue := exp

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Tiger Expression and Declaration

- More Tiger expressions:

If-then-else	<code>if exp₁ then exp₂ else exp₃</code>
If-then	<code>if exp₁ then exp₂</code>
While-expression	<code>while exp₁ do exp₂</code>
For-expression	<code>for id:=exp₁ to exp₂ do exp₃</code>
Break-expression	<code>break</code>
Let-expression	<code>let decsq in {exp} end</code>

- A Tiger declaration sequence is a sequence of type, variable, and function declarations:

```
dec -> tydec | vardec | fundec
decsq -> decsq dec | ?
```

Tiger Type Declaration

- Tiger Type declarations:

```
tydec -> type id = ty
ty -> id | { tyfields } | array of id
tyfields -> ? | id : type-id { ,id: type-id }
```

- You can define mutually-recursive types using a **consecutive sequence** of type declarations

```
type tree = {key : int, children : treelist}
type treelist = {hd : tree, tl : treelist}
```

recursion cycle must pass through a record or array type !

Variable and Function Declaration

- Tiger Variable declarations:

```
short-form: vardec -> var id := exp
long-form: vardec -> var id : type-id := exp

"var x := 3" in Toy is equivalent to "val x = ref 3" in ML
```

- Tiger Function declarations:

```
procedure: fundec -> function id (tyfields) := exp
function: fundec -> function id (tyfields):type-id := exp
```

- Function declarations may be mutually recursive --- must be declared in a **sequence of consecutive** function declarations! Variable declarations cannot be mutually recursive !

Tiger Absyn “Hack”

- When translating from **Concrete Syntax** to **Abstract Syntax**, we can do certain syntactic transformations

```
MINUS exp ===> 0 MINUS exp
exp1 & exp2 ===> if exp1 then exp2 else 0
exp1 | exp2 ===> if exp1 then 1 else exp2
```

This can make **Abstract Syntax** even simpler.

Toy does not support Macros. If the source language supports macros, they can be processed here.

Tiger Semantics

- **nil** --- a value belong to every record type.
- Scope rule --- similar to PASCAL, Algol ---- support nested scope for types, variables, and functions; redeclaration will hide the same name.

```
function f(v : int) =
  let var v := 6
    in print(v);
    let var v := 7 in print(v) end;
    print(v);
    let var v := 8 in print(v) end;
    print(v)
  end
```

- Support two different **name space**: one for types, and one for functions and variables. You can have a type called `foo` and a variable `foo` in scope at same time.

An Example

```
(* A program to solve the 8-queens problem, see Appel's book *)

let
  var N := 8

  type intArray = array of int
  var row := intArray [ N ] of 0
  var col := intArray [ N ] of 0
  var diag1 := intArray [N+N-1] of 0
  var diag2 := intArray [N+N-1] of 0

  function printboard() =
    (for i := 0 to N-1
      do (for j := 0 to N-1
          do print(if col[i]=j then " O" else " .");
          print("\n"));
      print("\n"))

  function try(c:int) =
    (* for i:= 0 to c do print("."); print("\n"); flush(); *)
    if c=N then printboard()
    else for r := 0 to N-1
      do if row[r]=0 & diag1[r+c]=0 & diag2[r+7-c]=0
        then (row[r]:=1; diag1[r+c]:=1; diag2[r+7-c]:=1;
              col[c]:=r; try(c+1));
              row[r]:=0; diag1[r+c]:=0; diag2[r+7-c]:=0)

  in try(0)
end
```